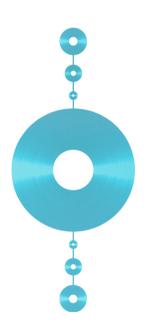
IAR C/C++ Compiler

Compiling and Linking

for Advanced RISC Machines Ltd's **ARM Cores**





DARM-9

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Preface

Welcome to the IAR C/C++ Development Guide for ARM. The purpose of this guide is to provide you with detailed reference information that can help you to use the build tools to best suit your application requirements. This guide also gives you suggestions on coding techniques so that you can develop applications with maximum efficiency.

Who should read this guide

Read this guide if you plan to develop an application using the C or C++ language for the ARM core and need detailed reference information on how to use the build tools. You should have working knowledge of:

- The architecture and instruction set of the ARM core. Refer to the documentation from Advanced RISC Machines Ltd for information about the ARM core
- The C or C++ programming language
- · Application development for embedded systems
- The operating system of your host computer.

How to use this guide

When you start using the IAR C/C++ compiler and linker for ARM, you should read *Part 1. Using the build tools* in this guide.

When you are familiar with the compiler and linker and have already configured your project, you can focus more on *Part 2. Reference information*.

If you are new to using the IAR Systems build tools, we recommend that you first study the *IDE Project Management and Building Guide for ARM*. This guide contains a product overview, conceptual and user information about the IDE and the IAR C-SPY® Debugger, and corresponding reference information.

What this guide contains

Below is a brief outline and summary of the chapters in this guide.

PART I. USING THE BUILD TOOLS

- Introduction to the IAR build tools gives an introduction to the IAR build tools, which includes an overview of the tools, the programming languages, the available device support, and extensions provided for supporting specific features of the ARM core.
- *Developing embedded applications* gives the information you need to get started developing your embedded software using the IAR build tools.
- Data storage describes how to store data in memory.
- *Functions* gives a brief overview of function-related extensions—mechanisms for controlling functions—and describes some of these mechanisms in more detail.
- *Linking using ILINK* describes the linking process using the IAR ILINK Linker and the related concepts.
- *Linking your application* lists aspects that you must consider when linking your application, including using ILINK options and tailoring the linker configuration file.
- *The DLIB runtime environment* describes the DLIB runtime environment in which an application executes. It covers how you can modify it by setting options, overriding default library modules, or building your own library. The chapter also describes system initialization introducing the file cstartup, how to use modules for locale, and file I/O.
- Assembler language interface contains information required when parts of an application are written in assembler language. This includes the calling convention.
- Using C gives an overview of the two supported variants of the C language and an overview of the compiler extensions, such as extensions to Standard C.
- Using C++ gives an overview of the two levels of C++ support: The industry-standard EC++ and IAR Extended EC++.
- *Application-related considerations* discusses a selected range of application issues related to using the compiler and linker.
- *Efficient coding for embedded applications* gives hints about how to write code that compiles to efficient code for an embedded application.

PART 2. REFERENCE INFORMATION

- *External interface details* provides reference information about how the compiler and linker interact with their environment—the invocation syntax, methods for passing options to the compiler and linker, environment variables, the include file search procedure, and the different types of compiler and linker output. The chapter also describes how the diagnostic system works.
- *Compiler options* explains how to set options, gives a summary of the options, and contains detailed reference information for each compiler option.

- *Linker options* gives a summary of the options, and contains detailed reference information for each linker option.
- *Data representation* describes the available data types, pointers, and structure types. This chapter also gives information about type and object attributes.
- *Extended keywords* gives reference information about each of the ARM-specific keywords that are extensions to the standard C/C++ language.
- Pragma directives gives reference information about the pragma directives.
- *Intrinsic functions* gives reference information about functions to use for accessing ARM-specific low-level features.
- *The preprocessor* gives a brief overview of the preprocessor, including reference information about the different preprocessor directives, symbols, and other related information.
- *Library functions* gives an introduction to the C or C++ library functions, and summarizes the header files.
- *The linker configuration file* describes the purpose of the linker configuration file and describes its contents.
- Section reference gives reference information about the use of sections.
- *Stack usage control files* describes the syntax and semantics of stack usage control files.
- *IAR utilities* describes the IAR utilities that handle the ELF and DWARF object formats.
- *Implementation-defined behavior for Standard C* describes how the compiler handles the implementation-defined areas of Standard C.
- *Implementation-defined behavior for C89* describes how the compiler handles the implementation-defined areas of the C language standard C89.

Other documentation

User documentation is available as hypertext PDFs and as a context-sensitive online help system in HTML format. You can access the documentation from the Information Center or from the **Help** menu in the IAR Embedded Workbench IDE. The online help system is also available via the F1 key.

USER AND REFERENCE GUIDES

The complete set of IAR Systems development tools is described in a series of guides. For information about:

- System requirements and information about how to install and register the IAR Systems products, refer to the booklet Quick Reference (available in the product box) and the *Installation and Licensing Guide*.
- Getting started using IAR Embedded Workbench and the tools it provides, see the guide *Getting Started with IAR Embedded Workbench*®.
- Using the IDE for project management and building, see the *IDE Project* Management and Building Guide for ARM.
- Using the IAR C-SPY® Debugger, see the C-SPY® Debugging Guide for ARM.
- Programming for the IAR Assembler for ARM, see the *IAR Assembler Reference Guide for ARM*.
- Using the IAR DLIB Library, see the *DLIB Library Reference information*, available in the online help system.
- Porting application code and projects created with a previous version of the IAR Embedded Workbench for ARM, see the *IAR Embedded Workbench*® *Migration Guide*.
- Developing safety-critical applications using the MISRA C guidelines, see the *IAR Embedded Workbench*® *MISRA C:2004 Reference Guide* or the *IAR Embedded Workbench*® *MISRA C:1998 Reference Guide*.

Note: Additional documentation might be available depending on your product installation.

THE ONLINE HELP SYSTEM

The context-sensitive online help contains:

- Information about debugging using the IAR C-SPY® Debugger
- Information about using the editor
- Reference information about the menus, windows, and dialog boxes in the IDE
- Compiler reference information
- Keyword reference information for the DLIB library functions. To obtain reference information for a function, select the function name in the editor window and press F1.

FURTHER READING

These books might be of interest to you when using the IAR Systems development tools:

- Seal, David, and David Jagger. ARM Architecture Reference Manual. Addison-Wesley
- Barr, Michael, and Andy Oram, ed. *Programming Embedded Systems in C and C++*. O'Reilly & Associates.
- Furber, Steve, ARM System-on-Chip Architecture. Addison-Wesley.
- Harbison, Samuel P. and Guy L. Steele (contributor). *C: A Reference Manual*. Prentice Hall.
- Josuttis, Nicolai M. *The C++ Standard Library: A Tutorial and Reference*. Addison-Wesley.
- Kernighan, Brian W. and Dennis M. Ritchie. *The C Programming Language*. *Prentice Hall*.
- Labrosse, Jean J. Embedded Systems Building Blocks: Complete and Ready-To-Use Modules in C. R&D Books.
- Lippman, Stanley B. and Josée Lajoie. C++ Primer. Addison-Wesley.
- Mann, Bernhard. C für Mikrocontroller. Franzis-Verlag. [Written in German.]
- Meyers, Scott. *Effective* C++: 50 Specific Ways to Improve Your Programs and Designs. Addison-Wesley.
- Meyers, Scott. *More Effective C++*. Addison-Wesley.
- Meyers, Scott. Effective STL. Addison-Wesley.
- Sloss, Andrew N. et al, ARM System Developer's Guide: Designing and Optimizing System Software. Morgan Kaufmann.
- Stroustrup, Bjarne. *The C++ Programming Language*. Addison-Wesley.
- Stroustrup, Bjarne. *Programming Principles and Practice Using C++*. Addison-Wesley.
- Sutter, Herb. Exceptional C++: 47 Engineering Puzzles, Programming Problems, and Solutions. Addison-Wesley.

WEB SITES

Recommended web sites:

- The Advanced RISC Machines Ltd web site, **www.arm.com**, that contains information and news about the ARM cores.
- The IAR Systems web site, **www.iar.com**, that holds application notes and other product information.

- The web site of the C standardization working group, www.open-std.org/jtc1/sc22/wg14.
- The web site of the C++ Standards Committee, www.open-std.org/jtc1/sc22/wg21.
- Finally, the Embedded C++ Technical Committee web site, www.caravan.net/ec2plus, that contains information about the Embedded C++ standard.

Document conventions

When, in this text, we refer to the programming language C, the text also applies to C++, unless otherwise stated.

When referring to a directory in your product installation, for example arm\doc, the full path to the location is assumed, for example c:\Program Files\IAR Systems\Embedded Workbench 6.n\arm\doc.

TYPOGRAPHIC CONVENTIONS

This guide uses the following typographic conventions:

Style	Used for	
computer	 Source code examples and file paths. Text on the command line. Binary, hexadecimal, and octal numbers. 	
parameter	A placeholder for an actual value used as a parameter, for example <i>filename</i> .h where <i>filename</i> represents the name of the file.	
[option]	An optional part of a linker directive, inline assembler statement, or a stack usage control directive, [and] are not part of the actual directive, but any [,], {, or } are part of the directive syntax.	
{option}	A mandatory part of a linker directive, $\{$ and $\}$ are not part of the actual directive, but any $[,], \{, or \}$ are part of the directive syntax	
[option]	An optional part of a command.	
[a b c]	An optional part of a command with alternatives.	
{a b c}	A mandatory part of a command with alternatives.	
bold	Names of menus, menu commands, buttons, and dialog boxes that appear on the screen.	
italic	 A cross-reference within this guide or to another guide. Emphasis. 	

Table 1: Typographic conventions used in this guide

Style	Used for	
	An ellipsis indicates that the previous item can be repeated an arbitrary number of times.	
X	Identifies instructions specific to the IAR Embedded Workbench $\ensuremath{\mathbb S}$ IDE interface.	
	Identifies instructions specific to the command line interface.	
	Identifies helpful tips and programming hints.	
<u>.</u>	Identifies warnings.	

Table 1: Typographic conventions used in this guide (Continued)

NAMING CONVENTIONS

The following naming conventions are used for the products and tools from IAR Systems® referred to in this guide:

Brand name	Generic term
IAR Embedded Workbench® for ARM	IAR Embedded Workbench®
IAR Embedded Workbench® IDE for ARM	the IDE
IAR C-SPY® Debugger for ARM	C-SPY, the debugger
IAR C-SPY® Simulator	the simulator
IAR C/C++ Compiler™ for ARM	the compiler
IAR Assembler [™] for ARM	the assembler
IAR ILINK Linker™	ILINK, the linker
IAR DLIB Library™	the DLIB library

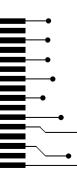
Table 2: Naming conventions used in this guide

Document conventions

Part I. Using the build tools

This part of the IAR C/C++ Development Guide for ARM includes these chapters:

- Introduction to the IAR build tools
- Developing embedded applications
- Data storage
- Functions
- Linking using ILINK
- Linking your application
- The DLIB runtime environment
- Assembler language interface
- Using C
- Using C++
- Application-related considerations
- Efficient coding for embedded applications.



Introduction to the IAR build tools

This chapter gives an introduction to the IAR build tools for the ARM core, which means you will get an overview of:

- The IAR build tools-the build interfaces, compiler, assembler, and linker
- The programming languages
- The available device support
- The extensions provided by the IAR C/C++ Compiler for ARM to support specific features of the ARM core.

The IAR build tools—an overview

In the IAR product installation you can find a set of tools, code examples, and user documentation, all suitable for developing software for ARM-based embedded applications. The tools allow you to develop your application in C, C++, or in assembler language.



IAR Embedded Workbench® is a very powerful Integrated Development Environment (IDE) that allows you to develop and manage complete embedded application projects. It provides an easy-to-learn and highly efficient development environment with maximum code inheritance capabilities, comprehensive and specific target support. IAR Embedded Workbench promotes a useful working methodology, and thus a significant reduction of the development time.

For information about the IDE, see the *IDE Project Management and Building Guide* for ARM.

The compiler, assembler, and linker can also be run from a command line environment, if you want to use them as external tools in an already established project environment.

IAR C/C++ COMPILER

The IAR C/C++ Compiler for ARM is a state-of-the-art compiler that offers the standard features of the C and C++ languages, plus extensions designed to take advantage of the ARM-specific facilities.

IAR ASSEMBLER

The IAR Assembler for ARM is a powerful relocating macro assembler with a versatile set of directives and expression operators. The assembler features a built-in C language preprocessor and supports conditional assembly.

The IAR Assembler for ARM uses the same mnemonics and operand syntax as the Advanced RISC Machines Ltd ARM Assembler, which simplifies the migration of existing code. For more information, see the *IAR Assembler Reference Guide for ARM*.

THE IAR ILINK LINKER

The IAR ILINK Linker for ARM is a powerful, flexible software tool for use in the development of embedded controller applications. It is equally well suited for linking small, single-file, absolute assembler programs as it is for linking large, relocatable input, multi-module, C/C++, or mixed C/C++ and assembler programs.

SPECIFIC ELF TOOLS

ILINK both uses and produces industry-standard ELF and DWARF as object format, additional IAR utilities that handle these formats are provided:

- The IAR Archive Tool—iarchive—creates and manipulates a library (archive) of several ELF object files
- The IAR ELF Tool—ielftool—performs various transformations on an ELF executable image (such as, fill, checksum, format conversion etc)
- The IAR ELF Dumper for ARM—ielfdumparm—creates a text representation of the contents of an ELF relocatable or executable image
- The IAR ELF Object Tool—iobjmanip—is used for performing low-level manipulation of ELF object files
- The IAR Absolute Symbol Exporter—isymexport—exports absolute symbols from a ROM image file, so that they can be used when linking an add-on application.

Note: These ELF utilities are well-suited for object files produced by the tools from IAR Systems. Thus, we recommend using them instead of the GNU binary utilities.

EXTERNAL TOOLS

For information about how to extend the tool chain in the IDE, see the *IDE Project* Management and Building Guide for ARM.

IAR language overview

There are two high-level programming languages you can use with the IAR C/C++ Compiler for:

- C, the most widely used high-level programming language in the embedded systems industry. You can build freestanding applications that follow these standards:
 - Standard C—also known as C99. Hereafter, this standard is referred to as *Standard C* in this guide.
 - C89—also known as C94, C90, C89, and ANSI C. This standard is required when MISRA C is enabled.
- C++, a modern object-oriented programming language with a full-featured library well suited for modular programming. Any of these standards can be used:
 - Standard C++—can be used with different levels of support for exceptions and runtime type information (RTTI).
 - Embedded C++ (EC++)—a subset of the C++ programming standard, which is intended for embedded systems programming. It is defined by an industry consortium, the Embedded C++ Technical committee. See the chapter *Using* C++.
 - IAR Extended Embedded C++ (EEC++)—EC++ with additional features such as full template support, multiple inheritance, namespace support, the new cast operators, as well as the Standard Template Library (STL).

Each of the supported languages can be used in *strict* or *relaxed* mode, or relaxed with IAR extensions enabled. The strict mode adheres to the standard, whereas the relaxed mode allows some common deviations from the standard.

For more information about C, see the chapter Using C.

For more information about C++, Embedded C++, and Extended Embedded C++, see the chapter Using C++.

For information about how the compiler handles the implementation-defined areas of the languages, see the chapter *Implementation-defined behavior for Standard C*.

It is also possible to implement parts of the application, or the whole application, in assembler language. See the *IAR Assembler Reference Guide for ARM*.

Device support

To get a smooth start with your product development, the IAR product installation comes with a wide range of device-specific support.

SUPPORTED ARM DEVICES

The IAR C/C++ Compiler for ARM supports several different ARM cores and devices based on the instruction sets version 4, 5, 6, 6M, and 7. The object code that the compiler generates is not always binary compatible between the cores. Therefore it is crucial to specify a processor option to the compiler. The default core is ARM7TDMI.

PRECONFIGURED SUPPORT FILES

The IAR product installation contains preconfigured files for supporting different devices. If you need additional files for device support, they can be created using one of the provided ones as a template.

Header files for I/O

Standard peripheral units are defined in device-specific I/O header files with the filename extension h. The product package supplies I/O files for all devices that are available at the time of the product release. You can find these files in the arm\inc\<vendor> directory. Make sure to include the appropriate include file in your application source files. If you need additional I/O header files, they can be created using one of the provided ones as a template. For detailed information about the header file format, see EWARM_HeaderFormat.pdf located in the arm\doc\ directory.

Device description files

The debugger handles several of the device-specific requirements, such as definitions of peripheral registers and groups of these, by using device description files. These files are located in the arm\inc directory and they have the filename extension ddf. For more information about these files, see the *C-SPY® Debugging Guide for ARM* and EWARM_DDFFormat.pdf located in the arm\doc\ directory.

EXAMPLES FOR GETTING STARTED

The arm\examples directory contains several hundreds of examples of working applications to give you a smooth start with your development. The complexity of the examples ranges from simple LED blink to USB mass storage controllers. Examples are provided for most of the supported devices.

Special support for embedded systems

This section briefly describes the extensions provided by the compiler to support specific features of the ARM core.

EXTENDED KEYWORDS

The compiler provides a set of keywords that can be used for configuring how the code is generated. For example, there are keywords for declaring special function types.

By default, language extensions are enabled in the IDE.

The command line option -e makes the extended keywords available, and reserves them so that they cannot be used as variable names. See, -e, page 242 for additional information.

For more information about the extended keywords, see the chapter Extended keywords.

PRAGMA DIRECTIVES

The pragma directives control the behavior of the compiler, for example how it allocates memory, whether it allows extended keywords, and whether it issues warning messages.

The pragma directives are always enabled in the compiler. They are consistent with standard C, and are very useful when you want to make sure that the source code is portable.

For more information about the pragma directives, see the chapter *Pragma directives*.

PREDEFINED SYMBOLS

With the predefined preprocessor symbols, you can inspect your compile-time environment, for example the CPU mode and time of compilation.

For more information about the predefined symbols, see the chapter The preprocessor.

SPECIAL FUNCTION TYPES

The special hardware features of the ARM core are supported by the compiler's special function types: software interrupts, interrupts, and fast interrupts. You can write a complete application without having to write any of these functions in assembler language.

For more information, see *Primitives for interrupts, concurrency, and OS-related programming*, page 59.

ACCESSING LOW-LEVEL FEATURES

For hardware-related parts of your application, accessing low-level features is essential. The compiler supports several ways of doing this: intrinsic functions, mixing C and assembler modules, and inline assembler. For information about the different methods, see *Mixing C and assembler*, page 137.

Special support for embedded systems

Developing embedded applications

This chapter provides the information you need to get started developing your embedded software for the ARM core using the IAR build tools.

First, you will get an overview of the tasks related to embedded software development, followed by an overview of the build process, including the steps involved for compiling and linking an application.

Next, the program flow of an executing application is described.

Finally, you will get an overview of the basic settings needed for a project.

Developing embedded software using IAR build tools

Typically, embedded software written for a dedicated microcontroller is designed as an endless loop waiting for some external events to happen. The software is located in ROM and executes on reset. You must consider several hardware and software factors when you write this kind of software.

MAPPING OF INTERNAL AND EXTERNAL MEMORY

Embedded systems typically contain various types of memory, such as on-chip RAM, external DRAM or SRAM, ROM, EEPROM, or flash memory.

As an embedded software developer, you must understand the features of the different types of memory. For example, on-chip RAM is often faster than other types of memories, and variables that are accessed often would in time-critical applications benefit from being placed here. Conversely, some configuration data might be accessed seldom but must maintain their value after power off, so they should be saved in EEPROM or flash memory.

For efficient memory usage, the compiler provides several mechanisms for controlling placement of functions and data objects in memory. For more information, see *Controlling data and function placement in memory*, page 201. The linker places sections of code and data in memory according to the directives you specify in the linker configuration file, see *Placing code and data—the linker configuration file*, page 73.

COMMUNICATION WITH PERIPHERAL UNITS

If external devices are connected to the microcontroller, you might need to initialize and control the signalling interface, for example by using chip select pins, and detect and handle external interrupt signals. Typically, this must be initialized and controlled at runtime. The normal way to do this is to use special function registers (SFR). These are typically available at dedicated addresses, containing bits that control the chip configuration.

Standard peripheral units are defined in device-specific I/O header files with the filename extension h. See *Device support*, page 33. For an example, see *Accessing special function registers*, page 214.

EVENT HANDLING

In embedded systems, using *interrupts* is a method for handling external events immediately; for example, detecting that a button was pressed. In general, when an interrupt occurs in the code, the core immediately stops executing the code it runs, and starts executing an interrupt routine instead.

The compiler supports the following processor exception types: interrupts, software interrupts, and fast interrupts, which means that you can write your interrupt routines in C, see *Interrupt functions*, page 59.

SYSTEM STARTUP

In all embedded systems, system startup code is executed to initialize the system—both the hardware and the software system—before the main function of the application is called. The CPU imposes this by starting execution from a fixed memory address.

As an embedded software developer, you must ensure that the startup code is located at the dedicated memory addresses, or can be accessed using a pointer from the vector table. This means that startup code and the initial vector table must be placed in non-volatile memory, such as ROM, EPROM, or flash.

A C/C++ application further needs to initialize all global variables. This initialization is handled by the linker and the system startup code in conjunction. For more information, see *Application execution—an overview*, page 42.

REAL-TIME OPERATING SYSTEMS

In many cases, the embedded application is the only software running in the system. However, using an RTOS has some advantages. For example, the timing of high-priority tasks is not affected by other parts of the program which are executed in lower priority tasks. This typically makes a program more deterministic and can reduce power consumption by using the CPU efficiently and putting the CPU in a lower-power state when idle.

Using an RTOS can make your program easier to read and maintain, and in many cases smaller as well. Application code can be cleanly separated in tasks which are truly independent of each other. This makes teamwork easier, as the development work can be easily split into separate tasks which are handled by one developer or a group of developers.

Finally, using an RTOS reduces the hardware dependence and creates a clean interface to the application, making it easier to port the program to different target hardware.

INTEROPERABILITY WITH OTHER BUILD TOOLS

The IAR compiler and linker provide support for AEABI, the Embedded Application Binary Interface for ARM. For more information about this interface specification, see the www.arm.com web site.

The advantage of this interface is the interoperability between vendors supporting it; an application can be built up of libraries of object files produced by different vendors and linked with a linker from any vendor, as long as they adhere to the AEABI standard.

AEABI specifies full compatibility for C and C++ object code, and for the C library. The AEABI does not include specifications for the C++ library.

For more information about the AEABI support in the IAR build tools, see *AEABI* compliance, page 192.

The ARM IAR build tools version 6.xx are not fully compatible with earlier versions of the product, see the *IAR Embedded Workbench*® *Migration Guide for ARM*® for more information.

For more information, see Linker optimizations, page 192.

The build process—an overview

This section gives an overview of the build process; how the various build tools—compiler, assembler, and linker—fit together, going from source code to an executable image.

To get familiar with the process in practice, you should run one or more of the tutorials available from the IAR Information Center.

THE TRANSLATION PROCESS

There are two tools in the IDE that translate application source files to intermediary object files. The IAR C/C++ Compiler and the IAR Assembler. Both produce relocatable object files in the industry-standard format ELF, including the DWARF format for debug information.

Note: The compiler can also be used for translating C/C++ source code into assembler source code. If required, you can modify the assembler source code which then can be assembled into object code. For more information about the IAR Assembler, see the *IAR Assembler Reference Guide for ARM*.

This illustration shows the translation process:

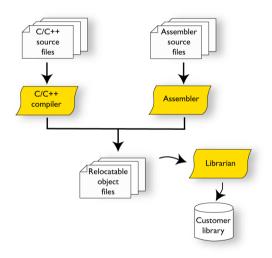


Figure 1: The build process before linking

After the translation, you can choose to pack any number of modules into an archive, or in other words, a library. The important reason you should use libraries is that each module in a library is conditionally linked in the application, or in other words, is only included in the application if the module is used directly or indirectly by a module supplied as an object file. Optionally, you can create a library; then use the IAR utility iarchive.

THE LINKING PROCESS

The relocatable modules, in object files and libraries, produced by the IAR compiler and assembler cannot be executed as is. To become an executable application, they must be *linked*.

Note: Modules produced by a toolset from another vendor can be included in the build as well. Be aware that this also might require a compiler utility library from the same vendor.

The IAR ILINK Linker (ilinkarm.exe) is used for building the final application. Normally, ILINK requires the following information as input:

- · Several object files and possibly certain libraries
- A program start label (set by default)
- The linker configuration file that describes placement of code and data in the memory of the target system.

This illustration shows the linking process:

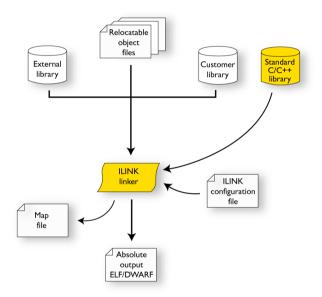


Figure 2: The linking process

Note: The Standard C/C++ library contains support routines for the compiler, and the implementation of the C/C++ standard library functions.

During the linking, ILINK might produce error messages and logging messages on stdout and stderr. The log messages are useful for understanding why an application was linked the way it was, for example, why a module was included or a section removed.

For more information about the procedure performed by ILINK, see *The linking process*, page 71.

AFTER LINKING

The IAR ILINK Linker produces an absolute object file in ELF format that contains the executable image. After linking, the produced absolute executable image can be used for:

- Loading into the IAR C-SPY Debugger or any other compatible external debugger that reads ELF and DWARF.
- Programming to a flash/PROM using a flash/PROM programmer. Before this is possible, the actual bytes in the image must be converted into the standard Motorola 32-bit S-record format or the Intel Hex-32 format. For this, use ielftool, see *The IAR ELF Tool—ielftool*, page 440.

This illustration shows the possible uses of the absolute output ELF/DWARF file:

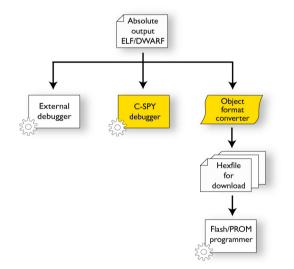


Figure 3: Possible uses of the absolute output ELF/DWARF file

Application execution—an overview

This section gives an overview of the execution of an embedded application divided into three phases, the:

Initialization phase

- Execution phase
- Termination phase.

THE INITIALIZATION PHASE

Initialization is executed when an application is started (the CPU is reset) but before the main function is entered. The initialization phase can for simplicity be divided into:

• Hardware initialization, which generally at least initializes the stack pointer.

The hardware initialization is typically performed in the system startup code <code>cstartup.s</code> and if required, by an extra low-level routine that you provide. It might include resetting/starting the rest of the hardware, setting up the CPU, etc, in preparation for the software C/C++ system initialization.

• Software C/C++ system initialization

Typically, this includes assuring that every global (statically linked) C/C++ symbol receives its proper initialization value before the main function is called.

• Application initialization

This depends entirely on your application. It can include setting up an RTOS kernel and starting initial tasks for an RTOS-driven application. For a bare-bone application, it can include setting up various interrupts, initializing communication, initializing devices, etc.

For a ROM/flash-based system, constants and functions are already placed in ROM. All symbols placed in RAM must be initialized before the main function is called. The linker has already divided the available RAM into different areas for variables, stack, heap, etc.

The following sequence of illustrations gives a simplified overview of the different stages of the initialization.

I When an application is started, the system startup code first performs hardware initialization, such as initialization of the stack pointer to point at the end of the predefined stack area:

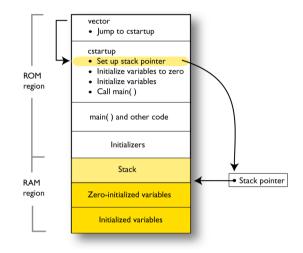
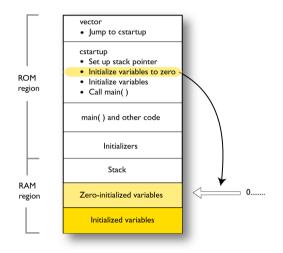


Figure 4: Initializing hardware

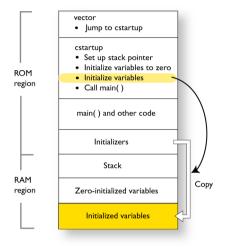
.....



2 Then, memories that should be zero-initialized are cleared, in other words, filled with zeros:

Figure 5: Zero-initializing variables

Typically, this is data referred to as *zero-initialized data*; variables declared as, for example, int i = 0;



3 For *initialized data*, data declared, for example, like int i = 6; the initializers are copied from ROM to RAM:

Figure 6: Initializing variables

4 Finally, the main function is called:

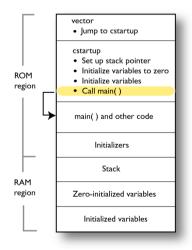


Figure 7: Calling main

For more information about each stage, see *System startup and termination*, page 113. For more information about initialization of data, see *Initialization at system startup*, page 76.

THE EXECUTION PHASE

The software of an embedded application is typically implemented as a loop which is either interrupt-driven or uses polling for controlling external interaction or internal events. For an interrupt-driven system, the interrupts are typically initialized at the beginning of the main function.

In a system with real-time behavior and where responsiveness is critical, a multi-task system might be required. This means that your application software should be complemented with a real-time operating system. In this case, the RTOS and the different tasks must also be initialized at the beginning of the main function.

THE TERMINATION PHASE

Typically, the execution of an embedded application should never end. If it does, you must define a proper end behavior.

To terminate an application in a controlled way, either call one of the Standard C library functions exit, _Exit, or abort, or return from main. If you return from main, the exit function is executed, which means that C++ destructors for static and global variables are called (C++ only) and all open files are closed.

Of course, in case of incorrect program logic, the application might terminate in an uncontrolled and abnormal way—a system crash.

For more information about this, see System termination, page 116.

Building applications—an overview

In the command line interface, this line compiles the source file myfile.c into the object file myfile.o using the default settings:

iccarm myfile.c

On the command line, this line can be used for starting ILINK:

ilinkarm myfile.o myfile2.o -o a.out --config my_configfile.icf

In this example, myfile.o and myfile2.o are object files, and my_configfile.icf is the linker configuration file. The option -o specifies the name of the output file.

Note: By default, the label where the application starts is <u>__iar_program_start</u>. You can use the --entry command line option to change this.

Basic project configuration

This section gives an overview of the basic settings for the project setup that are needed to make the compiler and linker generate the best code for the ARM device you are using. You can specify the options either from the command line interface or in the IDE.

You need to make settings for:

- Processor configuration, that is processor variant, CPU mode, interworking, VFP and floating-point arithmetic, and byte order
- Optimization settings
- Runtime environment
- Customizing the ILINK configuration, see the chapter *Linking your application*.

In addition to these settings, many other options and settings can fine-tune the result even further. For information about how to set options and for a list of all available options, see the chapters *Compiler options*, *Linker options*, and the *IDE Project Management and Building Guide for ARM*, respectively.

PROCESSOR CONFIGURATION

To make the compiler generate optimum code, you should configure it for the ARM core you are using.

Processor variant

The IAR C/C++ Compiler for ARM supports several different ARM cores and devices based on the instruction sets version 4, 5, 6, and 7. All supported cores support Thumb instructions and 64-bit multiply instructions. The object code that the compiler generates is not always binary compatible between the cores. Therefore it is crucial to specify a processor option to the compiler. The default core is ARM7TDMI.



See the *IDE Project Management and Building Guide for ARM*, for information about
setting the **Processor variant** option in the IDE.

Use the --cpu option to specify the ARM core; see --*arm*, page 234 and --*thumb*, page 265, for syntax information.

Interworking

When code is compiled with the --interwork option, ARM and Thumb code can be freely mixed. Interworking functions can be called from both ARM and Thumb code. Interworking is default for devices based on the instruction sets version 5, 6, and 7, or when using the --aeabi compiler option. See also *Veneers*, page 95.



See the *IDE Project Management and Building Guide for ARM* for information about setting the **Generate interwork code** option in the IDE.

Use the --interwork option to specify interworking capabilities for your project; see *--interwork*, page 248, for syntax information.

VFP and floating-point arithmetic

If you are using an ARM core that contains a Vector Floating Point (VFP) coprocessor, you can use the --fpu option to generate code that carries out floating-point operations utilizing the coprocessor, instead of using the software floating-point library routines.



See the *IDE Project Management and Building Guide for ARM*, for information about setting the **FPU** option in the IDE.



Use the --fpu option to specify the ARM core; see --*fpu*, page 246 for syntax information.

Byte order

The compiler supports the big-endian and little-endian byte order. All user and library modules in your application must use the same byte order.



See the *IDE Project Management and Building Guide for ARM* for information about setting the **Endian mode** option in the IDE.



Use the --endian option to specify the byte order for your project; see --endian, page 244, for syntax information.

OPTIMIZATION FOR SPEED AND SIZE

The compiler's optimizer performs, among other things, dead-code elimination, constant propagation, inlining, common sub-expression elimination, static clustering, instruction scheduling, and precision reduction. It also performs loop optimizations, such as unrolling and induction variable elimination.

You can decide between several optimization levels and for the highest level you can choose between different optimization goals—*size, speed*, or *balanced*. Most optimizations will make the application both smaller and faster. However, when this is not the case, the compiler uses the selected optimization goal to decide how to perform the optimization.

The optimization level and goal can be specified for the entire application, for individual files, and for individual functions. In addition, some individual optimizations, such as function inlining, can be disabled.

For information about compiler optimizations and for more information about efficient coding techniques, see the chapter *Efficient coding for embedded applications*.

RUNTIME ENVIRONMENT

To create the required runtime environment you should choose a runtime library and set library options. You might also need to override certain library modules with your own customized versions.

Two different sets of runtime libraries are provided:

The runtime library provided is the IAR DLIB Library, which supports Standard C and C++. This library also supports floating-point numbers in IEEE 754 format and it can be configured to include different levels of support for locale, file descriptors, multibyte characters, etc.

The runtime library contains the functions defined by the C and the C++ standards, and include files that define the library interface (the system header files).

The runtime library you choose can be one of the prebuilt libraries, or a library that you customized and built yourself. The IDE provides a library project template for both libraries, that you can use for building your own library version. This gives you full control of the runtime environment. If your project only contains assembler source code, you do not need to choose a runtime library.

For more information about the runtime environments, see the chapters *The DLIB runtime environment*, respectively.



Setting up for the runtime environment in the IDE

The library is automatically chosen according to the settings you make in **Project>Options>General Options**, on the pages **Target**, **Library Configuration**, **Library Options**. A correct include path is automatically set up for the system header files and for the device-specific include files.

Note that for the DLIB library there are different configurations—Normal, and Full—which include different levels of support for locale, file descriptors, multibyte characters, etc. See *Library configurations*, page 118, for more information.



Setting up for the runtime environment from the command line

You do not have to specify a library file explicitly, as ILINK automatically uses the correct library file.

A library configuration file that matches the library object file is automatically used. To explicitly specify a library configuration, use the --dlib_config option.

In addition to these options you might want to specify any application-specific linker options or the include path to application-specific header files by using the -I option, for example:

-I MyApplication\inc

.

For information about the prebuilt library object files, see *Using prebuilt libraries*, page 101 (DLIB).

Setting library and runtime environment options

You can set certain options to reduce the library and runtime environment size:

- The formatters used by the functions printf, scanf, and their variants, see *Choosing formatters for printf and scanf*, page 105 (DLIB).
- The size of the stack and the heap, see *Setting up the stack*, page 89, and *Setting up the heap*, page 90, respectively.

Basic project configuration

Data storage

This chapter gives a brief introduction to the memory layout of the ARM core and the fundamental ways data can be stored in memory: on the stack, in static (global) memory, or in heap memory. Finally, detailed information about data storage on the stack and the heap is provided.

Introduction

An ARM core can address 4 Gbytes of continuous memory, ranging from 0x00000000 to 0xFFFFFFFF. Different types of physical memory can be placed in the memory range. A typical application will have both read-only memory (ROM) and read/write memory (RAM). In addition, some parts of the memory range contain processor control registers and peripheral units.

DIFFERENT WAYS TO STORE DATA

In a typical application, data can be stored in memory in three different ways:

- Auto variables
- All variables that are local to a function, except those declared static, are stored either in registers or on the stack. These variables can be used as long as the function executes. When the function returns to its caller, the memory space is no longer valid. For more information, see *Auto variables—on the stack*, page 54.Global variables, module-static variables, and local variables declared static

In this case, the memory is allocated once and for all. The word static in this context means that the amount of memory allocated for this kind of variables does not change while the application is running. The ARM core has one single address space and the compiler supports full memory addressing.

• Dynamically allocated data.

An application can allocate data on the *heap*, where the data remains valid until it is explicitly released back to the system by the application. This type of memory is useful when the number of objects is not known until the application executes. Note that there are potential risks connected with using dynamically allocated data in systems with a limited amount of memory, or systems that are expected to run for a long time. For more information, see *Dynamic memory on the heap*, page 55.

Auto variables—on the stack

Variables that are defined inside a function—and not declared static—are named auto variables by the C standard. A few of these variables are placed in processor registers; the rest are placed on the stack. From a semantic point of view, this is equivalent. The main differences are that accessing registers is faster, and that less memory is required compared to when variables are located on the stack.

Auto variables can only live as long as the function executes; when the function returns, the memory allocated on the stack is released.

THE STACK

The stack can contain:

- · Local variables and parameters not stored in registers
- Temporary results of expressions
- The return value of a function (unless it is passed in registers)
- Processor state during interrupts
- Processor registers that should be restored before the function returns (callee-save registers).

The stack is a fixed block of memory, divided into two parts. The first part contains allocated memory used by the function that called the current function, and the function that called it, etc. The second part contains free memory that can be allocated. The borderline between the two areas is called the top of stack and is represented by the stack pointer, which is a dedicated processor register. Memory is allocated on the stack by moving the stack pointer.

A function should never refer to the memory in the area of the stack that contains free memory. The reason is that if an interrupt occurs, the called interrupt function can allocate, modify, and—of course—deallocate memory on the stack.

Advantages

The main advantage of the stack is that functions in different parts of the program can use the same memory space to store their data. Unlike a heap, a stack will never become fragmented or suffer from memory leaks.

It is possible for a function to call itself either directly or indirectly—a recursive function—and each invocation can store its own data on the stack.

Potential problems

The way the stack works makes it impossible to store data that is supposed to live after the function returns. The following function demonstrates a common programming mistake. It returns a pointer to the variable x, a variable that ceases to exist when the function returns.

```
int *MyFunction()
{
    int x;
    /* Do something here. */
    return &x; /* Incorrect */
}
```

Another problem is the risk of running out of stack. This will happen when one function calls another, which in turn calls a third, etc., and the sum of the stack usage of each function is larger than the size of the stack. The risk is higher if large data objects are stored on the stack, or when recursive functions are used.

Dynamic memory on the heap

Memory for objects allocated on the heap will live until the objects are explicitly released. This type of memory storage is very useful for applications where the amount of data is not known until runtime.

In C, memory is allocated using the standard library function malloc, or one of the related functions calloc and realloc. The memory is released again using free.

In C++, a special keyword, new, allocates memory and runs constructors. Memory allocated with new must be released using the keyword delete.

For related information, see --basic_heap, page 272.

POTENTIAL PROBLEMS

Applications that are using heap-allocated objects must be designed very carefully, because it is easy to end up in a situation where it is not possible to allocate objects on the heap.

The heap can become exhausted if your application uses too much memory. It can also become full if memory that no longer is in use was not released.

For each allocated memory block, a few bytes of data for administrative purposes is required. For applications that allocate a large number of small blocks, this administrative overhead can be substantial.

There is also the matter of fragmentation; this means a heap where small sections of free memory is separated by memory used by allocated objects. It is not possible to allocate a new object if no piece of free memory is large enough for the object, even though the sum of the sizes of the free memory exceeds the size of the object.

Unfortunately, fragmentation tends to increase as memory is allocated and released. For this reason, applications that are designed to run for a long time should try to avoid using memory allocated on the heap.

Functions

This chapter contains information about functions. It gives a brief overview of function-related extensions—mechanisms for controlling functions—and describes some of these mechanisms in more detail.

Function-related extensions

In addition to supporting Standard C, the compiler provides several extensions for writing functions in C. Using these, you can:

- Generate code for the different CPU modes ARM and Thumb
- Use primitives for interrupts, concurrency, and OS-related programming
- Inline functions
- Facilitate function optimization
- Access hardware features.

The compiler uses compiler options, extended keywords, pragma directives, and intrinsic functions to support this.

For more information about optimizations, see *Efficient coding for embedded applications*, page 197. For information about the available intrinsic functions for accessing hardware operations, see the chapter *Intrinsic functions*.

ARM and Thumb code

The IAR C/C++ Compiler for ARM can generate code for either the 32-bit ARM, or the 16-bit Thumb or Thumb2 instruction set. Use the --cpu_mode option, alternatively the --arm or --thumb options, to specify which instruction set should be used for your project. For individual functions, it is possible to override the project setting by using the extended keywords __arm and __thumb. You can freely mix ARM and thumb code in the same application, as long as the code is interworking.

The size of all code pointers is 4 bytes. There are restrictions to implicit and explicit casts from code pointers to data pointers or integer types or vice versa. For further information about the restrictions, see *Pointer types*, page 305.

In the chapter *Assembler language interface*, the generated code is studied in more detail in the description of calling C functions from assembler language and vice versa.

Execution in RAM

The __ramfunc keyword makes a function execute in RAM. In other words it places the function in a section that has read/write attributes. The function is copied from ROM to RAM at system startup just like any initialized variable, see *System startup and termination*, page 113.

The keyword is specified before the return type:

```
__ramfunc void foo(void);
```

If a function declared ___ramfunc tries to access ROM, the compiler will issue a warning.

If the whole memory area used for code and constants is disabled—for example, when the whole flash memory is being erased—only functions and data stored in RAM may be used. Interrupts must be disabled unless the interrupt vector and the interrupt service routines are also stored in RAM.

String literals and other constants can be avoided by using initialized variables. For example, the following lines:

```
__ramfunc void test()
{
   /* myc: initializer in ROM */
   const int myc[] = { 10, 20 };
   /* string literal in ROM */
   msg("Hello");
}
```

can be rewritten to:

```
_ramfunc void test()
{
   /* myc: initialized by cstartup */
   static int myc[] = { 10, 20 };
   /* hello: initialized by cstartup */
   static char hello[] = "Hello";
   msg(hello);
}
```

For more information, see Initializing code-copying ROM to RAM, page 93.

Primitives for interrupts, concurrency, and OS-related programming

The IAR C/C++ Compiler for ARM provides the following primitives related to writing interrupt functions, concurrent functions, and OS-related functions:

- The extended keywords __irq, __fiq, __swi, and __nested
- The intrinsic functions __enable_interrupt, __disable_interrupt, __get_interrupt_state, and __set_interrupt_state.

Note: ARM Cortex-M has a different interrupt mechanism than other ARM devices, and for these devices a different set of primitives is available. For more information, see *Interrupts for ARM Cortex-M*, page 64.

INTERRUPT FUNCTIONS

In embedded systems, using interrupts is a method for handling external events immediately; for example, detecting that a button was pressed.

Interrupt service routines

In general, when an interrupt occurs in the code, the core immediately stops executing the code it runs, and starts executing an interrupt routine instead. It is important that the environment of the interrupted function is restored after the interrupt is handled (this includes the values of processor registers and the processor status register). This makes it possible to continue the execution of the original code after the code that handled the interrupt was executed.

The compiler supports interrupts, software interrupts, and fast interrupts. For each interrupt type, an interrupt routine can be written.

All interrupt functions must be compiled in ARM mode; if you are using Thumb mode, use the __arm extended keyword or the #pragma type_attribute=__arm directive to override the default behavior. This is not applicable for Cortex-M devices.

Interrupt vectors and the interrupt vector table

Each interrupt routine is associated with a vector address/instruction in the exception vector table, which is specified in the ARM cores documentation. The interrupt vector is the address in the exception vector table. For the ARM cores, the exception vector table starts at address 0×0 .

Defining an interrupt function-an example

To define an interrupt function, the __irq or the __fiq keyword can be used. For example:

```
__irq __arm void IRQ_Handler(void)
{
   /* Do something */
}
```

See the ARM cores documentation for more information about the interrupt vector table.

Note: An interrupt function must have the return type void, and it cannot specify any parameters.

INSTALLING EXCEPTION FUNCTIONS

All interrupt functions and software interrupt handlers must be installed in the vector table. This is done in assembler language in the system startup file cstartup.s.

The default implementation of the ARM exception vector table in the standard runtime library jumps to predefined functions that implement an infinite loop. Any exception that occurs for an event not handled by your application will therefore be caught in the infinite loop (B_{\cdot}) .

The predefined functions are defined as weak symbols. A weak symbol is only included by the linker as long as no duplicate symbol is found. If another symbol is defined with the same name, it will take precedence. Your application can therefore simply define its own exception function by just defining it using the correct name.

These exception function names are defined in cstartup.s and referred to by the library exception vector code:

Undefined_Handler SWI_Handler Prefetch_Handler Abort_Handler IRQ_Handler FIQ_Handler

To implement your own exception handler, define a function using the appropriate exception function name from the list above.

For example to add an interrupt function in C, it is sufficient to define an interrupt function named IRQ_Handler:

__irq __arm void IRQ_Handler()
{
}

An interrupt function must have C linkage, read more in Calling convention, page 150.

If you use C++, an interrupt function could look, for example, like this:

```
extern "C"
{
    __irq __arm void IRQ_Handler(void);
}
__irq __arm void IRQ_Handler(void)
{
}
```

No other changes are needed.

INTERRUPTS AND FAST INTERRUPTS

The interrupt and fast interrupt functions are easy to handle as they do not accept parameters or have a return value.

- To declare an interrupt function, use the __irq extended keyword or the #pragma type_attribute=__irq directive. For syntax information, see__irq, page 318 and type_attribute, page 342, respectively.
- To declare a fast interrupt function, use the __fiq extended keyword or the #pragma type_attribute=__fiq directive. For syntax information, see __fiq, page 318, and type_attribute, page 342, respectively.

Note: An interrupt function (irq) and a fast interrupt function (fiq) must have a return type of void and cannot have any parameters. A software interrupt function (swi) may have parameters and return values. By default, only four registers, R0-R3, can be used for parameters and only the registers R0-R1 can be used for return values.

NESTED INTERRUPTS

Interrupts are automatically disabled by the ARM core prior to entering an interrupt handler. If an interrupt handler re-enables interrupts, calls functions, and another interrupt occurs, then the return address of the interrupted function—stored in LR—is overwritten when the second IRQ is taken. In addition, the contents of SPSR will be destroyed when the second interrupt occurs. The __irq keyword itself does not save and restore LR and SPSR. To make an interrupt handler perform the necessary steps needed when handling nested interrupts, the keyword __nested must be used in addition to __irq. The function prolog—function entrance sequence—that the compiler generates for nested interrupt handlers will switch from IRQ mode to system mode. Make sure that both the IRQ stack and system stack is set up. If you use the default cstartup.s file, both stacks are correctly set up.

Compiler-generated interrupt handlers that allow nested interrupts are supported for IRQ interrupts only. The FIQ interrupts are designed to be serviced quickly, which in most cases mean that the overhead of nested interrupts would be too high.

This example shows how to use nested interrupts with the ARM vectored interrupt controller (VIC):

```
__irq __nested __arm void interrupt_handler(void)
{
    void (*interrupt_task)();
    unsigned int vector;
    /* Get interrupt vector. */
    vector = VICVectAddr;
    interrupt_task = (void(*)()) vector;
    /* Allow other IRQ interrupts to be serviced. */
    __enable_interrupt();
    /* Execute the task associated with this interrupt. */
    (*interrupt_task)();
}
```

Note: The __nested keyword requires the processor mode to be in either User or System mode.

SOFTWARE INTERRUPTS

Software interrupt functions are slightly more complex than other interrupt functions, in the way that they need a software interrupt handler (a dispatcher), are invoked (called) from running application software, and that they accept arguments and have return values. The mechanisms for calling a software interrupt function and how the software interrupt handler dispatches the call to the actual software interrupt function is described here.

Calling a software interrupt function

To call a software interrupt function from your application source code, the assembler instruction SVC #immed is used, where immed is an integer value that is referred to as the software interrupt number—or swi_number—in this guide. The compiler provides an easy way to implicitly generate this instruction from C/C++ source code, by using the __swi keyword and the #pragma swi_number directive when declaring the function.

A __swi function can for example be declared like this:

#pragma swi_number=0x23
__swi int swi_function(int a, int b);

In this case, the assembler instruction SVC 0x23 will be generated where the function is called.

Software interrupt functions follow the same calling convention regarding parameters and return values as an ordinary function, except for the stack usage, see *Calling convention*, page 150.

For more information, see __swi, page 322, and swi_number, page 341, respectively.

The software interrupt handler and functions

The interrupt handler, for example SWI_Handler works as a dispatcher for software interrupt functions. It is invoked from the interrupt vector and is responsible for retrieving the software interrupt number and then calling the proper software interrupt function. The SWI_Handler must be written in assembler as there is no way to retrieve the software interrupt number from C/C++ source code.

The software interrupt functions

The software interrupt functions can be written in C or C++. Use the __swi keyword in a function definition to make the compiler generate a return sequence suited for a specific software interrupt function. The #pragma swi_number directive is not needed in the interrupt function definition.

For more information, see __swi, page 322.

Setting up the software interrupt stack pointer

If software interrupts will be used in your application, then the software interrupt stack pointer (SVC_STACK) must be set up and some space must be allocated for the stack. The SVC_STACK pointer can be set up together with the other stacks in the cstartup.s file. As an example, see the set up of the interrupt stack pointer. Relevant space for the SVC_STACK pointer is set up in the linker configuration file, see *Setting up the stack*, page 89.

INTERRUPT OPERATIONS

An interrupt function is called when an external event occurs. Normally it is called immediately while another function is executing. When the interrupt function has finished executing, it returns to the original function. It is imperative that the environment of the interrupted function is restored; this includes the value of processor registers and the processor status register. When an interrupt occurs, the following actions are performed:

- The operating mode is changed corresponding to the particular exception
- The address of the instruction following the exception entry instruction is saved in R14 of the new mode
- The old value of the CPSR is saved in the SPSR of the new mode
- Interrupt requests are disabled by setting bit 7 of the CPSR and, if the exception is a
 fast interrupt, further fast interrupts are disabled by setting bit 6 of the CPSR
- The PC is forced to begin executing at the relevant vector address.

For example, if an interrupt for vector 0×18 occurs, the processor will start to execute code at address 0×18 . The memory area that is used as start location for interrupts is called the interrupt vector table. The content of the interrupt vector is normally a branch instruction jumping to the interrupt routine.

Note: If the interrupt function enables interrupts, the special processor registers needed to return from the interrupt routine must be assumed to be destroyed. For this reason they must be stored by the interrupt routine to be restored before it returns. This is handled automatically if the __nested keyword is used.

INTERRUPTS FOR ARM CORTEX-M

ARM Cortex-M has a different interrupt mechanism than previous ARM architectures, which means the primitives provided by the compiler are also different.

On ARM Cortex-M, an interrupt service routine enters and returns in the same way as a normal function, which means no special keywords are required. Thus, the keywords __irq, __fiq, and __nested are not available when you compile for ARM Cortex-M.

These exception function names are defined in cstartup_M.c and cstartup_M.s. They are referred to by the library exception vector code:

NMI_Handler HardFault_Handler MemManage_Handler BusFault_Handler UsageFault_Handler SVC_Handler DebugMon_Handler PendSV_Handler SysTick_Handler

The vector table is implemented as an array. It should always have the name __vector_table, because cmain refers to that symbol and C-SPY looks for that symbol when determining where the vector table is located.

The predefined exception functions are defined as weak symbols. A weak symbol is only included by the linker as long as no duplicate symbol is found. If another symbol is defined with the same name, it will take precedence. Your application can therefore simply define its own exception function by just defining it using the correct name from the list above. If you need other interrupts or other exception handlers, you must make a copy of the cstartup_M.c or cstartup_M.s file and make the proper addition to the vector table.

The intrinsic functions __get_CPSR and __set_CPSR are not available when you compile for ARM Cortex-M. Instead, if you need to get or set values of these or other registers, you can use inline assembler. For more information, see *Passing values between C and assembler objects*, page 215.

C++ AND SPECIAL FUNCTION TYPES

C++ member functions can be declared using special function types, with the exception that interrupt member functions must be static. When a non-static member function is called, it must be applied to an object. When an interrupt occurs and the interrupt function is called, there is no object available to apply the member function to.

Special function types can be used for static member functions. For example, in the following example, the function handler is declared as an interrupt function:

```
class Device
{
   static __arm __irq void handler();
};
```

Inlining functions

Function inlining means that a function, whose definition is known at compile time, is integrated into the body of its caller to eliminate the overhead of the function call. This optimization, which is performed at optimization level High, normally reduces execution time, but might increase the code size. The resulting code might become more difficult to debug. Whether the inlining actually occurs is subject to the compiler's heuristics.

The compiler heuristically decides which functions to inline. Different heuristics are used when optimizing for speed, size, or when balancing between size and speed. Normally, code size does not increase when optimizing for size.

C VERSUS C++ SEMANTICS

In C++, all definitions of a specific inline function in separate translation units must be exactly the same. If the function is not inlined in one or more of the translation units, then one of the definitions from these translation units will be used as the function implementation.

In C, you must manually select one translation unit that includes the non-inlined version of an inline function. You do this by explicitly declaring the function as extern in that translation unit. If you declare the function as extern in more than one translation unit, the linker will issue a *multiple definition* error. In addition, in C, inline functions cannot refer to static variables or functions.

For example:

```
// In a header file.
static int sX;
inline void F(void)
{
    //static int sY; // Cannot refer to statics.
    //sX; // Cannot refer to statics.
}
// In one source file.
// Declare this F as the non-inlined version to use.
extern inline void F();
```

FEATURES CONTROLLING FUNCTION INLINING

There are several mechanisms for controlling function inlining:

• The inline keyword advises the compiler that the function defined immediately after the directive should be inlined.

If you compile your function in C or C++ mode, the keyword will be interpreted according to its definition in Standard C or Standard C++, respectively.

The main difference in semantics is that in Standard C you cannot (in general) simply supply an inline definition in a header file. You must supply an external definition in one of the compilation units, by designating the inline definition as being external in that compilation unit.

• #pragma inline is similar to the inline keyword, but with the difference that the compiler always uses C++ inline semantics.

By using the #pragma inline directive you can also disable the compiler's heuristics to either force inlining or completely disable inlining. For more information, see *inline*, page 332.

- --use_c++_inline forces the compiler to use C++ semantics when compiling a Standard C source code file.
- --no_inline, #pragma optimize=no_inline, and #pragma inline=never all disable function inlining. By default, function inlining is enabled at optimization level High.

The compiler can only inline a function if the definition is known. Normally, this is restricted to the current translation unit. However, when the --mfc compiler option for multi-file compilation is used, the compiler can inline definitions from all translation units in the multi-file compilation unit. For more information, see *Multi-file compilation units*, page 206.

For more information about the function inlining optimization, see *Function inlining*, page 209.

Inlining functions

Linking using ILINK

This chapter describes the linking process using the IAR ILINK Linker and the related concepts—first with an overview and then in more detail.

Linking—an overview

The IAR ILINK Linker is a powerful, flexible software tool for use in the development of embedded applications. It is equally well suited for linking small, single-file, absolute assembler programs as it is for linking large, relocatable, multi-module, C/C++, or mixed C/C++ and assembler programs.

ILINK combines one or more relocatable object files—produced by the IAR Systems compiler or assembler—with selected parts of one or more object libraries to produce an executable image in the industry-standard format *Executable and Linking Format* (ELF).

ILINK will automatically load only those library modules—user libraries and Standard C or C++ library variants—that are actually needed by the application you are linking. Further, ILINK eliminates duplicate sections and sections that are not required.

ILINK can link both ARM and Thumb code, as well as a combination of them. By automatically inserting additional instructions (veneers), ILINK will assure that the destination will be reached for any calls and branches, and that the processor state is switched when required. For more details about how to generate veneers, see *Veneers*, page 95.

ILINK uses a *configuration file* where you can specify separate locations for code and data areas of your target system memory map. This file also supports automatic handling of the application's initialization phase, which means initializing global variable areas and code areas by copying initializers and possibly decompressing them as well.

The final output produced by ILINK is an absolute object file containing the executable image in the ELF (including DWARF for debug information) format. The file can be downloaded to C-SPY or any other compatible debugger that supports ELF/DWARF, or it can be stored in EPROM or flash.

To handle ELF files, various tools are included. For information about included utilities, see *Specific ELF tools*, page 32.

Modules and sections

Each relocatable object file contains one module, which consists of:

- Several sections of code or data
- Runtime attributes specifying various types of information, for example the version of the runtime environment
- Optionally, debug information in DWARF format
- A symbol table of all global symbols and all external symbols used.

A *section* is a logical entity containing a piece of data or code that should be placed at a physical location in memory. A section can consist of several *section fragments*, typically one for each variable or function (symbols). A section can be placed either in RAM or in ROM. In a normal embedded application, sections that are placed in RAM do not have any content, they only occupy space.

Each section has a name and a type attribute that determines the content. The type attribute is used (together with the name) for selecting sections for the ILINK configuration. The most commonly used attributes are:

code	Executable code
readonly	Constant variables
readwrite	Initialized variables
zeroinit	Zero-initialized variables

Note: In addition to these section types—sections that contain the code and data that are part of your application—a final object file will contain many other types of sections, for example sections that contain debugging information or other type of meta information.

A section is the smallest linkable unit; but if possible, ILINK can exclude smaller units—section fragments—from the final application. For more information, see *Keeping modules*, page 89, and *Keeping symbols and sections*, page 89.

At compile time, data and functions are placed in different sections. At link time, one of the most important functions of the linker is to assign addresses to the various sections used by the application.

The IAR build tools have many predefined section names. See the chapter *Section reference* for more information about each section.

The linking process

The relocatable modules in object files and libraries, produced by the IAR compiler and assembler, cannot be executed as is. To become an executable application, they must be *linked*.

Note: Modules produced by a toolset from another vendor can be included in the build as well, as long as the module is AEABI (ARM Embedded Application Binary Interface) compliant. Be aware that this also might require a compiler utility library from the same vendor.

The IAR ILINK Linker is used for the link process. It normally performs the following procedure (note that some of the steps can be turned off by command line options or by directives in the linker configuration file):

- Determine which modules to include in the application. Modules provided in object files are always included. A module in a library file is only included if it provides a definition for a global symbol that is referenced from an included module.
- Select which standard library files to use. The selection is based on attributes of the included modules. These libraries are then used for satisfying any still outstanding undefined symbols.
- Determine which sections/section fragments from the included modules to include in the application. Only those sections/section fragments that are actually needed by the application are included. There are several ways to determine of which sections/section fragments that are needed, for example, the __root object attribute, the #pragma required directive, and the keep linker directive. In case of duplicate sections, only one is included.
- Where appropriate, arrange for the initialization of initialized variables and code in RAM. The initialize directive causes the linker to create extra sections to enable copying from ROM to RAM. Each section that will be initialized by copying is divided into two sections, one for the ROM part and one for the RAM part. If manual initialization is not used, the linker also arranges for the startup code to perform the initialization.
- Determine where to place each section according to the section placement directives in the *linker configuration file*. Sections that are to be initialized by copying appear twice in the matching against placement directives, once for the ROM part and once for the RAM part, with different attributes. During the placement, the linker also adds any required veneers to make a code reference reach its destination or to switch CPU modes.
- Produce an absolute file that contains the executable image and any debug information provided. The contents of each needed section in the relocatable input files is calculated using the relocation information supplied in its file and the addresses determined when placing sections. This process can result in one or more

relocation failures if some of the requirements for a particular section are not met, for instance if placement resulted in the destination address for a PC-relative jump instruction being out of range for that instruction.

• Optionally, produce a map file that lists the result of the section placement, the address of each global symbol, and finally, a summary of memory usage for each module and library.

This illustration shows the linking process:

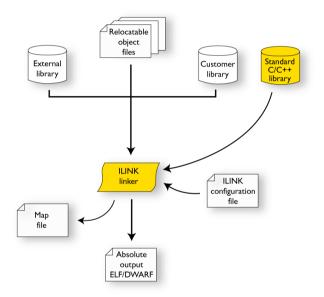


Figure 8: The linking process

During the linking, ILINK might produce error messages and logging messages on stdout and stderr. The log messages are useful for understanding why an application was linked as it was. For example, why a module or section (or section fragment) was included.

Note: To see the actual content of an ELF object file, use ielfdumparm. See *The IAR ELF Dumper for arm—ielfdumparm*, page 442.

Placing code and data-the linker configuration file

The placement of sections in memory is performed by the IAR ILINK Linker. It uses the *linker configuration file* where you can define how ILINK should treat each section and how they should be placed into the available memories.

A typical linker configuration file contains definitions of:

- Available addressable memories
- · Populated regions of those memories
- How to treat input sections
- Created sections
- How to place sections into the available regions.

The file consists of a sequence of declarative directives. This means that the linking process will be governed by all directives at the same time.

To use the same source code with different derivatives, just rebuild the code with the appropriate configuration file.

A SIMPLE EXAMPLE OF A CONFIGURATION FILE

A simple configuration file can look like this:

```
/* The memory space denoting the maximum possible amount
   of addressable memory */
define memory Mem with size = 4G;
/* Memory regions in an address space */
define region ROM = Mem: [from 0x00000 size 0x10000];
define region RAM = Mem: [from 0x20000 size 0x10000];
/* Create a stack */
define block STACK with size = 0x1000, alignment = 8 { };
/* Handle initialization */
do not initialize { section .noinit };
initialize by copy { readwrite }; /* Initialize RW sections,
                                     exclude zero-initialized
                                     sections */
/* Place startup code at a fixed address */
place at start of ROM { readonly section .cstartup };
/* Place code and data */
place in ROM { readonly }; /* Place constants and initializers in
                              ROM: .rodata and .data_init */
```

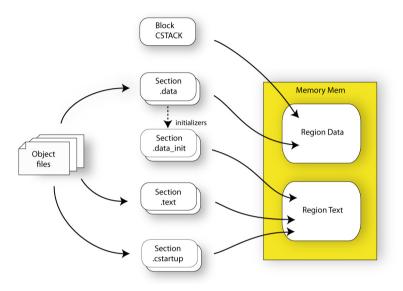
This configuration file defines one addressable memory Mem with the maximum of 4 Gbytes of memory. Further, it defines a ROM region and a RAM region in Mem, namely ROM and RAM. Each region has the size of 64 Kbytes.

The file then creates an empty block called STACK with a size of 4 Kbytes in which the application stack will reside. To create a *block* is the basic method which you can use to get detailed control of placement, size, etc. It can be used for grouping sections, but also as in this example, to specify the size and placement of an area of memory.

Next, the file defines how to handle the initialization of variables, read/write type (readwrite) sections. In this example, the initializers are placed in ROM and copied at startup of the application to the RAM area. By default, ILINK may compress the initializers if this appears to be advantageous.

The last part of the configuration file handles the actual placement of all the sections into the available regions. First, the startup code—defined to reside in the read-only (readonly) section .cstartup—is placed at the start of the ROM region, that is at address 0x10000. Note that the part within {} is referred to as *section selection* and it selects the sections for which the directive should be applied to. Then the rest of the read-only section .cstartup } takes precedence over the more generic section selection { readonly section .cstartup } takes precedence over the more generic section selection { readonly }.

Finally, the read/write (readwrite) sections and the STACK block are placed in the RAM region.



This illustration gives a schematic overview of how the application is placed in memory:

Figure 9: Application in memory

In addition to these standard directives, a configuration file can contain directives that define how to:

- Map a memory that can be addressed in multiple ways
- Handle conditional directives
- Create symbols with values that can be used in the application
- More in detail, select the sections a directive should be applied to
- More in detail, initialize code and data.

For more details and examples about customizing the linker configuration file, see the chapter *Linking your application*.

For more information about the linker configuration file, see the chapter *The linker configuration file*.

Initialization at system startup

In Standard C, all static variables—variables that are allocated at a fixed memory address—must be initialized by the runtime system to a known value at application startup. This value is either an explicit value assigned to the variable, or if no value is given, it is cleared to zero. In the compiler, there is one exception to this rule and that is variables declared __no_init which are not initialized at all.

The compiler generates a specific type of section for each type of variable initialization:

Categories of declared data	Source	Section type	Section name	Section content
Zero-initialized data	int i;	Read/write data, zero-init	.bss	None
Zero-initialized data	int i = 0;	Read/write data, zero-init	.bss	None
Initialized data (non-zero)	int i = 6;	Read/write data	.data	The initializer
Non-initialized data	no_init int i;	Read/write data, zero-init	.noinit	None
Constants	<pre>const int i = 6;</pre>	Read-only data	.rodata	The constant
Code	<pre>ramfunc void myfunc() {}</pre>	Read/write code	.textrw	The code

Table 3: Sections holding initialized data

Note: Clustering of static variables might group zero-initialized variables together with initialized data in .data. The compiler can decide to place constants in the .text section to avoid loading the address of a constant from a constant table.

For information about all supported sections, see the chapter Section reference.

THE INITIALIZATION PROCESS

Initialization of data is handled by ILINK and the system startup code in conjunction.

To configure the initialization of variables, you must consider these issues:

 Sections that should be zero-initialized are handled automatically by ILINK; they should only be placed in RAM • Sections that should be initialized, except for zero-initialized sections, should be listed in an initialize directive

Normally during linking, a section that should be initialized is split into two sections, where the original initialized section will keep the name. The contents are placed in the new initializer section, which will get the original name suffixed with _init. The initializers should be placed in ROM and the initialized sections in RAM, by means of placement directives. The most common example is the .data section which the linker splits into .data and .data_init.

- Sections that contains constants should not be initialized; they should only be placed in flash/ROM
- Sections holding __no_init declared variables should not be initialized and thus should be listed in a do not initialize directive. They should also be placed in RAM.

In the linker configuration file, it can look like this:

Note: When compressed initializers are used (see *Initialize directive*, page 408), the contents sections (that is, sections with the _init suffix) are not listed as separate sections in the map file. Instead, they are combined into aggregates of "initializer bytes". You can place the contents sections the usual way in the linker configuration file; however, this affects the placement (and possibly the number) of the "initializer bytes" aggregates.

For more information about and examples of how to configure the initialization, see *Linking considerations*, page 85.

C++ DYNAMIC INITIALIZATION

The compiler places subroutine pointers for performing C++ dynamic initialization into sections of the ELF section types SHT_PREINIT_ARRAY and SHT_INIT_ARRAY. By default, the linker will place these into a linker-created block, ensuring that all sections of the section type SHT_PREINIT_ARRAY are placed before those of the type SHT_INIT_ARRAY. If any such sections were included, code to call the routines will also be included.

The linker-created blocks are only generated if the linker configuration does not contain section selector patterns for the preinit_array and init_array section types. The effect of the linker-created blocks will be very similar to what happens if the linker configuration file contains this:

If you put this into your linker configuration file, you must also mention the CPP_INIT block in one of the section placement directives. If you wish to select where the linker-created block is placed, you can use a section selector with the name ".init_array".

See also Section-selectors, page 413.

Stack usage analysis

Under the right circumstances, the linker can accurately calculate the maximum stack usage for each call graph root (each function that is not called from another function).

If you enable stack usage analysis, a stack usage chapter will be added to the linker map file, listing for each call graph root the particular call chain which results in the maximum stack depth.

This is only accurate if there is accurate stack usage information for each function in the application.

In general, the compiler will generate this information for each C function, but if there are indirect calls (calls using function pointers) in your application, you must supply a list of possible functions that can be called from each calling function. You can do this by using pragma directives in the source file, or by using a separate stack usage control file when linking.

If you use a stack usage control file, you can also supply stack usage information for functions in modules that do not have stack usage information.

You can use the check that directive in your linker configuration file to check that the stack usage calculated by the linker does not exceed the stack space you have allocated.

LIMITATIONS

Apart from missing or incorrect stack usage information, there are also other sources of inaccuracy in the analysis:

- The linker might not always be able to identify all functions in object modules that lack stack usage information. In particular this might be a problem with object modules written in assembler or produced by non-IAR tools.
- If you use inline assembler to change the frame size or to perform function calls, this will not be reflected in the analysis.
- Extra space consumed by other sources (the processor, an operating system, etc) is not accounted for.
- C++ source code that uses exceptions is not supported.
- If you use other forms of function calls, like software interrupts, they will not be reflected in the call graph.
- Using multi-file compilation (--mfc) can interfere with using a stack usage control file to specify properties of module-local functions in the involved files.

Note that stack usage analysis produces a worst case result. The program might not actually ever end up in the maximum call chain, by design, or by coincidence. In particular, the set of possible destinations for a virtual function call in C++ might sometimes include implementations of the function in question which cannot, in fact be called from that point in the code.



Stack usage analysis is only a complement to actual measurement. If the result is important, you need to perform independent validation of the results of the analysis.

STACK USAGE CONTROL FILES

A stack usage control file contains stack usage control directives.

Using stack usage control files, you can:

- Specify complete stack usage information (call graph root category, stack usage, and possible calls) for a function, by using the stack usage control directive function.
- Exclude certain functions from stack usage analysis, by using the stack usage control directive exclude.
- Specify the possible destinations for indirect calls in a function, by using the stack usage control directive possible calls.

- Specify that functions are call graph roots, including an optional call graph root category, by using the stack usage control directive call graph root.
- Specify a maximum recursion depth for a recursion nest (a set of cycles in the call graph with at least one common node).
- Selectively suppress the warning about unmentioned functions referenced by a module for which you have supplied stack usage information in the stack usage control file.

If your interrupt functions have not already been designated as call graph roots by the compiler, you must do so manually. You can do this either by using the #pragma call_graph_root directive in your source code or by using a simple stack usage control file, which might look something like this:

call graph root [interrupt]: Irq1Handler, Irq2Handler;

For more information, see *call_graph_root*, page 328 and the chapter *Stack usage control files*, page 429.

SOURCE ANNOTATION

As an alternative to specifying possible calls in a stack usage control file, you can instead annotate the source code.

In C files, at the point of an indirect call, you can use the #pragma calls directive to list the possible destinations for that call.

You can also, at the definition of a function, specify that it is a call graph root by using the #pragma call_graph_root directive.

SITUATIONS WHERE WARNINGS ARE ISSUED

When stack usage analysis is enabled in the linker, warnings will be generated in the following circumstances:

- There is at least one function without stack usage information.
- There is at least one indirect call site in the application for which a list of possible called functions has not been supplied.
- There are no known indirect calls, but there is at least one uncalled function that is not known to be a call graph root.
- The application contains recursion (a cycle in the call graph) for which no maximum recursion depth has been supplied, or which is of a form for which the linker is unable to calculate a reliable estimate of stack usage.
- There are calls to a function declared as a call graph root.
- You have used the stack usage control file to supply stack usage information for functions in a module that does not have such information, and there are functions

referenced by that module which have not been mentioned as being called in the stack usage control file.

MAP FILE CONTENTS

When stack usage analysis is enabled, the linker map file contains a stack usage chapter with a summary of the stack usage for each call graph root category, and lists the call chain that results in the maximum stack depth for each call graph root. This is an example of what the stack usage chapter in the map file might look like:

Call Graph Root Category	Max Use	Total U	se		
interrupt Program entry		1:			
Program entry "iar_program_start": 0x Maximum call chain	000085ac	168	bytes		
"iar_program_start" "cmain" "printf" "_PrintfTiny" "_Prout" "putchar" "write" "dwrite" "iar_sh_stdout" "iar_get_ttio"		0 0 8 24 56 16 16 0 0 24 24 24			
interrupt "FaultHandler": 0x0000843	4				
Maximum call chain		32	bytes		
"FaultHandler"		32			

```
interrupt
"IRQHandler": 0x00008424
Maximum call chain 104 bytes
"IRQHandler" 24
"do_something" in suexample.o [1] 80
```

The summary contains the depth of the deepest call chain in each category as well as the sum of the depths of the deepest call chains in that category.

In this case, the maximum stack depth for the program entry (<u>__iar_program_start</u>) is 168 bytes, and occurs inside the system library printf function. Public functions are listed by name, while module-local functions also include the name of the module (like do_something above).

CHECKING THAT THE STACK IS LARGE ENOUGH

You can use the check that directive in your linker configuration file to check that the stack is large enough.

For example:

When linking, the linker emits an error if the expression is false (zero). In this example there would be an error if the sum of 168 (the maximum stack usage of the program entry), 136 (the sum of the maximum stack usages in category "interrupt"), and 100 (a safety margin) is greater than the size of the CSTACK block.

CALL GRAPH LOG

To help you interpret the results of the stack usage analysis, there is a log output option that produces a simple text representation of the call graph (--log call_graph).

Example output:

```
Program entry:
0 ___iar_program_start [168]
  0 _____cmain [168]
    0 iar data init3 [16]
      8 __iar_zero_init3 [8]
       16 - [0]
      8 ___iar_copy_init3 [8]
        16 - [0]
    0 __low_level_init [0]
    0 main [168]
      8 printf [160]
        32 _PrintfTiny [136]
          88 _Prout [80]
            104 putchar [64]
              120 __write [48]
                120 dwrite [48]
                  120 ____iar__sh__stdout [48]
                    144 ____iar_get_ttio [24]
                      168 iar lookup ttioh [0]
                  120 ____iar__sh_write [24]
                    144 - [0]
          88 __aeabi_uidiv [0]
            88 __aeabi_idiv0 [0]
          88 strlen [0]
    0 exit [8]
      0 _exit [8]
        0 __exit [8]
          0 __iar_close_ttio [8]
            8 __iar_lookup_ttioh [0] ***
    0 __exit [8] ***
```

Each line consists of this information:

- The stack usage at the point of call of the function
- The name of the function, or a single '-' to indicate usage in a function at a point with no function call (typically in a leaf function)
- The stack usage along the deepest call chain from that point. If no such value could be calculated, "[---]" is output instead. "***" marks functions that have already been shown.

CALL GRAPH XML OUTPUT

The linker can also produce a call graph file in XML format. This file contains one node for each function in your application, with the stack usage and call information relevant to that function. It is intended to be input for post-processing tools and is not particularly human-readable.

For more information about the XML format used, see the callGraph.txt file in your product installation.

Linking your application

This chapter lists aspects that you must consider when linking your application. This includes using ILINK options and tailoring the linker configuration file.

Finally, this chapter provides some hints for troubleshooting.

Linking considerations

Before you can link your application, you must set up the configuration required by ILINK. Typically, you must consider:

- Defining your own memory areas
- Placing sections
- Keeping modules in the application
- · Keeping symbols and sections in the application
- Application startup
- Setting up the stack and heap
- Setting up the atexit limit
- Changing the default initialization
- Symbols for controlling the application
- Standard library handling
- Other output formats than ELF/DWARF
- Veneers.

CHOOSING A LINKER CONFIGURATION FILE

The config directory contains two ready-made templates for the linker configuration files:

- generic.icf, designed for all cores except for Cortex-M cores
- generic_cortex.icf, designed for all Cortex-M cores.

The files contain the information required by ILINK. The only change, if any, you will normally have to make to the supplied configuration file is to customize the start and end addresses of each region so they fit the target system memory map. If, for example, your application uses additional external RAM, you must also add details about the external RAM memory area.

For some devices, device-specific configuration files are automatically selected.

To edit a linker configuration file, use the editor in the IDE, or any other suitable editor. Alternatively, choose **Project>Options>Linker** and click the **Edit** button on the **Config** page to open the dedicated linker configuration file editor.

Do not change the original template file. We recommend that you make a copy in the working directory, and modify the copy instead. If you are using the linker configuration file editor in the IDE, the IDE will make a copy for you.

Each project in the IDE should have a reference to one, and only one, linker configuration file. This file can be edited, but for the majority of all projects it is sufficient to configure the vital parameters in **Project>Options>Linker>Config**.

DEFINING YOUR OWN MEMORY AREAS

The default configuration file that you selected has predefined ROM and RAM regions. This example will be used as a starting-point for all further examples in this chapter:

```
/* Define the addressable memory */
define memory Mem with size = 4G;
/* Define a region named ROM with start address 0 and to be 64
Kbytes large */
```

define region ROM = Mem: [from 0 size 0x10000];

```
/* Define a region named RAM with start address 0x20000 and to be
64 Kbytes large */
define region RAM = Mem:[from 0x20000 size 0x10000];
```

Each region definition must be tailored for the actual hardware.

To find out how much of each memory that was filled with code and data after linking, inspect the memory summary in the map file (command line option --map).

Adding an additional region

To add an additional region, use the define region directive, for example:

```
/* Define a 2nd ROM region to start at address 0x80000 and to be
128 Kbytes large */
define region ROM2 = Mem:[from 0x80000 size 0x20000];
```

Merging different areas into one region

If the region is comprised of several areas, use a region expression to merge the different areas into one region, for example:

or equivalently

PLACING SECTIONS

The default configuration file that you selected places all predefined sections in memory, but there are situations when you might want to modify this. For example, if you want to place the section that holds constant symbols in the CONSTANT region instead of in the default place. In this case, use the place in directive, for example:

```
/* Place sections with readonly content in the ROM region */
place in ROM {readonly};
```

/* Place the constant symbols in the CONSTANT region */
place in CONSTANT {readonly section .rodata};

Note: Placing a section—used by the IAR build tools—in a different memory which use a different way of referring to its content, will fail.

For the result of each placement directive after linking, inspect the placement summary in the map file (the command line option --map).

Placing a section at a specific address in memory

To place a section at a specific address in memory, use the place at directive, for example:

```
/* Place section .vectors at address 0 */
place at address Mem:0x0 {readonly section .vectors};
```

Placing a section first or last in a region

To place a section first or last in a region is similar, for example:

```
/* Place section .vectors at start of ROM */
place at start of ROM {readonly section .vectors};
```

Declare and place your own sections

To declare new sections—in addition to the ones used by the IAR build tools—to hold specific parts of your code or data, use mechanisms in the compiler and assembler. For example:

```
/* Place a variable in that section. */
const short MyVariable @ "MYOWNSECTION" = 0xF0F0;
```

This is the corresponding example in assembler language:

name	createSection		
section	MYOWNSECTION: CONST	;	Create a section,
		;	and fill it with
dc16	0xF0F0	;	constant bytes.
end			

To place your new section, the original place in ROM {readonly}; directive is sufficient.

However, to place the section MyOwnSection explicitly, update the linker configuration file with a place in directive, for example:

```
/* Place MyOwnSection in the ROM region */
place in ROM {readonly section MyOwnSection};
```

RESERVING SPACE IN RAM

Often, an application must have an empty uninitialized memory area to be used for temporary storage, for example a heap or a stack. It is easiest to achieve this at link time. You must create a block with a specified size and then place it in a memory.

In the linker configuration file, it can look like this:

```
define block TempStorage with size = 0x1000, alignment = 4 { };
place in RAM { block TempStorage };
```

To retrieve the start of the allocated memory from the application, the source code could look like this:

```
/* Define a section for temporary storage. */
#pragma section = "TEMPSTORAGE"
char *GetTempStorageStartAddress()
{
    /* Return start address of section TEMPSTORAGE. */
    return __section_begin("TEMPSTORAGE");
}
```

KEEPING MODULES

If a module is linked as an object file, it is always kept. That is, it will contribute to the linked application. However, if a module is part of a library, it is included only if it is symbolically referred to from other parts of the application. This is true, even if the library module contains a root symbol. To assure that such a library module is always included, use iarchive to extract the module from the library, see *The IAR Archive Tool—iarchive*, page 437.

For information about included and excluded modules, inspect the log file (the command line option --log modules).

For more information about modules, see Modules and sections, page 70.

KEEPING SYMBOLS AND SECTIONS

By default, ILINK removes any sections, section fragments, and global symbols that are not needed by the application. To retain a symbol that does not appear to be needed—or actually, the section fragment it is defined in—you can either use the root attribute on the symbol in your C/C++ or assembler source code, or use the ILINK option --keep. To retain sections based on attribute names or object names, use the directive keep in the linker configuration file.

To prevent ILINK from excluding sections and section fragments, use the command line options --no_remove or --no_fragments, respectively.

For information about included and excluded symbols and sections, inspect the log file (the command line option --log sections).

For more information about the linking procedure for keeping symbols and sections, see *The linking process*, page 71.

APPLICATION STARTUP

By default, the point where the application starts execution is defined by the __iar_program_start label, which is defined to point at the start of the cstartup.s file. The label is also communicated via ELF to any debugger that is used.

To change the start point of the application to another label, use the ILINK option --entry; see --entry, page 279.

SETTING UP THE STACK

The size of the CSTACK block is defined in the linker configuration file. To change the allocated amount of memory, change the block definition for CSTACK:

define block CSTACK with size = 0x2000, alignment = 8{ };

Specify an appropriate size for your application.

For more information about the stack, see Stack considerations, page 183.

SETTING UP THE HEAP

The size of the heap is defined in the linker configuration file as a block:

define block HEAP with size = 0x1000, alignment = 8{ };
place in RAM {block HEAP};

Specify the appropriate size for your application. If you use a heap, you must allocate at least 50 bytes for it.

SETTING UP THE ATEXIT LIMIT

By default, the atexit function can be called a maximum of 32 times from your application. To either increase or decrease this number, add a line to your configuration file. For example, to reserve room for 10 calls instead, write:

```
define symbol __iar_maximum_atexit_calls = 10;
```

CHANGING THE DEFAULT INITIALIZATION

By default, memory initialization is performed during application startup. ILINK sets up the initialization process and chooses a suitable packing method. If the default initialization process does not suit your application and you want more precise control over the initialization process, these alternatives are available:

- Suppressing initialization
- Choosing the packing algorithm
- Manual initialization
- Initializing code—copying ROM to RAM.

For information about the performed initializations, inspect the log file (the command line option --log initialization).

Suppressing initialization

If you do not want the linker to arrange for initialization by copying, for some or all sections, make sure that those sections do not match a pattern in an initialize by copy directive (or use an except clause to exclude them from matching). If you do not want any initialization by copying at all, you can omit the initialize by copy directive entirely.

This can be useful if your application, or just your variables, are loaded into RAM by some other mechanism before application startup.

Choosing a packing algorithm

To override the default packing algorithm, write for example:

initialize by copy with packing = lzw { readwrite };

For more information about the available packing algorithms, see *Initialize directive*, page 408.

Manual initialization

In the usual case, the initialize by copy directive is used for making the linker arrange for initialization by copying (with or without packing) of sections with content at application startup. The linker achieves this by logically creating an initialization section for each such section, holding the content of the section, and turning the original section into a section without content. Then, the linker adds table elements to the initialization table so that the initialization will be performed at application startup. You can use initialize manually to suppress the creation of table elements to take control over when and how the elements are copied. This is useful for overlays, but also in a number of other circumstances.

For sections without content (zero-initialized sections), the situation is reversed. The linker arranges for zero initialization of all such sections at application startup, except for those that have been mentioned in a do not initialize directive. Usually, only .noinit sections are specified in a do not initialize directive, but you can add any zero-initialized sections you like, and take direct control over when and how these sections are initialized.

Simple copying example with an implicit block

Assume that you have some initialized variables in MYSECTION. If you add this directive to your linker configuration file:

initialize manually { section MYSECTION };

you can use this source code example to initialize the section:

```
#pragma section = "MYSECTION"
#pragma section = "MYSECTION_init"
void DoInit()
{
    char * from = __section_begin("MYSECTION_init");
    char * to = __section_begin("MYSECTION");
    memcpy(to, from, __section_size("MYSECTION"));
}
```

This piece of source code takes advantage of the fact that if you use __section_begin (and related operators) with a section name, a synthetic block is created by the linker for those sections.

Example with explicit blocks

Assume that you instead of needing manual initialization for variables in a specific section, you need it for all initialized variables from a particular library. In that case, you must create explicit blocks for both the variables and the content. Like this:

```
initialize manually { section .data object mylib.a };
define block MYBLOCK { section .data object mylib.a };
define block MYBLOCK_init { section .data_init object mylib.a };
```

You must also place the two new blocks using one of the section placement directives, the block MYBLOCK in RAM and the block MYBLOCK_init in ROM.

Then you can initialize the sections using the same source code as in the previous example, only with MYBLOCK instead of MYSECTION.

Overlay example

This is a simple overlay example that takes advantage of automatic block creation:

initialize manually { section MYOVERLAY* };

define overlay MYOVERLAY { section MYOVERLAY1 }; define overlay MYOVERLAY { section MYOVERLAY2 };

You must also place overlay MYOVERLAY somewhere in RAM. The copying could look like this:

```
#pragma section = "MYOVERLAY"
#pragma section = "MYOVERLAY1_init"
#pragma section = "MYOVERLAY2_init"
void SwitchToOverlay1()
{
    char * from = __section_begin("MYOVERLAY1_init");
    char * to = __section_begin("MYOVERLAY1_init"));
    woid SwitchToOverlay2()
    {
        char * from = __section_begin("MYOVERLAY2_init");
        char * to = __section_begin("MYOVERLAY2_init");
        memcpy(to, from, __section_size("MYOVERLAY2_init"));
    }
```

Initializing code—copying ROM to RAM

Sometimes, an application copies pieces of code from flash/ROM to RAM. You can direct the linker to arrange for this to be done automatically at application startup, or do it yourself at some later time using the techniques described in *Manual initialization*, page 91.

You need to list the code sections that should be copied in an initialize by copy directive. The easiest way is usually to place the relevant functions in a particular section (for example, RAMCODE), and add section RAMCODE to your initialize by copy directive. For example:

initialize by copy { rw, section RAMCODE };

If you need to place the RAMCODE functions in some particular location, you must mention them in a placement directive, otherwise they will be placed together with other read/write sections.

If you need to control the manner and/or time of copying, you must use an initialize manually directive instead. See *Manual initialization*, page 91.

If the functions need to run without accessing the flash/ROM, you can use the ____ramfunc keyword when compiling. See *Execution in RAM*, page 58.

Running all code from RAM

If you want to copy the entire application from ROM to RAM at program startup, use the initilize by copy directive, for example:

initialize by copy { readonly, readwrite };

The readwrite pattern will match all statically initialized variables and arrange for them to be initialized at startup. The readonly pattern will do the same for all read-only code and data, except for code and data needed for the initialization.

To reduce the ROM space that is needed, it might be useful to compress the data with one of the available packing algorithms. For example,

initialize by copy with packing = lzw { readonly, readwrite };

For more information about the available compression algorithms, see *Initialize directive*, page 408.

Because the function __low_level_init, if present, is called before initialization, it, and anything it needs, will not be copied from ROM to RAM either. In some circumstances—for example, if the ROM contents are no longer available to the program after startup—you might need to avoid using the same functions during startup and in the rest of the code.

If anything else should not be copied, include it in an except clause. This can apply to, for example, the interrupt vector table.

It is also recommended to exclude the C++ dynamic initialization table from being copied to RAM, as it is typically only read once and then never referenced again. For example, like this:

INTERACTION BETWEEN ILINK AND THE APPLICATION

ILINK provides the command line options --config_def and --define_symbol to define symbols which can be used for controlling the application. You can also use symbols to represent the start and end of a continuous memory area that is defined in the linker configuration file. For more information, see *Interaction between the tools and your application*, page 186.

To change a reference to one symbol to another symbol, use the ILINK command line option --redirect. This is useful, for example, to redirect a reference from a non-implemented function to a stub function, or to choose one of several different implementations of a certain function, for example, how to choose the DLIB formatter for the standard library functions printf and scanf.

The compiler generates mangled names to represent complex C/C++ symbols. If you want to refer to these symbols from assembler source code, you must use the mangled names.

For information about the addresses and sizes of all global (statically linked) symbols, inspect the entry list in the map file (the command line option --map).

For more information, see Interaction between the tools and your application, page 186.

STANDARD LIBRARY HANDLING

By default, ILINK determines automatically which variant of the standard library to include during linking. The decision is based on the sum of the runtime attributes available in each object file and the library options passed to ILINK.

To disable the automatic inclusion of the library, use the option --no_library_search. In this case, you must explicitly specify every library file to be included. For information about available library files, see *Using prebuilt libraries*, page 101.

PRODUCING OTHER OUTPUT FORMATS THAN ELF/DWARF

ILINK can only produce an output file in the ELF/DWARF format. To convert that format into a format suitable for programming PROM/flash, see *The IAR ELF Tool—ielftool*, page 440.

VENEERS

The ARM cores need to use veneers on two occasions:

- When calling an ARM function from Thumb mode or vice versa; the veneer then changes the state of the microprocessor. The called function must be an interworking function, see *Interworking*, page 48. If the core supports the BLX instruction, a veneer is not needed for changing modes.
- When calling a function that it cannot normally reach; the veneer introduces code which makes the call successfully reach the destination.

Code for veneers can be inserted between any caller and called function. As a result, the R12 register must be treated as a scratch register at function calls, including functions written in assembler. This also applies to jumps.

For more information, see --no_veneers, page 289.

Hints for troubleshooting

ILINK has several features that can help you manage code and data placement correctly, for example:

- Messages at link time, for examples when a relocation error occurs
- The --log option that makes ILINK log information to stdout, which can be useful to understand why an executable image became the way it is, see --*log*, page 284
- The --map option that makes ILINK produce a memory map file, which contains the result of the linker configuration file, see *--map*, page 285.

RELOCATION ERRORS

For each instruction that cannot be relocated correctly, ILINK will generate a *relocation error*. This can occur for instructions where the target is out of reach or is of an incompatible type, or for many other reasons.

A relocation error produced by ILINK can look like this:

```
Error[Lp002]: relocation failed: out of range or illegal value
Kind : R_XXX_YYY[0x1]
Location : 0x40000448
    "myfunc" + 0x2c
Module: somecode.o
Section: 7 (.text)
Offset: 0x2c
Destination: 0x900000c
    "read"
    Module: read.o(iolib.a)
Section: 6 (.text)
Offset: 0x0
```

The message entries are described in this table:

Message entry	Description			
Kind	The relocation directive that failed. The directive depends on the instruction used.			
Location	 The location where the problem occurred, described with the following details: The instruction address, expressed both as a hexadecimal value and as a label with an offset. In this example, 0x40000448 and "myfunc" + 0x2c. The module, and the file. In this example, the module somecode.o. The section number and section name. In this example, section number 7 with the name.text. The offset, specified in number of bytes, in the section. In this 			
	example, 0x2c.			
Destination	 The target of the instruction, described with the following details: The instruction address, expressed both as a hexadecimal value and as a label with an offset. In this example, 0x9000000c and "read" (thus, no offset). The module, and when applicable the library. In this example, the module read.o and the library iolib.a. The section number and section name. In this example, section number 6 with the name .text. The offset, specified in number of bytes, in the section. In this example, 0x0. 			

Table 4: Description of a relocation error

Possible solutions

In this case, the distance from the instruction in myfunc to __read is too long for the branch instruction.

Possible solutions include ensuring that the two .text sections are allocated closer to each other or using some other calling mechanism that can reach the required distance. It is also possible that the referring function tried to refer to the wrong target and that this caused the range error.

Different range errors have different solutions. Usually, the solution is a variant of the ones presented above, in other words modifying either the code or the section placement.

Hints for troubleshooting

The DLIB runtime environment

This chapter describes the runtime environment in which an application executes. In particular, the chapter covers the DLIB runtime library and how you can optimize it for your application.

Introduction to the runtime environment

The runtime environment is the environment in which your application executes. The runtime environment depends on the target hardware, the software environment, and the application code.

RUNTIME ENVIRONMENT FUNCTIONALITY

The *runtime environment* supports Standard C and C++, including the standard template library. The runtime environment consists of the *runtime library*, which contains the functions defined by the C and the C++ standards, and include files that define the library interface (the system header files).

The runtime library is delivered both as prebuilt libraries and (depending on your product package) as source files, and you can find them in the product subdirectories arm\lib and arm\src\lib, respectively.

The runtime environment also consists of a part with specific support for the target system, which includes:

- Support for hardware features:
 - Direct access to low-level processor operations by means of *intrinsic* functions, such as functions for interrupt mask handling
 - Peripheral unit registers and interrupt definitions in include files
 - The Vector Floating Point (VFP) coprocessor.
- Runtime environment support, that is, startup and exit code and low-level interface to some library functions.
- A floating-point environment (fenv) that contains floating-point arithmetics support, see *fenv.h*, page 395.
- Special compiler support, for instance functions for switch handling or integer arithmetics.

For more information about AEABI compliance, see AEABI compliance, page 192.

For more information about the library, see the chapter Library functions.

SETTING UP THE RUNTIME ENVIRONMENT

The IAR DLIB runtime environment can be used as is together with the debugger. However, to run the application on hardware, you must adapt the runtime environment. Also, to configure the most code-efficient runtime environment, you must determine your application and hardware requirements. The more functionality you need, the larger your code will become.

This is an overview of the steps involved in configuring the most efficient runtime environment for your target hardware:

• Choose which runtime library object file to use

It is not necessary to specify a library file explicitly, as ILINK automatically uses the correct library file. See *Using prebuilt libraries*, page 101.

Choose which predefined runtime library configuration to use—Normal or Full

You can configure the level of support for certain library functionality, for example, locale, file descriptors, and multibyte characters. If you do not specify anything, a default library configuration file that matches the library object file is automatically used. To specify a library configuration explicitly, use the --dlib_config compiler option. See *Library configurations*, page 118.

• Optimize the size of the runtime library

You can specify the formatters used by the functions printf, scanf, and their variants, see *Choosing formatters for printf and scanf*, page 105. You can also specify the size and placement of the stacks and the heap, see *Setting up the stack*, page 89, and *Setting up the heap*, page 90, respectively.

• Include debug support for runtime and I/O debugging

The library offers support for mechanisms like redirecting standard input and output to the C-SPY Terminal I/O window and accessing files on the host computer, see *Application debug support*, page 107.

Adapt the library for target hardware

The library uses a set of low-level functions for handling accesses to your target system. To make these accesses work, you must implement your own version of these functions. For example, to make printf write to an LCD display on your board, you must implement a target-adapted version of the low-level function __write, so that it can write characters to the display. To customize such functions, you need a good understanding of the library low-level interface, see *Adapting the library for target hardware*, page 110.

• Override library modules

If you have customized the library functionality, you need to make sure your versions of the library modules are used instead of the default modules. This can be done without rebuilding the entire library, see *Overriding library modules*, page 111.

• Customize system initialization

It is likely that you need to customize the source code for system initialization, for example, your application might need to initialize memory-mapped special function registers, or omit the default initialization of data sections. You do this by customizing the routine __low_level_init, which is executed before the data sections are initialized. See *System startup and termination*, page 113 and *Customizing system initialization*, page 117.

• Configure your own library configuration files

In addition to the prebuilt library configurations, you can make your own library configuration, but that requires that you *rebuild* the library. This gives you full control of the runtime environment. See *Building and using a customized library*, page 112.

• Manage a multithreaded environment

In a multithreaded environment, you must adapt the runtime library to treat all library objects according to whether they are global or local to a thread. See *Managing a multithreaded environment*, page 130.

Check module consistency

You can use runtime model attributes to ensure that modules are built using compatible settings, see *Checking module consistency*, page 136.

Using prebuilt libraries

The prebuilt runtime libraries are configured for different combinations of these features:

- Architecture
- CPU mode
- Byte order
- Library configuration—Normal or Full.
- Floating-point implementation.

The linker will automatically include the correct library object file and library configuration file. To explicitly specify a library configuration, use the --dlib_config option. For more information, see *Runtime environment*, page 50.

LIBRARY FILENAME SYNTAX

The names of the libraries are constructed from these elements:

{architecture}	is the name of the architecture. It can be one of $4t$, $5E$, $6M$, or $7M$ for the ARM architectures v4T, v5TE, v6M, or v7M, respectively. Libraries built for the v5TE architecture are also used for the v6 architecture and later (except for v6M and v7M).	
{cpu_mode}	is one of t or a, for Thumb and ARM, respectively.	
{byte_order}	is one of 1 or $\ensuremath{\mathtt{b}}$ for little-endian and big-endian, respectively	
{fp_implementation}	is _ when the library is compiled without VFP support, that is, using a software implementation floating-point functions,	
	and v when the library is compiled with VFP support for architectures VFPv2 or later. Libraries compiled with VFP support have two entries for each function with floating-point signature. One entry is compliant with the VFP variant of AAPCS base standard. The linker will use the VFP calling convention, and the base standard entry for other modules.	
{language}	is c when the library is compiled for Standard C++ support, and e when compiled for Embedded C++ support.	
{lib_config>	is one of n or ${\tt f}$ for normal and full, respectively.	
{debug_interface}	is one of s, b or i, for the SWI/SVC mechanism, the BKPT mechanism, and the IAR-specific breakpoint mechanism, respectively. For more information, see <i>semihosting</i> , page 292.	
{rwpi}	is s when the library contains read/write position-independent code, see <i>rwpi</i> , page 263.	

Note: There are two library configuration files: DLib_Config_Normal.h and DLib_Config_Full.h.

You can find the library object files and the library configuration files in the subdirectory $arm\lib\$.

GROUPS OF LIBRARY FILES

The libraries are delivered in groups of library functions:

Library files for C library functions

These are the functions defined by Standard C, for example functions like printf and scanf. Note that this library does not include math functions.

The names of the library files are constructed in the following way:

dl{architecture}_{cpu_mode}{byte_order}{lib-config}{rwpi}.a

which more specifically means

 $dl{4t|5E|6M|7M}_{a|t}{1|b}{n|f}{s}.a$

Library files for C++ and Embedded C++ library functions

These are the functions defined by C++, compiled with support for either Standard C++ or Embedded C++.

The names of the library files are constructed in the following way:

dlpp{architecture}_{cpu_mode}{byte_order}{fp_implementation}
{lib-config}{language}.a

which more specifically means

 $dlpp{4t|5E|6M|7M}_{a|t}{1|b}{_|v}{n|f}{c|e}.a$

Library files for math functions

These are the functions for floating-point arithmetic and functions with a floating-point type in its signature as defined by Standard C, for example functions like sqrt.

The names of the library files are constructed in the following way:

m{architecture}_{cpu_mode} {byte_order} {fp_implementation}.a

which more specifically means

 $m\{4t|5E|6M|7M\}_{a|t}\{1|b\}\{_v\}.a$

Library files for runtime support functions

These are functions for system startup, initialization, non floating-point AEABI support routines, and some of the functions that are part of Standard C and C++.

The names of the library files are constructed in the following way:

rt{architecture}_{cpu_mode}{byte_order}.a

which more specifically means

 $rt{4t|5E|6M|7M}_{a|t><1|b}.a$

Library files for debug support functions

These are functions for debug support for the semihosting interface. The names of the library files are constructed in the following way:

sh{debug_interface}_{byte_order}.a

which more specifically means

 $sh{s|b|i}_{1|b}.a$

CUSTOMIZING A PREBUILT LIBRARY WITHOUT REBUILDING

The prebuilt libraries delivered with the compiler can be used as is. However, you can customize parts of a library without rebuilding it.

These items can be customized:

Items that can be customized	Described in			
Formatters for printf and scanf	Choosing formatters for printf and scanf, page 105			
Startup and termination code	System startup and termination, page 113			
Low-level input and output	Standard streams for input and output, page 119			
File input and output	File input and output, page 123			
Low-level environment functions	Environment interaction, page 126			
Low-level signal functions	Signal and raise, page 127			
Low-level time functions	Time, page 127			
Some library math functions	Math functions, page 127			
Size of heaps, stacks, and sections	Stack considerations, page 183 Heap considerations, page 185 Placing code and data—the linker configuration file, page 73			

Table 5: Customizable items

For information about how to override library modules, see *Overriding library modules*, page 111.

Choosing formatters for printf and scanf

The linker automatically chooses an appropriate formatter for printf- and scanf-related function based on information from the compiler. If that information is missing or insufficient, for example if printf is used through a function pointer, if the object file is old, etc, then the automatic choice is the Full formatter. In this case you might want to choose a formatter manually.

To override the default formatter for all the printf- and scanf-related functions, except for wprintf and wscanf variants, you simply set the appropriate library options. This section describes the different options available.

Note: If you rebuild the library, you can optimize these functions even further, see *Configuration symbols for printf and scanf*, page 121.

CHOOSING A PRINTF FORMATTER

The printf function uses a formatter called _Printf. The full version is quite large, and provides facilities not required in many embedded applications. To reduce the memory consumption, three smaller, alternative versions are also provided in the Standard C/EC++ library.

This table summarizes the capabilities of the different formatters:

Formatting capabilities		Small/	Large/	Full/
		SmallNoMb	LargeNoMb	FullNoMb
Basic specifiers c, d, i, o, p, s, u, X, x, and %	Yes	Yes	Yes	Yes
Multibyte support	No	Yes/No	Yes/No	Yes/No
Floating-point specifiers $a, and \mathbb{A}$	No	No	No	Yes
Floating-point specifiers e, E, f, F, g, and ${\rm G}$	No	No	Yes	Yes
Conversion specifier n	No	No	Yes	Yes
Format flag +, -, #, 0, and space	No	Yes	Yes	Yes
Length modifiers h, 1, L, s, t, and ${\tt Z}$	No	Yes	Yes	Yes
Field width and precision, including *	No	Yes	Yes	Yes
long long support	No	No	Yes	Yes

Table 6: Formatters for printf

For information about how to fine-tune the formatting capabilities even further, see *Configuration symbols for printf and scanf*, page 121.



Manually specifying the print formatter in the IDE

To specify a formatter manually, choose **Project>Options** and select the **General Options** category. Select the appropriate option on the **Library options** page.



Manually specifying the printf formatter from the command line

To specify a formatter manually, use one of these ILINK command line options:

--redirect _Printf=_PrintfFull
--redirect _Printf=_PrintfFullNoMb
--redirect _Printf=_PrintfLarge
--redirect _Printf=_PrintfSmall
--redirect _Printf=_PrintfSmallNoMb
--redirect _Printf=_PrintfTiny
--redirect _Printf=_PrintfTinyNoMb

CHOOSING A SCANF FORMATTER

In a similar way to the printf function, scanf uses a common formatter, called _Scanf. The full version is quite large, and provides facilities that are not required in many embedded applications. To reduce the memory consumption, two smaller, alternative versions are also provided in the Standard C/C++ library.

This table summarizes the capabilities of the different formatters:

	Small/	Large/	Full/
Formatting capabilities	SmallNoMB	LargeNoMb	FullNoMb
Basic specifiers c, d, i, o, p, s, u, X, x, and $\%$	Yes	Yes	Yes
Multibyte support	Yes/No	Yes/No	Yes/No
Floating-point specifiers $a, and {\mathbb A}$	No	No	Yes
Floating-point specifiers e, E, f, F, g, and ${\tt G}$	No	No	Yes
Conversion specifier n	No	No	Yes
Scan set [and]	No	Yes	Yes
Assignment suppressing *	No	Yes	Yes
long long support	No	No	Yes

Table 7: Formatters for scanf

For information about how to fine-tune the formatting capabilities even further, see *Configuration symbols for printf and scanf*, page 121.



Manually specifying the scanf formatter in the IDE

To specify a formatter manually, choose **Project>Options** and select the **General Options** category. Select the appropriate option on the **Library options** page.



Manually specifying the scanf formatter from the command line

To specify a formatter manually, use one of these ILINK command line options:

- --redirect _Scanf=_ScanfFull --redirect _Scanf=_ScanfFullNoMb --redirect _Scanf=_ScanfLarge --redirect _Scanf=_ScanfLargeNoMb --redirect _Scanf=_ScanfSmall
- --redirect _Scanf=_ScanfSmallNoMb

Application debug support

In addition to the tools that generate debug information, there is a debug version of the library low-level interface (typically, I/O handling and basic runtime support). Using the debug library, your application can perform things like opening a file on the host computer and redirecting stdout to the debugger Terminal I/O window.

INCLUDING C-SPY DEBUGGING SUPPORT

You can make the library provide debugging support for:

- Handling program abort, exit, and assertions
- I/O handling, which means that stdin and stdout are redirected to the C-SPY Terminal I/O window, and that it is possible to access files on the host computer during debugging.

If you build your application project with the ILINK option **Semihosted** (--semihosted) or **IAR breakpoint** (--semihosting=iar_breakpoint), certain functions in the library are replaced by functions that communicate with the debugger.



To set linker options for debug support in the IDE, choose **Project>Options>General Options**. On the **Library configuration** page, select the Semihosted option or the **IAR breakpoint** option.

Note that for some Cortex-M devices it is also possible to direct stdout/stderr via SWO. This can significantly improve stdout/stderr performance compared to semihosting. For hardware requirement, see the *C-SPY® Debugging Guide for ARM*.



To enable stdout via SWO on the command line, use the linker option --redirect __iar_sh_stdout=__iar_sh_stdout_swo.



To enable stdout via SWO in the IDE, choose **Project>Options>General Options**. On the **Library configuration** page, select the **Semihosted** option and the **stdout/stderr via SWO** option.

THE DEBUG LIBRARY FUNCTIONALITY

The debug library is used for communication between the application being debugged and the debugger itself. The debugger provides runtime services to the application via the low-level DLIB interface; services that allow capabilities like file and terminal I/O to be performed on the host computer.

These capabilities can be valuable during the early development of an application, for example in an application that uses file I/O before any flash file system I/O drivers are implemented. Or, if you need to debug constructions in your application that use stdin and stdout without the actual hardware device for input and output being available. Another use is producing debug printouts.

The low-level debugger runtime interface provided by DLIB is compatible with the semihosting interface provided by ARM Limited. When an application invokes a semihosting call, the execution stops at a debugger breakpoint. The debugger then handles the call, performs any necessary actions on the host computer and then resumes the execution.

The semihosting mechanism

There are three variants of semihosting mechanisms available:

- · For Cortex-M, the interface uses BKPT instructions to perform semihosting calls
- For other ARM cores, SVC instructions are used for the semihosting calls
- IAR breakpoint, which is an IAR-specific alternative to semihosting that uses SVC.

To support semihosting via SVC, the debugger must set its semihosting breakpoint on the Supervisor Call vector to catch SVC calls. If your application uses SVC calls for other purposes than semihosting, the handling of this breakpoint will cause a severe performance penalty for each such call. IAR breakpoint is a way to get around this. By using a special function call instead of an SVC instruction to perform semihosting, the semihosting breakpoint can be set on that special function instead. This means that semihosting will not interfere with other uses of the Supervisor Call vector.

Note that IAR breakpoint is an IAR-specific extension of the semihosting standard. If you link your application with libraries built with toolchains from other vendors than IAR Systems and use IAR breakpoint, semihosting calls made from code in those libraries will not work.

THE C-SPY TERMINAL I/O WINDOW

To make the Terminal I/O window available, the application must be linked with support for I/O debugging. This means that when the functions __read or __write are called to perform I/O operations on the streams stdin, stdout, or stderr, data will be sent to or read from the C-SPY Terminal I/O window.

Note: The Terminal I/O window is not opened automatically just because __read or __write is called; you must open it manually.

For more information about the Terminal I/O window, see the *C-SPY*® *Debugging Guide for ARM*.

Speeding up terminal output

On some systems, terminal output might be slow because the host computer and the target hardware must communicate for each character.

For this reason, a replacement for the __write function called __write_buffered is included in the DLIB library. This module buffers the output and sends it to the debugger one line at a time, speeding up the output. Note that this function uses about 80 bytes of RAM memory.

To use this feature you can either choose **Project>Options>General Options>Library Options** and select the option **Buffered terminal output** in the IDE, or add this to the linker command line:

--redirect __write=__write_buffered

LOW-LEVEL FUNCTIONS IN THE DEBUG LIBRARY

Note: The debug library contains implementations of the following low-level functions:

Function in DLIB low-level interface	Action	
abort	Exits the application	
clock	Returns the clock on the host computer	
close	Closes the associated host file on the host computer	
exit	Notifies that the end of the application was reached	
iar_ReportAssert	Prints an assert message to terminal I/O	
lseek	Searches in the associated host file on the host computer	
open	Opens a file on the host computer	
read	Directs stdin to the Terminal I/O window; all other files will read the associated host file	
remove	Removes a file on the host computer	
rename	Renames a file on the host computer	
system	Writes a message to the Debug Log window and returns -1	
time	Returns the time on the host computer	
write	Directs stdout and stderr to the Terminal I/O window; all other files will write to the associated host file	

Table 8: Functions with special meanings when linked with debug library

You should not use the low-level interface functions prefixed with _ or __ directly in your application. Instead you should use the high-level functions that use these functions. For more information, see *Library low-level interface*, page 111.

Adapting the library for target hardware

The library uses a set of low-level functions for handling accesses to your target system. To make these accesses work, you must implement your own version of these functions. These low-level functions are referred to as the *library low-level interface*.

When you have implemented your low-level interface, you must add your version of these functions to your project. For information about this, see *Overriding library modules*, page 111.

LIBRARY LOW-LEVEL INTERFACE

The library uses a set of low-level functions to communicate with the target system. For example, printf and all other standard output functions use the low-level function __write to send the actual characters to an output device. Most of the low-level functions, like __write, have no implementation. Instead, you must implement them yourself to match your hardware.

However, the library contains a debug version of the library low-level interface, where the low-level functions are implemented so that they interact with the host computer via the debugger, instead of with the target hardware. If you use the debug library, your application can perform tasks like writing to the Terminal I/O window, accessing files on the host computer, getting the time from the host computer, etc. For more information, see *The debug library functionality*, page 108.

Note that your application should not use the low-level functions directly. Instead you should use the corresponding standard library function. For example, to write to stdout, you should use standard library functions like printf or puts, instead of __write.

The library files that you can override with your own versions are located in the arm\src\lib directory.

The low-level interface is further described in these sections:

- Standard streams for input and output, page 119
- File input and output, page 123
- Signal and raise, page 127
- Time, page 127
- Assert, page 130.

Overriding library modules

To use a library low-level interface that you have implemented, add it to your application. See *Adapting the library for target hardware*, page 110. Or, you might want to override a default library routine with your customized version. In both cases, follow this procedure:

- Use a template source file—a library source file or another template—and copy it to your project directory.
- **2** Modify the file.
- **3** Add the customized file to your project, like any other source file.

Note: If you have implemented a library low-level interface and added it to a project that you have built with debug support, your low-level functions will be used and not the C-SPY debug support modules. For example, if you replace the debug support module __write with your own version, the C-SPY Terminal I/O window will not be supported.

The library files that you can override with your own versions are located in the arm\src\lib directory.

Building and using a customized library

Building a customized library is a complex process. Therefore, consider carefully whether it is really necessary. You must build your own C/C++ standard library when you want to define your own library configuration with support for locale, file descriptors, multibyte characters, etc.

In those cases, you must:

- Set up a library project
- Make the required library modifications
- Build your customized library
- Finally, make sure your application project will use the customized library.

Note: To build IAR Embedded Workbench projects from the command line, use the IAR Command Line Build Utility (iarbuild.exe). However, no make or batch files for building the library from the command line are provided.

For information about the build process and the IAR Command Line Build Utility, see the *IDE Project Management and Building Guide for ARM*.

SETTING UP A LIBRARY PROJECT

The IDE provides a library project template which can be used for customizing the runtime environment configuration. This library template uses the Full library configuration, see Table 9, *Library configurations*.



In the IDE, modify the generic options in the created library project to suit your application, see *Basic project configuration*, page 48.

Note: There is one important restriction on setting options. If you set an option on file level (file level override), no options on higher levels that operate on files will affect that file.

MODIFYING THE LIBRARY FUNCTIONALITY

You must modify the library configuration file and build your own library if you want to modify support for, for example, locale, file descriptors, and multibyte characters. This will include or exclude certain parts of the runtime environment.

The library functionality is determined by a set of *configuration symbols*. The default values of these symbols are defined in the file DLib_Defaults.h. This read-only file describes the configuration possibilities. In addition, your library has its own library configuration file, which sets up that specific library with the required library configuration. For more information, see Table 5, *Customizable items*.

The library configuration file is used for tailoring a build of the runtime library, and for tailoring the system header files.

Modifying the library configuration file

In your library project, open the library configuration file and customize it by setting the values of the configuration symbols according to the application requirements.

When you are finished, build your library project with the appropriate project options.

USING A CUSTOMIZED LIBRARY

After you build your library, you must make sure to use it in your application project.

In the IDE you must do these steps:

- I Choose **Project>Options** and click the **Library Configuration** tab in the **General Options** category.
- 2 Choose Custom DLIB from the Library drop-down menu.
- **3** In the **Configuration file** text box, locate your library configuration file.
- **4** Click the **Library** tab, also in the **Linker** category. Use the **Additional libraries** text box to locate your library file.

System startup and termination

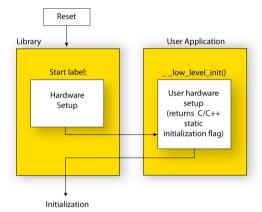
This section describes the runtime environment actions performed during startup and termination of your application.

The code for handling startup and termination is located in the source files cstartup.s, cmain.s, cexit.s, and low_level_init.c or low_level_init.s located in the arm\src\lib directory.

For information about how to customize the system startup code, see *Customizing system initialization*, page 117.

SYSTEM STARTUP

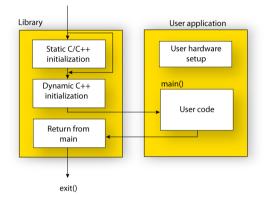
During system startup, an initialization sequence is executed before the main function is entered. This sequence performs initializations required for the target hardware and the C/C++ environment.



For the hardware initialization, it looks like this:

- *Target hardware initialization phase* When the CPU is reset it will jump to the program entry label __iar_program_start in the system startup code.
- Exception stack pointers are initialized to the end of each corresponding section
- The stack pointer is initialized to the end of the CSTACK block
- The function __low_level_init is called if you defined it, giving the application a chance to perform early initializations.

Note: For Cortex-M devices, the second bullet in the above list does not apply. The first and the third bullets are handled slightly differently. At reset, a Cortex-M CPU initializes PC and SP from the vector table (__vector_table), which is defined in the cstartup_M.c file.



For the C/C++ initialization, it looks like this:

- *C/C++ initialization phase*Static and global variables are initialized. That is, zero-initialized variables are cleared and the values of other initialized variables are copied from ROM to RAM memory. This step is skipped if __low_level_init returns zero. For more information, see *Initialization at system startup*, page 76
- Static C++ objects are constructed
- The main function is called, which starts the application.

For information about the initialization phase, see *Application execution—an overview*, page 42.

SYSTEM TERMINATION

This illustration shows the different ways an embedded application can terminate in a controlled way:

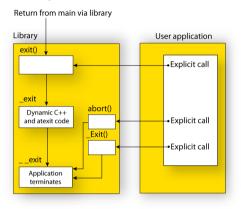


Figure 10: System termination phase

An application can terminate normally in two different ways:

- Return from the main function
- Call the exit function.

Because the C standard states that the two methods should be equivalent, the system startup code calls the exit function if main returns. The parameter passed to the exit function is the return value of main.

The default exit function is written in C. It calls a small assembler function _exit that will perform these operations:

- Call functions registered to be executed when the application ends. This includes C++ destructors for static and global variables, and functions registered with the standard function atexit
- Close all open files
- Call __exit
- When __exit is reached, stop the system.

An application can also exit by calling the abort or the _Exit function. The abort function just calls __exit to halt the system, and does not perform any type of cleanup. The _Exit function is equivalent to the abort function, except for the fact that _Exit takes an argument for passing exit status information.

If you want your application to do anything extra at exit, for example resetting the system, you can write your own implementation of the __exit(int) function.

C-SPY interface to system termination

If your project is linked with the semihosted interface, the normal __exit function is replaced with a special one. C-SPY will then recognize when this function is called and can take appropriate actions to simulate program termination. For more information, see *Application debug support*, page 107.

Customizing system initialization

It is likely that you need to customize the code for system initialization. For example, your application might need to initialize memory-mapped special function registers (SFRs), or omit the default initialization of data sections performed by cstartup.

You can do this by providing a customized version of the routine __low_level_init, which is called from cmain.s before the data sections are initialized. Modifying the file cstartup.s directly should be avoided.

The code for handling system startup is located in the source files cstartup.s and low_level_init.c, located in the arm\src\lib directory.

Note: Normally, you do not need to customize either of the files cmain.s or cexit.s.

If you intend to rebuild the library, the source files are available in the template library project, see *Building and using a customized library*, page 112.

Note: Regardless of whether you modify the routine __low_level_init or the file cstartup.s, you do not have to rebuild the library.

__LOW_LEVEL_INIT

Two skeleton low-level initialization files are supplied with the product: a C source file, <code>low_level_init.c</code> and an alternative assembler source file, <code>low_level_init.s</code>. The latter is part of the prebuilt runtime environment. The only limitation using the C source version is that static initialized variables cannot be used within the file, as variable initialization has not been performed at this point.

The value returned by __low_level_init determines whether or not data sections should be initialized by the system startup code. If the function returns 0, the data sections will not be initialized.

MODIFYING THE FILE CSTARTUP.S

As noted earlier, you should not modify the file cstartup.s if a customized version of __low_level_init is enough for your needs. However, if you do need to modify the file cstartup.s, we recommend that you follow the general procedure for creating a modified copy of the file and adding it to your project, see *Overriding library modules*, page 111.

Note that you must make sure that the linker uses the start label used in your version of cstartup.s. For information about how to change the start label used by the linker, see *--entry*, page 279.

For Cortex-M, you must create a modified copy of cstartup_M.s or cstartup_M.c to use interrupts or other exception handlers.

Library configurations

It is possible to configure the level of support for, for example, locale, file descriptors, multibyte characters.

The runtime library configuration is defined in the *library configuration file*. It contains information about what functionality is part of the runtime environment. The configuration file is used for tailoring a build of a runtime library, and tailoring the system header files used when compiling your application. The less functionality you need in the runtime environment, the smaller it becomes.

The library functionality is determined by a set of *configuration symbols*. The default values of these symbols are defined in the file DLib_Defaults.h. This read-only file describes the configuration possibilities.

These predefined library configurations are available:

Library configuration	Description	
Normal DLIB (default)	No locale interface, C locale, no file descriptor support, and no	
	multibyte characters in printf and scanf.	
Full DLIB	Full locale interface, C locale, file descriptor support, and multibyte characters in printf and scanf.	

Table 9: Library configurations

CHOOSING A RUNTIME CONFIGURATION

To choose a runtime configuration, use one of these methods:

- Default prebuilt configuration—if you do not specify a library configuration explicitly you will get the default configuration. A configuration file that matches the runtime library object file will automatically be used.
- Prebuilt configuration of your choice—to specify a runtime configuration explicitly, use the --dlib_config compiler option. See --dlib_config, page 242.
- Your own configuration—you can define your own configurations, which means that you must modify the configuration file. Note that the library configuration file describes how a library was built and thus cannot be changed unless you rebuild the library. For more information, see *Building and using a customized library*, page 112.

The prebuilt libraries are based on the default configurations, see Table 9, *Library configurations*.

Standard streams for input and output

Standard communication channels (streams) are defined in stdio.h. If any of these streams are used by your application, for example by the functions printf and scanf, you must customize the low-level functionality to suit your hardware.

There are low-level I/O functions, which are the fundamental functions through which C and C++ performs all character-based I/O. For any character-based I/O to be available, you must provide definitions for these functions using whatever facilities the hardware environment provides. For more information about implementing low-level functions, see *Adapting the library for target hardware*, page 110.

IMPLEMENTING LOW-LEVEL CHARACTER INPUT AND OUTPUT

To implement low-level functionality of the stdin and stdout streams, you must write the functions __read and __write, respectively. You can find template source code for these functions in the arm\src\lib directory.

If you intend to rebuild the library, the source files are available in the template library project, see *Building and using a customized library*, page 112. Note that customizing the low-level routines for input and output does not require you to rebuild the library.

Note: If you write your own variants of <u>__read or __write</u>, special considerations for the C-SPY runtime interface are needed, see *Application debug support*, page 107.

Example of using __write

The code in this example uses memory-mapped I/O to write to an LCD display, whose port is assumed to be located at address 0x1000:

```
#include <stddef.h>
__no_init volatile unsigned char lcdIO @ 0x1000;
size_t __write(int handle,
               const unsigned char *buf,
               size_t bufSize)
{
 size_t nChars = 0;
  /* Check for the command to flush all handles */
 if (handle == -1)
  {
   return 0;
 }
  /* Check for stdout and stderr
     (only necessary if FILE descriptors are enabled.) */
 if (handle != 1 && handle != 2)
  {
   return -1;
 }
 for (/* Empty */; bufSize > 0; --bufSize)
  {
   lcdIO = *buf;
   ++buf;
    ++nChars;
 }
 return nChars;
}
```

Note: When DLIB calls __write, DLIB assumes the following interface: a call to __write where buf has the value NULL is a command to flush the stream. When the handle is -1, all streams should be flushed.

Example of using __read

The code in this example uses memory-mapped I/O to read from a keyboard, whose port is assumed to be located at 0x1000:

```
#include <stddef.h>
__no_init volatile unsigned char kbIO @ 0x1000;
size_t __read(int handle,
              unsigned char *buf,
              size_t bufSize)
{
 size_t nChars = 0;
  /* Check for stdin
     (only necessary if FILE descriptors are enabled) */
 if (handle != 0)
  {
   return -1;
  3
 for (/*Empty*/; bufSize > 0; --bufSize)
  ł
   unsigned char c = kbI0;
   if (c == 0)
     break;
    *buf++ = c;
    ++nChars;
  }
 return nChars;
}
```

For information about the @ operator, see *Controlling data and function placement in memory*, page 201.

Configuration symbols for printf and scanf

When you set up your application project, you typically need to consider what printf and scanf formatting capabilities your application requires, see *Choosing formatters* for printf and scanf, page 105.

If the provided formatters do not meet your requirements, you can customize the full formatters. However, that means you must rebuild the runtime library.

The default behavior of the printf and scanf formatters are defined by configuration symbols in the file DLib_Defaults.h.

These configuration symbols determine what capabilities the function printf should have:

Printf configuration symbols	Includes support for
_DLIB_PRINTF_MULTIBYTE	Multibyte characters
_DLIB_PRINTF_LONG_LONG	Long long (11 qualifier)
_DLIB_PRINTF_SPECIFIER_FLOAT	Floating-point numbers
_DLIB_PRINTF_SPECIFIER_A	Hexadecimal floating-point numbers
_DLIB_PRINTF_SPECIFIER_N	Output count (%n)
_DLIB_PRINTF_QUALIFIERS	Qualifiers h, 1, L, v, t, and ${\rm z}$
_DLIB_PRINTF_FLAGS	Flags -, +, #, and 0
_DLIB_PRINTF_WIDTH_AND_PRECISION	Width and precision
_DLIB_PRINTF_CHAR_BY_CHAR	Output char by char or buffered

Table 10: Descriptions of printf configuration symbols

When you build a library, these configurations determine what capabilities the function scanf should have:

Scanf configuration symbols	Includes support for
_DLIB_SCANF_MULTIBYTE	Multibyte characters
_DLIB_SCANF_LONG_LONG	Long long (11 qualifier)
_DLIB_SCANF_SPECIFIER_FLOAT	Floating-point numbers
_DLIB_SCANF_SPECIFIER_N	Output count (%n)
_DLIB_SCANF_QUALIFIERS	Qualifiers h, j, l, t, z, and L
_DLIB_SCANF_SCANSET	Scanset ([*])
_DLIB_SCANF_WIDTH	Width
_DLIB_SCANF_ASSIGNMENT_SUPPRESSING	Assignment suppressing ([*])

Table 11: Descriptions of scanf configuration symbols

CUSTOMIZING FORMATTING CAPABILITIES

To customize the formatting capabilities, you must;

- I Set up a library project, see *Building and using a customized library*, page 112.
- **2** Define the configuration symbols according to your application requirements.

File input and output

The library contains a large number of powerful functions for file I/O operations, such as fopen, fclose, fprintf, fputs, etc. All these functions call a small set of low-level functions, each designed to accomplish one particular task; for example, ___open opens a file, and __write outputs characters. Before your application can use the library functions for file I/O operations, you must implement the corresponding low-level function to suit your target hardware. For more information, see *Adapting the library for target hardware*, page 110.

Note that file I/O capability in the library is only supported by libraries with the full library configuration, see *Library configurations*, page 118. In other words, file I/O is supported when the configuration symbol __DLIB_FILE_DESCRIPTOR is enabled. If not enabled, functions taking a FILE * argument cannot be used.

Template code for these I/O files is included in the product:

I/O function	File	Description
close	close.c	Closes a file.
lseek	lseek.c	Sets the file position indicator.
open	open.c	Opens a file.
read	read.c	Reads a character buffer.
write	write.c	Writes a character buffer.
remove	remove.c	Removes a file.
rename	rename.c	Renames a file.

Table 12: Low-level I/O files

The low-level functions identify I/O streams, such as an open file, with a file descriptor that is a unique integer. The I/O streams normally associated with stdin, stdout, and stderr have the file descriptors 0, 1, and 2, respectively.

Note: If you link your application with I/O debug support, C-SPY variants of the low-level I/O functions are linked for interaction with C-SPY. For more information, see *Application debug support*, page 107.

Locale

Locale is a part of the C language that allows language- and country-specific settings for several areas, such as currency symbols, date and time, and multibyte character encoding.

Depending on what runtime library you are using you get different level of locale support. However, the more locale support, the larger your code will get. It is therefore necessary to consider what level of support your application needs.

The DLIB library can be used in two main modes:

- With locale interface, which makes it possible to switch between different locales during runtime
- Without locale interface, where one selected locale is hardwired into the application.

LOCALE SUPPORT IN PREBUILT LIBRARIES

The level of locale support in the prebuilt libraries depends on the library configuration.

- All prebuilt libraries support the C locale only
- All libraries with *full library configuration* have support for the locale interface. For prebuilt libraries with locale interface, it is by default only supported to switch multibyte character encoding at runtime.
- Libraries with *normal library configuration* do not have support for the locale interface.

If your application requires a different locale support, you must rebuild the library.

CUSTOMIZING THE LOCALE SUPPORT

If you decide to rebuild the library, you can choose between these locales:

- The Standard C locale
- The POSIX locale
- A wide range of European locales.

Locale configuration symbols

The configuration symbol _DLIB_FULL_LOCALE_SUPPORT, which is defined in the library configuration file, determines whether a library has support for a locale interface or not. The locale configuration symbols _LOCALE_USE_LANG_REGION and _ENCODING_USE_ENCODING define all the supported locales and encodings:

#define _DLIB_FULL_LOCALE_	_SUPPORT 1
#define _LOCALE_USE_C	/* C locale */
<pre>#define _LOCALE_USE_EN_US</pre>	/* American English */
<pre>#define _LOCALE_USE_EN_GB</pre>	/* British English */
<pre>#define _LOCALE_USE_SV_SE</pre>	/* Swedish in Sweden */

See DLib_Defaults.h for a list of supported locale and encoding settings.

If you want to customize the locale support, you simply define the locale configuration symbols required by your application. For more information, see *Building and using a customized library*, page 112.

Note: If you use multibyte characters in your C or assembler source code, make sure that you select the correct locale symbol (the local host locale).

Building a library without support for locale interface

The locale interface is not included if the configuration symbol __DLIB_FULL_LOCALE_SUPPORT is set to 0 (zero). This means that a hardwired locale is used—by default the Standard C locale—but you can choose one of the supported locale configuration symbols. The setlocale function is not available and can therefore not be used for changing locales at runtime.

Building a library with support for locale interface

Support for the locale interface is obtained if the configuration symbol __DLIB_FULL_LOCALE_SUPPORT is set to 1. By default, the Standard C locale is used, but you can define as many configuration symbols as required. Because the setlocale function will be available in your application, it will be possible to switch locales at runtime.

CHANGING LOCALES AT RUNTIME

The standard library function setlocale is used for selecting the appropriate portion of the application's locale when the application is running.

The setlocale function takes two arguments. The first one is a locale category that is constructed after the pattern LC_CATEGORY. The second argument is a string that describes the locale. It can either be a string previously returned by setlocale, or it can be a string constructed after the pattern:

lang_REGION

or

lang_REGION.encoding

The *lang* part specifies the language code, and the *REGION* part specifies a region qualifier, and *encoding* specifies the multibyte character encoding that should be used.

The *lang_REGION* part matches the _LOCALE_USE_*LANG_REGION* preprocessor symbols that can be specified in the library configuration file.

Example

This example sets the locale configuration symbols to Swedish to be used in Finland and UTF8 multibyte character encoding:

```
setlocale (LC_ALL, "sv_FI.Utf8");
```

Environment interaction

According to the C standard, your application can interact with the environment using the functions getenv and system.

Note: The putenv function is not required by the standard, and the library does not provide an implementation of it.

THE GETENV FUNCTION

The getenv function searches the string, pointed to by the global variable __environ, for the key that was passed as argument. If the key is found, the value of it is returned, otherwise 0 (zero) is returned. By default, the string is empty.

To create or edit keys in the string, you must create a sequence of null terminated strings where each string has the format:

key=value\0

End the string with an extra null character (if you use a C string, this is added automatically). Assign the created sequence of strings to the __environ variable.

For example:

const char MyEnv[] = "Key=Value\0Key2=Value2\0"; __environ = MyEnv;

If you need a more sophisticated environment variable handling, you should implement your own getenv, and possibly putenv function. This does not require that you rebuild the library. You can find source templates in the files getenv.c and environ.c in the arm\src\lib directory. For information about overriding default library modules, see *Overriding library modules*, page 111.

THE SYSTEM FUNCTION

If you need to use the system function, you must implement it yourself. The system function available in the library simply returns -1.

If you decide to rebuild the library, you can find source templates in the library project template. For more information, see *Building and using a customized library*, page 112.

Note: If you link your application with support for I/O debugging, the system function is replaced by a C-SPY variant. For more information, see *Application debug support*, page 107.

Signal and raise

Default implementations of the functions signal and raise are available. If these functions do not provide the functionality that you need, you can implement your own versions.

This does not require that you rebuild the library. You can find source templates in the files signal.c and raise.c in the arm\src\lib directory. For information about overriding default library modules, see *Overriding library modules*, page 111.

If you decide to rebuild the library, you can find source templates in the library project template. For more information, see *Building and using a customized library*, page 112.

Time

To make the __time32, __time64, and date functions work, you must implement the functions clock, __time32, __time64, and __getzone. Whether you use __time32 or __time64 depends on which interface you use for time_t, see *time.h*, page 396.

To implement these functions does not require that you rebuild the library. You can find source templates in the files clock.c, time.c, time64.c, and getzone.c in the arm\src\lib directory. For information about overriding default library modules, see *Overriding library modules*, page 111.

If you decide to rebuild the library, you can find source templates in the library project template. For more information, see *Building and using a customized library*, page 112.

The default implementation of __getzone specifies UTC (Coordinated Universal Time) as the time zone.

Note: If you link your application with support for I/O debugging, the functions clock and time are replaced by C-SPY variants that return the host clock and time respectively. For more information, see *Application debug support*, page 107.

Math functions

Some library math functions are also available size-optimized versions, and in more accurate versions.

SMALLER VERSIONS

The functions cos, exp, log, log10, pow, sin, tan, and __iar_Sin (a help function for sin and cos) exist in additional, smaller versions in the library. They are about 20% smaller and about 20% faster than the default versions. The functions handle INF and NaN values. The drawbacks are that they almost always lose some precision and they do not have the same input range as the default versions.

The names of the functions are constructed like:

___iar_xxx_small<f|1>

where f is used for float variants, l is used for long double variants, and no suffix is used for double variants.

To use these functions, the default function names must be redirected to these names when linking, using the following options:

```
--redirect sin=__iar_sin_small
--redirect cos=__iar_cos_small
--redirect tan=__iar_tan_small
--redirect log=__iar_log_small
--redirect log10=__iar_log10_small
--redirect exp= iar exp small
--redirect pow=__iar_pow_small
--redirect __iar_Sin=__iar_Sin_small
--redirect sinf=__iar_sin_smallf
--redirect cosf=__iar_cos_smallf
--redirect tanf=__iar_tan_smallf
--redirect logf=__iar_log_smallf
--redirect log10f=__iar_log10_smallf
--redirect expf=__iar_exp_smallf
--redirect powf=__iar_pow_smallf
--redirect __iar_Sinf=__iar_Sin_smallf
--redirect sinl=__iar_sin_small1
--redirect cosl=__iar_cos_small1
--redirect tanl=__iar_tan_small1
--redirect log1=__iar_log_small1
--redirect log101=__iar_log10_small1
--redirect expl=__iar_exp_small1
--redirect powl=__iar_pow_small1
--redirect __iar_Sinl=__iar_Sin_small1
```

Note that if cos or sin is redirected, __iar_Sin must be redirected as well.

MORE ACCURATE VERSIONS

The functions cos, pow, sin, and tan, and the help functions __iar_Sin and __iar_Pow exist in versions in the library that are more exact and can handle larger argument ranges. The drawback is that they are larger and slower than the default versions.

The names of the functions are constructed like:

```
___iar_xxx_accurate<f|1>
```

where f is used for float variants, l is used for long double variants, and no suffix is used for double variants.

To use these functions, the default function names must be redirected to these names when linking, using the following options:

--redirect sin=__iar_sin_accurate --redirect cos=__iar_cos_accurate --redirect tan= iar tan accurate --redirect pow=__iar_pow_accurate --redirect __iar_Sin=__iar_Sin_accurate --redirect __iar_Pow=__iar_Pow_accurate --redirect sinf=__iar_sin_accuratef --redirect cosf=__iar_cos_accuratef --redirect tanf=__iar_tan_accuratef --redirect powf=__iar_pow_accuratef --redirect __iar_Sinf=__iar_Sin_accuratef --redirect __iar_Powf=__iar_Pow_accuratef --redirect sinl=__iar_sin_accuratel --redirect cosl=__iar_cos_accuratel --redirect tanl=__iar_tan_accuratel --redirect powl=__iar_pow_accuratel --redirect __iar_Sinl=__iar_Sin_accuratel --redirect __iar_Powl=__iar_Pow_accuratel

Note that if cos or sin is redirected, __iar_Sin must be redirected as well. The same applies to pow and __iar_Pow.

Assert

If you linked your application with support for runtime debugging, an assert will print a message on stdout. If this is not the behavior you require, you must add the source file xreportassert.c to your application project. The __iar_ReportAssert function generates the assert notification. You can find template code in the arm\src\lib

directory. For more information, see *Building and using a customized library*, page 112. To turn off assertions, you must define the symbol NDEBUG.



In the IDE, this symbol NDEBUG is by default defined in a Release project and *not* defined in a Debug project. If you build from the command line, you must explicitly define the symbol according to your needs. See *NDEBUG*, page 388.

Atexit

The linker allocates a static memory area for atexit function calls. By default, the number of calls to the atexit function are limited to 32. To change this limit, see *Setting up the atexit limit*, page 90.

Managing a multithreaded environment

In a multithreaded environment, the standard library must treat all library objects according to whether they are global or local to a thread. If an object is a true global object, any updates of its state must be guarded by a locking mechanism to make sure that only one thread can update it at any given time. If an object is local to a thread, the static variables containing the object state must reside in a variable area local to that thread. This area is commonly named *thread-local storage* (TLS).

There are three possible scenarios, and you need to consider which one that applies to you:

- If you are using an RTOS that supports the multithreading provided by the DLIB library, the RTOS and the DLIB library will handle multithreading which means you do not need to adapt the DLIB library.
- If you are using an RTOS that does not support or only partly supports the multithreading provided by the DLIB library, you probably need to adapt both the RTOS and the DLIB library.
- If you are not using an RTOS, you must adapt the DLIB library to get support for multithreading.

MULTITHREAD SUPPORT IN THE DLIB LIBRARY

The DLIB library uses two kinds of locks—*system locks* and *file stream locks*. The file stream locks are used as guards when the state of a file stream is updated, and are only needed in the Full library configuration. The following library objects are guarded with system locks:

• The heap, in other words when malloc, new, free, delete, realloc, or calloc is used.

- The file system (only available in the Full library configuration), but not the file streams themselves. The file system is updated when a stream is opened or closed, in other words when fopen, fclose, fdopen, fflush, or freopen is used.
- The signal system, in other words when signal is used.
- The temporary file system, in other words when tmpnam is used.
- Initialization of static function objects.

These library objects use TLS:

localeconv, setlocale	
asctime, localtime, gmtime, mktime	
mbrlen, mbrtowc, mbsrtowc, mbtowc, wcrtomb, wcsrtomb, wctomb	
rand, srand	
atexit, strtok	

Table 13: Library objects using TLS

ENABLING MULTITHREAD SUPPORT

To enable multithread support in the library, you must:

- Implement code for the library's system locks interface
- If file streams are used, implement code for the library's file stream locks interface or redirect the interface to the system locks interface (using the linker option --redirect)
- Implement source code that handles thread creation, thread destruction, and TLS access methods for the library
- Modify the linker configuration file accordingly
- If any of the C++ variants are used, use the compiler option --guard_calls. Otherwise, function static variables with dynamic initializers might be initialized simultaneously by several threads.

You can find the required declaration of functions and definitions of macros in the DLib_Threads.h file, which is included by yvals.h.

System locks interface

This interface must be fully implemented for system locks to work:

The lock and unlock implementation must survive nested calls.

File streams locks interface

This interface is only needed for the Full library configuration. If file streams are used, they can either be fully implemented or they can be redirected to the system locks interface. This interface must be implemented for file streams locks to work:

```
typedef void *__iar_Rmtx; /* Lock info object */
void __iar_file_Mtxinit(__iar_Rmtx *);/*Initialize a file lock */
void __iar_file_Mtxdst(__iar_Rmtx *); /* Destroy a file lock */
void __iar_file_Mtxlock(__iar_Rmtx *); /* Lock a file lock */
void __iar_file_Mtxunlock(__iar_Rmtx *); /* Unlock a file lock */
```

The lock and unlock implementation must survive nested calls.

DLIB lock usage

The number of locks that the DLIB library assumes exist are:

- _FOPEN_MAX—the maximum number of file stream locks. These locks are only used in the Full library configuration, in other words only if both the macro symbols _DLIB_FILE_DESCRIPTOR and _FILE_OP_LOCKS are true.
- _MAX_LOCK—the maximum number of system locks.

Note that even if the application uses fewer locks, the DLIB library will initialize and destroy all of the locks above.

For information about the initialization and destruction code, see xsyslock.c.

TLS handling

The DLIB library supports TLS memory areas for two types of threads: the *main thread* (the main function including the system startup and exit code) and *secondary threads*.

The main thread's TLS memory area:

- Is automatically created and initialized by your application's startup sequence
- Is automatically destructed by the application's destruct sequenceIs located in the section __DLIB_PERTHREAD
- Exists also for non-threaded applications.

Each secondary thread's TLS memory area:

- Must be manually created and initialized
- Must be manually destructed
- Is located in a manually allocated memory area.

If you need the runtime library to support secondary threads, you must override the function:

void *__iar_dlib_perthread_access(void *symbp);

The parameter is the address to the TLS variable to be accessed—in the main thread's TLS area—and it should return the address to the symbol in the current TLS area.

Two interfaces can be used for creating and destroying secondary threads. You can use the following interface that allocates a memory area on the heap and initializes it. At deallocation, it destroys the objects in the area and then frees the memory.

```
void *__iar_dlib_perthread_allocate(void);
void __iar_dlib_perthread_deallocate(void *);
```

Alternatively, if the application handles the TLS allocation, you can use this interface for initializing and destroying the objects in the memory area:

```
void ___iar_dlib_perthread_initialize(void *);
void ___iar_dlib_perthread_destroy(void *);
```

These macros can be helpful when you implement an interface for creating and destroying secondary threads:

Macro	Description
IAR_DLIB_PERTHREAD_SIZE	The size needed for the TLS memory area.
IAR_DLIB_PERTHREAD_INIT_SIZE	The initializer size for the TLS memory area. You should initialize the rest of the TLS memory area, up to IAR_DLIB_PERTHREAD_SIZE to zero.
IAR_DLIB_PERTHREAD_SYMBOL_ OFFSET(<i>symbolptr</i>)	The offset to the symbol in the TLS memory area.

Table 14: Macros for implementing TLS allocation

Note that the size needed for TLS variables depends on which DLIB resources your application uses.

This is an example of how you can handle threads:

```
#include <vvals.h>
/* A thread's TLS pointer */
void _DLIB_TLS_MEMORY *TLSp;
/* Are we in a secondary thread? */
int InSecondaryThread = 0;
/* Allocate a thread-local TLS memory
   area and store a pointer to it in TLSp. */
void AllocateTLS()
{
 TLSp = __iar_dlib_perthread_allocate();
}
/* Deallocate the thread-local TLS memory area. */
void DeallocateTLS()
{
  ___iar_dlib_perthread_deallocate(TLSp);
}
/* Access an object in the
   thread-local TLS memory area. */
void _DLIB_TLS_MEMORY *__iar_dlib_perthread_access(
 void _DLIB_TLS_MEMORY *symbp)
{
 char _DLIB_TLS_MEMORY *p = 0;
 if (InSecondaryThread)
   p = (char _DLIB_TLS_MEMORY *) TLSp;
 else
   p = (char _DLIB_TLS_MEMORY *)
      __segment_begin("__DLIB_PERTHREAD");
 p += __IAR_DLIB_PERTHREAD_SYMBOL_OFFSET(symbp);
 return (void _DLIB_TLS_MEMORY *) p;
}
```

The TLSp variable is unique for each thread, and must be exchanged by the RTOS or manually whenever a thread switch occurs.

TLS IN THE LINKER CONFIGURATION FILE

Normally, the linker automatically chooses how to initialize static data. If threads are used, the main thread's TLS memory area must be initialized by plain copying because the initializers are used for each secondary thread's TLS memory area as well. This is controlled by the following statement in your linker configuration file:

initialize by copy with packing = none {section __DLIB_PERTHREAD };

Checking module consistency

This section introduces the concept of runtime model attributes, a mechanism used by the tools provided by IAR Systems to ensure that modules that are linked into an application are compatible, in other words, are built using compatible settings. The tools use a set of predefined runtime model attributes. In addition to these, you can define your own that you can use to ensure that incompatible modules are not used together.

Note: In addition to the predefined attributes, compatibility is also checked against the AEABI runtime attributes. These attributes deal mainly with object code compatibility, etc. They reflect compilation settings and are not user-configurable.

RUNTIME MODEL ATTRIBUTES

A runtime attribute is a pair constituted of a named key and its corresponding value. In general, two modules can only be linked together if they have the same value for each key that they both define.

There is one exception: if the value of an attribute is *, then that attribute matches any value. The reason for this is that you can specify this in a module to show that you have considered a consistency property, and this ensures that the module does not rely on that property.

Note: For IAR predefined runtime model attributes, the linker checks them in several ways.

Example

In this table, the object files could (but do not have to) define the two runtime attributes color and taste:

Object file	Color	Taste
file1	blue	not defined
file2	red	not defined

Table 15: Example of runtime model attributes

Object file	Color	Taste
file3	red	*
file4	red	spicy
file5	red	lean

Table 15: Example of runtime model attributes (Continued)

In this case, file1 cannot be linked with any of the other files, since the runtime attribute color does not match. Also, file4 and file5 cannot be linked together, because the taste runtime attribute does not match.

On the other hand, file2 and file3 can be linked with each other, and with either file4 or file5, but not with both.

USING RUNTIME MODEL ATTRIBUTES

To ensure module consistency with other object files, use the #pragma rtmodel directive to specify runtime model attributes in your C/C++ source code. For example, if you have a UART that can run in two modes, you can specify a runtime model attribute, for example uart. For each mode, specify a value, for example model and mode2. Declare this in each module that assumes that the UART is in a particular mode. This is how it could look like in one of the modules:

#pragma rtmodel="uart", "model"

Alternatively, you can also use the rtmodel assembler directive to specify runtime model attributes in your assembler source code. For example:

rtmodel "uart", "model"

Note that key names that start with two underscores are reserved by the compiler. For more information about the syntax, see *rtmodel*, page 338 and the *IAR Assembler Reference Guide for ARM*, respectively.

At link time, the IAR ILINK Linker checks module consistency by ensuring that modules with conflicting runtime attributes will not be used together. If conflicts are detected, an error is issued.

Assembler language interface

When you develop an application for an embedded system, there might be situations where you will find it necessary to write parts of the code in assembler, for example when using mechanisms in the ARM core that require precise timing and special instruction sequences.

This chapter describes the available methods for this and some C alternatives, with their advantages and disadvantages. It also describes how to write functions in assembler language that work together with an application written in C or C++.

Finally, the chapter covers how functions are called, and how you can implement support for call frame information in your assembler routines for use in the C-SPY® Call Stack window.

Mixing C and assembler

The IAR C/C++ Compiler for ARM provides several ways to access low-level resources:

- Modules written entirely in assembler
- Intrinsic functions (the C alternative)
- Inline assembler.

It might be tempting to use simple inline assembler. However, you should carefully choose which method to use.

INTRINSIC FUNCTIONS

The compiler provides a few predefined functions that allow direct access to low-level processor operations without having to use the assembler language. These functions are known as intrinsic functions. They can be very useful in, for example, time-critical routines.

An intrinsic function looks like a normal function call, but it is really a built-in function that the compiler recognizes. The intrinsic functions compile into inline code, either as a single instruction, or as a short sequence of instructions.

The advantage of an intrinsic function compared to using inline assembler is that the compiler has all necessary information to interface the sequence properly with register allocation and variables. The compiler also knows how to optimize functions with such sequences; something the compiler is unable to do with inline assembler sequences. The result is that you get the desired sequence properly integrated in your code, and that the compiler can optimize the result.

For more information about the available intrinsic functions, see the chapter *Intrinsic functions*.

MIXING C AND ASSEMBLER MODULES

It is possible to write parts of your application in assembler and mix them with your C or C++ modules. This gives several benefits compared to using inline assembler:

- The function call mechanism is well-defined
- The code will be easy to read
- The optimizer can work with the C or C++ functions.

This causes some overhead in the form of a function call and return instruction sequences, and the compiler will regard some registers as scratch registers. However, the compiler will also assume that all scratch registers are destroyed by an inline assembler instruction. In many cases, the overhead of the extra instructions can be removed by the optimizer.

An important advantage is that you will have a well-defined interface between what the compiler produces and what you write in assembler. When using inline assembler, you will not have any guarantees that your inline assembler lines do not interfere with the compiler generated code.

When an application is written partly in assembler language and partly in C or C++, you are faced with several questions:

- How should the assembler code be written so that it can be called from C?
- Where does the assembler code find its parameters, and how is the return value passed back to the caller?
- How should assembler code call functions written in C?
- How are global C variables accessed from code written in assembler language?
- Why does not the debugger display the call stack when assembler code is being debugged?

The first question is discussed in the section *Calling assembler routines from C*, page 147. The following two are covered in the section *Calling convention*, page 150.

The answer to the final question is that the call stack can be displayed when you run assembler code in the debugger. However, the debugger requires information about the call frame, which must be supplied as annotations in the assembler source file. For more information, see *Call frame information*, page 157.

The recommended method for mixing C or C++ and assembler modules is described in *Calling assembler routines from C*, page 147, and *Calling assembler routines from* C++, page 149, respectively.

INLINE ASSEMBLER

Inline assembler can be used for inserting assembler instructions directly into a C or C++ function. Typically, this can be useful if you need to:

- Access hardware resources that are not accessible in C (in other words, when there
 is no definition for an SFR or there is no suitable intrinsic function available).
- Manually write a time-critical sequence of code that if written in C will not have the right timing.
- Manually write a speed-critical sequence of code that if written in C will be too slow.

An inline assembler statement is similar to a C function in that it can take input arguments (input operands), have return values (output operands), and read or write to C symbols (via the operands). An inline assembler statement can also declare *clobbered resources* (that is, values in registers and memory that have been overwritten).

Limitations

Most things you can to do in normal assembler language are also possible with inline assembler, with the following differences:

- Alignment cannot be controlled; this means, for example, that DC32 directives might be misaligned.
- The only accepted register synonyms are SP (for R13), LR (for R14), and PC (for R15).
- In general, assembler directives will cause errors or have no meaning. However, data definition directives will work as expected.
- Resources used (registers, memory, etc) that are also used by the C compiler must be declared as operands or clobbered resources.
- If you do not want to risk that the inline assembler statement to be optimized away by the compiler, you must declare it volatile.

- Accessing a C symbol or using a constant expression requires the use of operands.
- Dependencies between the expressions for the operands might result in an error.

Risks with inline assembler

Without operands and clobbered resources, inline assembler statements have no interface with the surrounding C source code. This makes the inline assembler code fragile, and might also become a maintenance problem if you update the compiler in the future. There are also several limitations to using inline assembler without operands and clobbered resources:

- The compiler's various optimizations will disregard any effects of the inline statements, which will not be optimized at all.
- The inline assembler statement will be volatile and *clobbered memory* is implied. This means that the compiler will not remove the assembler statement. It will simply be inserted at the given location in the program flow. The consequences or side-effects that the insertion might have on the surrounding code are not taken into consideration. If, for example, registers or memory locations are altered, they might have to be restored within the sequence of inline assembler instructions for the rest of the code to work properly.



The following example (for ARM mode) demonstrates the risks of using the asm keyword without operands and clobbers:

```
int Add(int term1, int term2)
{
   asm("adds r0,r0,r1");
   return term1;
}
```

In this example:

- The function Add assumes that values are passed and returned in registers in a way that they might not always be, for example if the function is inlined.
- The s in the adds instruction implies that the condition flags are updated, which should be specified using the cc clobber operand. Otherwise, the compiler will assume that the condition flags are not modified.

Inline assembler without using operands or clobbered resources is therefore often best avoided.

Reference information for inline assembler

The asm and __asm keywords both insert inline assembler instructions. However, when you compile C source code, the asm keyword is not available when the option --strict is used. The __asm keyword is always available.

```
Syntax
                         The syntax of an inline assembler statement is (similar to the one used by GNU gcc):
                         asm [volatile]( string [assembler-interface])
                         string can contain one or more valid assembler instructions or data definition
                         assembler directives, separated by n.
                         For example:
                         asm("label:nop\n"
                             "b label");
                         Note that you can define and use local labels in inline assembler instructions.
                         assembler-interface is:
                           : comma-separated list of output operands /* optional */
                           : comma-separated list of input operands
                                                                               /* optional */
                           : comma-separated list of clobbered resources /* optional */
Operands
                         An inline assembler statement can have one input and one output comma-separated list
                         of operands. Each operand consists of an optional symbolic name in brackets, a quoted
                         constraint, followed by a C expression in parentheses.
Syntax of operands
                         [[ symbolic-name ]] "[modifiers]constraint" (expr)
                         For example:
                         int Add(int term1, int term2)
                         {
                           int sum;
                           asm("add %0,%1,%2"
                               : "=r"(sum)
                               : "r" (term1), "r" (term2));
                           return sum;
                         }
                         In this example, the assembler instruction uses one output operand, sum, two input
                         operands, term1 and term2, and no clobbered resources.
                         It is possible to omit any list by leaving it empty. For example:
                         int matrix[M][N];
                         void MatrixPreloadRow(int row)
                         {
                           asm volatile ("pld [%0]" : : "r" (&matrix[row][0]));
                         }
```

Operand constraints

Operand constraints specify what kind of operand to use, such as a register or an immediate constant. A constraint can also specify what type of register to use, or which constant values are valid.

Constraint	Description
r	Uses a general purpose register for the expression: R0-R12, R14 (for ARM and Thumb2) R0-R7 (for Thumb1)
1	R0-R7 (only valid for Thumb1)
Rp	Uses a pair of general purpose registers, for example ${\tt R0, R1}$
i	An immediate integer operand with a constant value. Symbolic constants are allowed.
j	A 16-bit constant suitable for a MOVW instruction (valid for ARM and Thumb2).
n	An immediate operand, alias for i.
I	An immediate constant valid for a data processing instruction (for ARM and Thumb2), or a constant in the range 0 to 255 (for Thumb1).
J	An immediate constant in the range -4095 to 4095 (for ARM and Thumb2), or a constant in the range -255 to -1 (for Thumb1).
K	An immediate constant that satisfies the I constraint if inverted (for ARM and Thumb2), or a constant that satisfies the I constraint multiplied by any power of 2 (for Thumb1).
L	An immediate constant that satisfies the ${\tt I}$ constraint if negated (for ARM and Thumb2), or a constant in the range -7 to 7 (for Thumb1).
М	An immediate constant that is a multiple of 4 in the range 0 to 1020 (only valid for Thumb1).
N	An immediate constant in the range 0 to 31 (only valid for Thumb1).
0	An immediate constant that is a multiple of 4 in the range -508 to 508 (only valid for Thumb1).
t	An S register.
W	A D register.
q	A Q register.
Dv	A 32-bit floating-point immediate constant for the \texttt{VMOV} . $\texttt{F32}$ instruction.
Dy	A 64-bit floating-point immediate constant for the \mathtt{VMOV} . $\mathtt{F64}$ instruction.

Table 16: Inline assembler operand constraints

	Constraint	Description
	v2s v4Q	A vector of 2, 3, or 4 consecutive S, D, or Q registers. For example,
		$v4\ensuremath{\mathbb{Q}}$ is a vector of four $\ensuremath{\mathbb{Q}}$ registers. The vectors do not overlap, so the
		available $v4\text{Q}$ register vectors are <code>Q0-Q3, Q4-Q7, Q8-Q11</code> , and
		Q12-Q15.
	Table 16: Inline assemble	er operand constraints (Continued)
Constraint modifiers		s can be used together with a constraint to modify its meaning. This rted constraint modifiers:
	Modifier	Description
	=	Write-only operand
	+	Read-write operand
	&	Early clobber output operand which is written to before the instruction
		has processed all the input operands.
	Table 17: Supported cons	traint modifiers
Referring to operands	Assembler instructions refer to operands by prefixing their order number with %. The first operand has order number 0 and is referred to by %0.	
% [operand.name]. Sym code and can be the same a		symbolic name, you can refer to it using the syntax . Symbolic operand names are in a separate namespace from C/C++ same as a C/C++ variable names. Each operand name must however ssembler statement. For example:
	<pre>int Add(int term1, int term2) {</pre>	
	int sum;	
	: [Rd]"=r	l],%[Rn],%[Rm]" "(sum) (term1), [Rm]"r" (term2));
	return sum; }	

Operand modifiers An operand modifier is a single letter between the % and the operand number, which is used for transforming the operand. In the example below, the modifiers L and H are used for accessing the least and most significant 16 bits, respectively, of an immediate operand:

```
int Mov32()
{
    int a;
    asm("movw %0,%L1 \n"
        "movt %0,%H1 \n" : "=r"(a) : "i"(0x12345678UL));
    return a;
}
```

Some operand modifiers can be combined, in which case each letter will transform the result from the previous modifier. This table describes the transformation performed by each valid modifier:

Modifier	Description
L	The lowest-numbered register of a register pair, or the low 16 bits of an immediate constant.
Н	The highest-numbered register of a register pair, or the high 16 bits of an immediate constant.
С	For an immediate operand, an integer or symbol address without a preceding # sign. Cannot be transformed by additional operand modifiers.
В	For an immediate operand, the bitwise inverse of integer or symbol without a preceding # sign. Cannot be transformed by additional operand modifiers.
Q	The least significant register of a register pair.
R	The most significant register of a register pair.
М	For a register or a register pair, the register list suitable for $1dm$ or stm. Cannot be transformed by additional operand modifiers.
a	Transforms a register Rn into a memory operand [Rn , #0] suitable for pld.
b	The low ${\rm S}$ register part of a ${\rm D}$ register.
р	The high ${\rm S}$ register part of a ${\rm D}$ register.
e	The low $\mathbb D$ register part of a $\mathbb Q$ register, or the low register in a vector of Neon registers.
f	The high ${\mathbb D}$ register part of a ${\mathbb Q}$ register, or the high register in a vector of Neon registers.

Table 18: Operand modifiers and transformations

	Modifier	Description
	h	For a (vector of) D or Q registers, the corresponding list of D registers within curly braces. For example, $Q0$ becomes { $D0$, $D1$ }. Cannot be transformed by additional operand modifiers.
	У	S register as indexed D register, for example S7 becomes D3 [1]. Cannot be transformed by additional operand modifiers.
	Table 18: Operand modij	fiers and transformations
Input operands	Input operands cannot have any modifiers, but they can have any valid C expression as long as the type of the expression fits the register.	
	-	vill be evaluated just before any of the assembler instructions in the tement and assigned to the constraint, for example a register.
Output operands	Output operands must have = as a modifier and the C expression must be an l-value and specify a writable location. For example, $=r$ for a write-only general purpose register. The constraint will be assigned to the evaluated C expression (as an l-value) immediately after the last assembler instruction in the inline assembler statement. Input operands are assumed to be consumed before output is produced and the compiler may use the same register for an input and output operand. To prohibit this, prefix the output constraint with & to make it an early clobber resource, for example $=\&r$. This will ensure that the output operand will be allocated in a different register than the input operands.	
Input/output operands	An operand that should be used both for input and output must be listed as an outpu operand and have the + modifier. The C expression must be an l-value and specify a writable location. The location will be read immediately before any assembler instructions and it will be written to right after the last assembler instruction.	
	This is an example of using a read-write operand:	
	int Double(int value)	
	{ asm("add %0,%0,%0" : "+r"(value)); return value; }	
		we, the input value for value will be placed in a general purpose assembler statement, the result from the ADD instruction will be placed .
Clobbered resources	An inline assemble	r statement can have a list of clobbered resources.
	"resource1", "ı	resource2",

Specify clobbered resources to inform the compiler about which resources the inline assembler statement destroys. Any value that resides in a clobbered resource and that is needed after the inline assembler statement will be reloaded.

Clobbered resources will not be used as input or output operands.

This is an example of how to use clobbered resources:

```
int Add(int term1, int term2)
{
    int sum;
    asm("adds %0,%1,%2"
        : "=r"(sum)
        : "r" (term1), "r" (term2)
        : "cc");
    return sum;
}
```

In this example the condition codes will be modified by the ADDS instruction. Therefore, "cc" must be listed in the clobber list.

This table lists valid clobbered resources:

Clobber	Description
R0-R12, R14 for ARM mode and Thumb2 R0-R7, R12, R14 for Thumb1	General purpose registers
S0-S31, D0-D31, Q0-Q15	Floating-point registers
сс	The condition flags (N, Z, V, and C)
memory	To be used if the instructions modify any memory. This will avoid keeping memory values cached in registers across the inline assembler statement.

Table 19: List of valid clobbers

THIS IS AN EXAMPLE OF HOW TO USE CLOBBERED MEMORY

```
int StoreExclusive(unsigned long * location, unsigned long value)
{
  int failed;
  asm("strex %0,%2,[%1]"
        : "=&r"(failed)
        : "r"(location), "r"(value)
        : "r"(location), "r"(value)
        : "memory");
  /* Note: 'strex' requires ARMv6 (ARM) or ARMv6T2 (THUMB) */
  return failed;
}
```

Calling assembler routines from C

An assembler routine that will be called from C must:

- Conform to the calling convention
- Have a PUBLIC entry-point label
- Be declared as external before any call, to allow type checking and optional promotion of parameters, as in these examples:

```
extern int foo(void);
or
extern int foo(int i, int j);
```

One way of fulfilling these requirements is to create skeleton code in C, compile it, and study the assembler list file.

CREATING SKELETON CODE

The recommended way to create an assembler language routine with the correct interface is to start with an assembler language source file created by the C compiler. Note that you must create skeleton code for each function prototype.

The following example shows how to create skeleton code to which you can easily add the functional body of the routine. The skeleton source code only needs to declare the variables required and perform simple accesses to them. In this example, the assembler routine takes an int and a char, and then returns an int:

```
extern int gInt;
extern char gChar;
int Func(int arg1, char arg2)
{
    int locInt = arg1;
    gInt = arg1;
    gChar = arg2;
    return locInt;
}
int main()
{
    int locInt = gInt;
    gInt = Func(locInt, gChar);
    return 0;
}
```

Note: In this example we use a low optimization level when compiling the code to show local and global variable access. If a higher level of optimization is used, the required references to local variables could be removed during the optimization. The actual function declaration is not changed by the optimization level.

COMPILING THE CODE



In the IDE, specify list options on file level. Select the file in the workspace window. Then choose **Project>Options**. In the **C/C++ Compiler** category, select **Override inherited settings**. On the **List** page, deselect **Output list file**, and instead select the **Output assembler file** option and its suboption **Include source**. Also, be sure to specify a low level of optimization.

Use these options to compile the skeleton code:

iccarm skeleton.c -lA .

The -1A option creates an assembler language output file including C or C++ source lines as assembler comments. The . (period) specifies that the assembler file should be named in the same way as the C or C++ module (skeleton), but with the filename extension s. Also remember to specify a low level of optimization, and -e for enabling language extensions.

The result is the assembler source output file skeleton.s.

Note: The -1A option creates a list file containing call frame information (CFI) directives, which can be useful if you intend to study these directives and how they are used. If you only want to study the calling convention, you can exclude the CFI directives from the list file.



In the IDE, choose **Project>Options>C/C++ Compiler>List** and deselect the suboption **Include call frame information**.

On the command line, use the option -1B instead of -IA. Note that CFI information must be included in the source code to make the C-SPY Call Stack window work.

The output file

The output file contains the following important information:

- The calling convention
- The return values
- The global variables
- The function parameters
- How to create space on the stack (auto variables)
- Call frame information (CFI).

The CFI directives describe the call frame information needed by the Call Stack window in the debugger. For more information, see *Call frame information*, page 157.

Calling assembler routines from C++

The C calling convention does not apply to C++ functions. Most importantly, a function name is not sufficient to identify a C++ function. The scope and the type of the function are also required to guarantee type-safe linkage, and to resolve overloading.

Another difference is that non-static member functions get an extra, hidden argument, the this pointer.

However, when using C linkage, the calling convention conforms to the C calling convention. An assembler routine can therefore be called from C++ when declared in this manner:

```
extern "C"
{
   int MyRoutine(int);
}
```

The following example shows how to achieve the equivalent to a non-static member function, which means that the implicit this pointer must be made explicit. It is also possible to "wrap" the call to the assembler routine in a member function. Use an inline member function to remove the overhead of the extra call—this assumes that function inlining is enabled:

```
class MyClass;
extern "C"
{
  void DoIt(MyClass *ptr, int arg);
}
class MyClass
{
public:
  inline void DoIt(int arg)
  {
    ::DoIt(this, arg);
  }
};
```

Calling convention

A calling convention is the way a function in a program calls another function. The compiler handles this automatically, but, if a function is written in assembler language, you must know where and how its parameters can be found, how to return to the program location from where it was called, and how to return the resulting value.

It is also important to know which registers an assembler-level routine must preserve. If the program preserves too many registers, the program might be ineffective. If it preserves too few registers, the result would be an incorrect program.

This section describes the calling conventions used by the compiler. These items are examined:

- Function declarations
- C and C++ linkage
- Preserved versus scratch registers
- Function entrance
- Function exit
- Return address handling.

At the end of the section, some examples are shown to describe the calling convention in practice.

The calling convention used by the compiler adheres to the Procedure Call Standard for the ARM architecture, AAPCS, a part of AEABI; see *AEABI compliance*, page 192. AAPCS is not fully described here. For example, the use of floating-point coprocessor registers when using the VFP calling convention is not covered.

FUNCTION DECLARATIONS

In C, a function must be declared in order for the compiler to know how to call it. A declaration could look as follows:

```
int MyFunction(int first, char * second);
```

This means that the function takes two parameters: an integer and a pointer to a character. The function returns a value, an integer.

In the general case, this is the only knowledge that the compiler has about a function. Therefore, it must be able to deduce the calling convention from this information.

USING C LINKAGE IN C++ SOURCE CODE

In C++, a function can have either C or C++ linkage. To call assembler routines from C++, it is easiest if you make the C++ function have C linkage.

This is an example of a declaration of a function with C linkage:

```
extern "C"
{
   int F(int);
}
```

It is often practical to share header files between C and C++. This is an example of a declaration that declares a function with C linkage in both C and C++:

```
#ifdef __cplusplus
extern "C"
{
#endif
int F(int);
#ifdef __cplusplus
}
#endif
```

PRESERVED VERSUS SCRATCH REGISTERS

The general ARM CPU registers are divided into three separate sets, which are described in this section.

Scratch registers

Any function is permitted to destroy the contents of a scratch register. If a function needs the register value after a call to another function, it must store it during the call, for example on the stack.

Any of the registers R0 to R3, and R12, can be used as a scratch register by the function. Note that R12 is a scratch register also when calling between assembler functions only because of automatically inserted instructions for veneers.

Preserved registers

Preserved registers, on the other hand, are preserved across function calls. The called function can use the register for other purposes, but must save the value before using the register and restore it at the exit of the function.

The registers R4 through to R11 are preserved registers. They are preserved by the called function.

Special registers

For some registers, you must consider certain prerequisites:

- The stack pointer register, R13/SP, must at all times point to or below the last element on the stack. In the eventuality of an interrupt, everything below the point the stack pointer points to, can be destroyed. At function entry and exit, the stack pointer must be 8-byte aligned. In the function, the stack pointer must always be word aligned. At exit, SP must have the same value as it had at the entry.
- The register R15/PC is dedicated for the Program Counter.
- The link register, R14/LR, holds the return address at the entrance of the function.

FUNCTION ENTRANCE

Parameters can be passed to a function using one of two basic methods: in registers or on the stack. It is much more efficient to use registers than to take a detour via memory, so the calling convention is designed to use registers as much as possible. Only a limited number of registers can be used for passing parameters; when no more registers are available, the remaining parameters are passed on the stack. These exceptions to the rules apply:

- Interrupt functions cannot take any parameters, except software interrupt functions that accept parameters and have return values
- Software interrupt functions cannot use the stack in the same way as ordinary functions. When an SVC instruction is executed, the processor switches to supervisor mode where the supervisor stack is used. Arguments can therefore not be

passed on the stack if your application is not running in supervisor mode previous to the interrupt.

Hidden parameters

In addition to the parameters visible in a function declaration and definition, there can be hidden parameters:

- If the function returns a structure larger than 32 bits, the memory location where the structure is to be stored is passed as an extra parameter. Notice that it is always treated as the *first parameter*.
- If the function is a non-static C++ member function, then the this pointer is passed as the first parameter (but placed after the return structure pointer, if there is one). For more information, see *Calling assembler routines from C*, page 147.

Register parameters

The registers available for passing parameters are R0-R3:

Parameters	Passed in registers
Scalar and floating-point values no larger than 32 bits,	Passed using the first free register:
and single-precision (32-bits) floating-point values	R0-R3
long long and double-precision (64-bit) values	Passed in first available register pair:
	R0:R1, or R2:R3

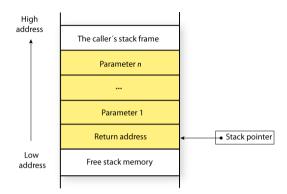
Table 20: Registers used for passing parameters

The assignment of registers to parameters is a straightforward process. Traversing the parameters from left to right, the first parameter is assigned to the available register or registers. Should there be no more available registers, the parameter is passed on the stack in reverse order.

When functions that have parameters smaller than 32 bits are called, the values are sign or zero extended to ensure that the unused bits have consistent values. Whether the values will be sign or zero extended depends on their type—signed or unsigned.

Stack parameters and layout

Stack parameters are stored in memory, starting at the location pointed to by the stack pointer. Below the stack pointer (towards low memory) there is free space that the called function can use. The first stack parameter is stored at the location pointed to by the stack pointer. The next one is stored at the next location on the stack that is divisible by four, etc. It is the responsibility of the caller to clean the stack after the called function has returned.



This figure illustrates how parameters are stored on the stack:

Figure 13: Stack image after the function call

FUNCTION EXIT

A function can return a value to the function or program that called it, or it can have the return type void.

The return value of a function, if any, can be scalar (such as integers and pointers), floating-point, or a structure.

Registers used for returning values

The registers available for returning values are R0 and R0:R1.

Return values	Passed in registers
Scalar and structure return values no larger than 32 bits, and single-precision (32-bit) floating-point return values	RO
The memory address of a structure return value larger than 32 bits	R0
long long and double-precision (64-bit) return values	R0:R1

Table 21: Registers used for returning values

If the returned value is smaller than 32 bits, the value is sign- or zero-extended to 32 bits.

Stack layout at function exit

It is the responsibility of the caller to clean the stack after the called function has returned.

Return address handling

A function written in assembler language should, when finished, return to the caller, by jumping to the address pointed to by the register LR.

At function entry, non-scratch registers and the LR register can be pushed with one instruction. At function exit, all these registers can be popped with one instruction. The return address can be popped directly to PC.

The following example shows what this can look like:

```
name call
section .text:CODE
extern func
push {r4-r6,lr} ; Preserve stack alignment 8
bl func
; Do something here.
pop {r4-r6,pc} ; return
end
```

EXAMPLES

The following section shows a series of declaration examples and the corresponding calling conventions. The complexity of the examples increases toward the end.

Example I

Assume this function declaration:

```
int add1(int);
```

This function takes one parameter in the register R0, and the return value is passed back to its caller in the register R0.

This assembler routine is compatible with the declaration; it will return a value that is one number higher than the value of its parameter:

```
name return
section .text:CODE
adds r0, r0, #1
bx lr
end
```

Example 2

This example shows how structures are passed on the stack. Assume these declarations:

```
struct MyStruct
{
    short a;
    short b;
    short c;
    short d;
    short e;
};
```

int MyFunction(struct MyStruct x, int y);

The values of the structure members a, b, c, and d are passed in registers R0-R3. The last structure member e and the integer parameter y are passed on the stack. The calling function must reserve eight bytes on the top of the stack and copy the contents of the two stack parameters to that location. The return value is passed back to its caller in the register R0.

Example 3

The function below will return a structure of type struct MyStruct.

```
struct MyStruct
{
    int mA;
};
struct MyStruct MyFunction(int x);
```

It is the responsibility of the calling function to allocate a memory location for the return value and pass a pointer to it as a hidden first parameter. The pointer to the location where the return value should be stored is passed in \mathbb{R}^0 . The parameter x is passed in \mathbb{R}^1 .

Assume that the function instead was declared to return a pointer to the structure:

struct MyStruct *MyFunction(int x);

In this case, the return value is a scalar, so there is no hidden parameter. The parameter x is passed in R0, and the return value is returned in R0.

Call frame information

When you debug an application using C-SPY, you can view the *call stack*, that is, the chain of functions that called the current function. To make this possible, the compiler supplies debug information that describes the layout of the call frame, in particular information about where the return address is stored.

If you want the call stack to be available when debugging a routine written in assembler language, you must supply equivalent debug information in your assembler source using the assembler directive CFI. This directive is described in detail in the *IAR Assembler Reference Guide for ARM*.

CFI DIRECTIVES

The CFI directives provide C-SPY with information about the state of the calling function(s). Most important of this is the return address, and the value of the stack pointer at the entry of the function or assembler routine. Given this information, C-SPY can reconstruct the state for the calling function, and thereby unwind the stack.

A full description about the calling convention might require extensive call frame information. In many cases, a more limited approach will suffice.

When describing the call frame information, the following three components must be present:

- A names block describing the available resources to be tracked
- A *common block* corresponding to the calling convention
- A *data block* describing the changes that are performed on the call frame. This typically includes information about when the stack pointer is changed, and when permanent registers are stored or restored on the stack.

This table lists all the resources defined in the names block used by the compiler:

Resource	Description	
CFA R13	The call frames of the stack	
R0-R12	Processor general-purpose 32-bit registers	
R13	Stack pointer, SP	
R14	Link register, LR	
D0-D31	Vector Floating Point (VFP) 64-bit coprocessor register	
CPSR	Current program status register	
SPSR	Saved program status register	

Table 22: Call frame information resources defined in a names block

CREATING ASSEMBLER SOURCE WITH CFI SUPPORT

The recommended way to create an assembler language routine that handles call frame information correctly is to start with an assembler language source file created by the compiler.

Start with suitable C source code, for example:

```
int F(int);
int cfiExample(int i)
{
  return i + F(i);
}
```

2 Compile the C source code, and make sure to create a list file that contains call frame information—the CFI directives.

On the command line, use the option -IA.



In the IDE, choose **Project>Options>C/C++ Compiler>List** and make sure the suboption **Include call frame information** is selected.

For the source code in this example, the list file looks like this:

```
NAME Cfi
        RTMODEL "___SystemLibrary", "DLib"
        EXTERN F
        PUBLIC cfiExample
        CFI Names cfiNames0
        CFI StackFrame CFA R13 DATA
        CFI Resource R0:32, R1:32, R2:32, R3:32, R4:32, R5:32,
R6:32, R7:32
       CFI Resource R8:32, R9:32, R10:32, R11:32, R12:32,
R13:32, R14:32
        CFI EndNames cfiNames0
        CFI Common cfiCommon0 Using cfiNames0
        CFI CodeAlign 4
        CFI DataAlign 4
        CFI ReturnAddress R14 CODE
        CFT CFA R13+0
        CFI R0 Undefined
        CFI R1 Undefined
        CFI R2 Undefined
        CFI R3 Undefined
        CFI R4 SameValue
        CFI R5 SameValue
        CFI R6 SameValue
        CFI R7 SameValue
        CFI R8 SameValue
        CFI R9 SameValue
        CFI R10 SameValue
        CFI R11 SameValue
        CFI R12 Undefined
        CFI R14 SameValue
        CFI EndCommon cfiCommon0
        SECTION `.text`:CODE:NOROOT(2)
        CFI Block cfiBlock0 Using cfiCommon0
        CFI Function cfiExample
        ARM
cfiExample:
        PUSH
                 \{R4, LR\}
        CFI R14 Frame(CFA, -4)
        CFI R4 Frame(CFA, -8)
```

```
CFI CFA R13+8
MOVS R4,R0
MOVS R0,R4
BL F
ADDS R0,R0,R4
POP {R4,PC} ;; return
CFI EndBlock cfiBlock0
```

END

Note: The header file Common.i contains the macros CFI_NAMES_BLOCK, CFI_COMMON_ARM, and CFI_COMMON_Thumb, which declare a typical names block and a typical common block. These two macros declare several resources, both concrete and virtual.

Using C

This chapter gives an overview of the compiler's support for the C language. The chapter also gives a brief overview of the IAR C language extensions.

C language overview

The IAR C/C++ Compiler for ARM supports the ISO/IEC 9899:1999 standard (including up to technical corrigendum No.3), also known as C99. In this guide, this standard is referred to as *Standard C* and is the default standard used in the compiler. This standard is stricter than C89.

In addition, the compiler also supports the ISO 9899:1990 standard (including all technical corrigenda and addenda), also known as C94, C90, C89, and ANSI C. In this guide, this standard is referred to as *C89*. Use the --c89 compiler option to enable this standard.

The C99 standard is derived from C89, but adds features like these:

- The inline keyword advises the compiler that the function defined immediately after the keyword should be inlined
- Declarations and statements can be mixed within the same scope
- A declaration in the initialization expression of a for loop
- The bool data type
- The long long data type
- The complex floating-point type
- C++ style comments
- Compound literals
- · Incomplete arrays at the end of structs
- Hexadecimal floating-point constants
- Designated initializers in structures and arrays
- The preprocessor operator _Pragma()
- Variadic macros, which are the preprocessor macro equivalents of printf style functions
- VLA (variable length arrays) must be explicitly enabled with the compiler option --vla
- Inline assembler using the asm or the __asm keyword, see *Inline assembler*, page 139.

Note: Even though it is a C99 feature, the IAR C/C++ Compiler for ARM does not support UCNs (universal character names).

Extensions overview

The compiler offers the features of Standard C and a wide set of extensions, ranging from features specifically tailored for efficient programming in the embedded industry to the relaxation of some minor standards issues.

This is an overview of the available extensions:

IAR C language extensions

For information about available language extensions, see *IAR C language extensions*, page 163. For more information about the extended keywords, see the chapter *Extended keywords*. For information about C++, the two levels of support for the language, and C++ language extensions; see the chapter *Using* C++.

Pragma directives

The #pragma directive is defined by Standard C and is a mechanism for using vendor-specific extensions in a controlled way to make sure that the source code is still portable.

The compiler provides a set of predefined pragma directives, which can be used for controlling the behavior of the compiler, for example how it allocates memory, whether it allows extended keywords, and whether it outputs warning messages. Most pragma directives are preprocessed, which means that macros are substituted in a pragma directive. The pragma directives are always enabled in the compiler. For several of them there is also a corresponding C/C++ language extension. For information about available pragma directives, see the chapter *Pragma directives*.

Preprocessor extensions

The preprocessor of the compiler adheres to Standard C. The compiler also makes several preprocessor-related extensions available to you. For more information, see the chapter *The preprocessor*.

Intrinsic functions

The intrinsic functions provide direct access to low-level processor operations and can be very useful in, for example, time-critical routines. The intrinsic functions compile into inline code, either as a single instruction or as a short sequence of instructions. For more information about using intrinsic functions, see *Mixing C and assembler*, page 137. For information about available functions, see the chapter *Intrinsic functions*.

• Library functions

The IAR DLIB Library provides the C and C++ library definitions that apply to embedded systems. For more information, see *IAR DLIB Library*, page 391.

Note: Any use of these extensions, except for the pragma directives, makes your source code inconsistent with Standard C.

ENABLING LANGUAGE EXTENSIONS

You can choose different levels of language conformance by means of project options:

Command line	IDE*	Description
strict	Strict	All IAR C language extensions are disabled; errors are issued for anything that is not part of Standard C.
None	Standard	All extensions to Standard C are enabled, but no extensions for embedded systems programming. For information about extensions, see IAR C language extensions, page 163.
-e	Standard with IAR extensions	All IAR C language extensions are enabled.

Table 23: Language extensions

* In the IDE, choose **Project>Options>** C/C++ Compiler>Language>Language conformance and select the appropriate option. Note that language extensions are enabled by default.

IAR C language extensions

The compiler provides a wide set of C language extensions. To help you to find the extensions required by your application, they are grouped like this in this section:

- *Extensions for embedded systems programming*—extensions specifically tailored for efficient embedded programming for the specific core you are using, typically to meet memory restrictions
- *Relaxations to Standard C*—that is, the relaxation of some minor Standard C issues and also some useful but minor syntax extensions, see *Relaxations to Standard C*, page 166.

EXTENSIONS FOR EMBEDDED SYSTEMS PROGRAMMING

The following language extensions are available both in the C and the C++ programming languages and they are well suited for embedded systems programming:

• Type attributes, and object attributes

For information about the related concepts, the general syntax rules, and for reference information, see the chapter *Extended keywords*.

• Placement at an absolute address or in a named section

The @ operator or the directive #pragma location can be used for placing global and static variables at absolute addresses, or placing a variable or function in a named section. For more information about using these features, see *Controlling data and function placement in memory*, page 201, and *location*, page 333.

• Alignment control

Each data type has its own alignment; for more information, see *Alignment*, page 297. If you want to change the alignment, the __packed data type attribute, and the #pragma pack, and #pragma data_alignment directives are available. If you want to check the alignment of an object, use the __ALIGNOF__() operator.

The __ALIGNOF__ operator is used for accessing the alignment of an object. It takes one of two forms:

- __ALIGNOF__ (type)
- __ALIGNOF__ (*expression*)

In the second form, the expression is not evaluated.

• Anonymous structs and unions

C++ includes a feature called anonymous unions. The compiler allows a similar feature for both structs and unions in the C programming language. For more information, see *Anonymous structs and unions*, page 199.

• Bitfields and non-standard types

In Standard C, a bitfield must be of the type int or unsigned int. Using IAR C language extensions, any integer type or enumeration can be used. The advantage is that the struct will sometimes be smaller. For more information, see *Bitfields*, page 300.

static_assert()

The construction static_assert(const-expression, "message"); can be used in C/C++. The construction will be evaluated at compile time and if const-expression is false, a message will be issued including the message string.

• Parameters in variadic macros

Variadic macros are the preprocessor macro equivalents of printf style functions. The preprocessor accepts variadic macros with no arguments, which means if no parameter matches the . . . parameter, the comma is then deleted in the ", ##__VA_ARGS__" macro definition. According to Standard C, the . . . parameter must be matched with at least one argument.

Dedicated section operators

The compiler supports getting the start address, end address, and size for a section with these built-in section operators:

section_begin	Returns the address of the first byte of the named section or block.
section_end	Returns the address of the first byte <i>after</i> the named section or block.
section_size	Returns the size of the named section or block in bytes.
Note: The aliasesse segment_size/s	gment_begin/sfb,segment_end/sfe, and fs can also be used.

The operators can be used on named sections or on named blocks defined in the linker configuration file.

These operators behave syntactically as if declared like:

```
void * __section_begin(char const * section)
void * __section_end(char const * section)
size_t * __section_size(char const * section)
```

When you use the @ operator or the #pragma location directive to place a data object or a function in a user-defined section, or when you use named blocks in the linker configuration file, the section operators can be used for getting the start and end address of the memory range where the sections or blocks were placed.

The named *section* must be a string literal and it must have been declared earlier with the #pragma *section* directive. The type of the __section_begin operator is a pointer to void. Note that you must enable language extensions to use these operators.

The operators are implemented in terms of *symbols* with dedicated names, and will appear in the linker map file under these names:

Operator	Symbol
section_begin(sec)	sec\$\$Base
section_end(sec)	sec\$\$Limit
section_size(sec)	sec\$\$Length

Table 24: Section operators and their symbols

Note that the linker will not necessarily place sections with the same name consecutively when these operators are not used. Using one of these operators (or the equivalent symbols) will cause the linker to behave as if the sections were in a named block. This is to assure that the sections are placed consecutively, so that the operators can be assigned meaningful values. If this is in conflict with the section placement as specified in the linker configuration file, the linker will issue an error.

Example

In this example, the type of the __section_begin operator is void *.

```
#pragma section="MYSECTION"
...
section_start_address = __section_begin("MYSECTION");
```

See also section, page 340, and location, page 333.

RELAXATIONS TO STANDARD C

This section lists and briefly describes the relaxation of some Standard C issues and also some useful but minor syntax extensions:

Arrays of incomplete types

An array can have an incomplete struct, union, or enum type as its element type. The types must be completed before the array is used (if it is), or by the end of the compilation unit (if it is not).

• Forward declaration of enum types

The extensions allow you to first declare the name of an enum and later resolve it by specifying the brace-enclosed list.

• Accepting missing semicolon at the end of a struct or union specifier

A warning—instead of an error—is issued if the semicolon at the end of a struct or union specifier is missing.

• Null and void

In operations on pointers, a pointer to void is always implicitly converted to another type if necessary, and a null pointer constant is always implicitly converted to a null

pointer of the right type if necessary. In Standard C, some operators allow this kind of behavior, while others do not allow it.

• Casting pointers to integers in static initializers

In an initializer, a pointer constant value can be cast to an integral type if the integral type is large enough to contain it. For more information about casting pointers, see *Casting*, page 306.

• Taking the address of a register variable

In Standard C, it is illegal to take the address of a variable specified as a register variable. The compiler allows this, but a warning is issued.

• long float means double

The type long float is accepted as a synonym for double.

• Repeated typedef declarations

Redeclarations of typedef that occur in the same scope are allowed, but a warning is issued.

Mixing pointer types

Assignment and pointer difference is allowed between pointers to types that are interchangeable but not identical; for example, unsigned char * and char *. This includes pointers to integral types of the same size. A warning is issued.

Assignment of a string constant to a pointer to any kind of character is allowed, and no warning is issued.

• Non-top level const

Assignment of pointers is allowed in cases where the destination type has added type qualifiers that are not at the top level (for example, int ** to int const **). Comparing and taking the difference of such pointers is also allowed.

• Non-lvalue arrays

A non-lvalue array expression is converted to a pointer to the first element of the array when it is used.

• Comments at the end of preprocessor directives

This extension, which makes it legal to place text after preprocessor directives, is enabled unless the strict Standard C mode is used. The purpose of this language extension is to support compilation of legacy code; we do not recommend that you write new code in this fashion.

• An extra comma at the end of enum lists

Placing an extra comma is allowed at the end of an enum list. In strict Standard C mode, a warning is issued.

• A label preceding a }

In Standard C, a label must be followed by at least one statement. Therefore, it is illegal to place the label at the end of a block. The compiler allows this, but issues a warning.

Note that this also applies to the labels of switch statements.

Empty declarations

An empty declaration (a semicolon by itself) is allowed, but a remark is issued (provided that remarks are enabled).

Single-value initialization

Standard C requires that all initializer expressions of static arrays, structs, and unions are enclosed in braces.

Single-value initializers are allowed to appear without braces, but a warning is issued. The compiler accepts this expression:

struct str { int a; x = 10;

{

}

Declarations in other scopes

External and static declarations in other scopes are visible. In the following example, the variable y can be used at the end of the function, even though it should only be visible in the body of the if statement. A warning is issued.

```
int test(int x)
  if (x)
  {
    extern int y;
    y = 1;
  }
  return y;
```

Expanding function names into strings with the function as context

Use any of the symbols __func__ or __FUNCTION__ inside a function body to make the symbol expand into a string that contains the name of the current function. Use the symbol __PRETTY_FUNCTION__ to also include the parameter types and return type. The result might, for example, look like this if you use the ___PRETTY_FUNCTION__ symbol:

```
"void func(char)"
```

These symbols are useful for assertions and other trace utilities and they require that language extensions are enabled, see -e, page 242.

• Static functions in function and block scopes

Static functions may be declared in function and block scopes. Their declarations are moved to the file scope.

• Numbers scanned according to the syntax for numbers

Numbers are scanned according to the syntax for numbers rather than the pp-number syntax. Thus, 0x123e+1 is scanned as three tokens instead of one valid token. (If the --strict option is used, the pp-number syntax is used instead.)

IAR C language extensions

Using C++

IAR Systems supports the C++ language. You can choose between these standards: Standard C++, the industry-standard Embedded C++, and Extended Embedded C++. This chapter describes what you need to consider when using the C++ language.

Overview—Embedded C++ and Extended Embedded C++

Embedded C++ is a proper subset of the ISO/IEC C++ standard which is intended for embedded systems programming. It was defined by an industry consortium, the Embedded C++ Technical Committee. Performance and portability are particularly important in embedded systems development, which was considered when defining the language. EC++ offers the same object-oriented benefits as C++, but without some features that can increase code size and execution time in ways that are hard to predict.

EMBEDDED C++

These C++ features are supported:

- Classes, which are user-defined types that incorporate both data structure and behavior; the essential feature of inheritance allows data structure and behavior to be shared among classes
- Polymorphism, which means that an operation can behave differently on different classes, is provided by virtual functions
- Overloading of operators and function names, which allows several operators or functions with the same name, provided that their argument lists are sufficiently different
- Type-safe memory management using the operators new and delete
- Inline functions, which are indicated as particularly suitable for inline expansion.

C++ features that are excluded are those that introduce overhead in execution time or code size that are beyond the control of the programmer. Also excluded are features added very late before Standard C++ was defined. Embedded C++ thus offers a subset of C++ which is efficient and fully supported by existing development tools.

Embedded C++ lacks these features of C++:

- Templates
- Multiple and virtual inheritance
- Exception handling

- Runtime type information
- New cast syntax (the operators dynamic_cast, static_cast, reinterpret_cast, and const_cast)
- Namespaces
- The mutable attribute.

The exclusion of these language features makes the runtime library significantly more efficient. The Embedded C++ library furthermore differs from the full C++ library in that:

- The standard template library (STL) is excluded
- Streams, strings, and complex numbers are supported without the use of templates
- Library features which relate to exception handling and runtime type information (the headers except, stdexcept, and typeinfo) are excluded.

Note: The library is not in the std namespace, because Embedded C++ does not support namespaces.

EXTENDED EMBEDDED C++

IAR Systems' Extended EC++ is a slightly larger subset of C++ which adds these features to the standard EC++:

- Full template support
- Multiple and virtual inheritance
- Namespace support
- The mutable attribute
- The cast operators static_cast, const_cast, and reinterpret_cast.

All these added features conform to the C++ standard.

To support Extended EC++, this product includes a version of the standard template library (STL), in other words, the C++ standard chapters utilities, containers, iterators, algorithms, and some numerics. This STL is tailored for use with the Extended EC++ language, which means no exceptions, no multiple inheritance, and no support for runtime type information (rtti). Moreover, the library is not in the std namespace.

Note: A module compiled with Extended EC++ enabled is fully link-compatible with a module compiled without Extended EC++ enabled.

Overview—Standard C++

The IAR C++ implementation fully complies with the ISO/IEC 1488:2003 C++ standard. In this guide, this standard is referred to as Standard C++.

The main reason for using Standard C++ instead of EC++ or EEC++ is when there is a need for either:

- Exception support
- Runtime type information (RTTI) support
- The standard C++ library (the EC++ library is a stripped version of the C++ library where streams and strings are not templates).

If code size is important and your application does not need any of these features, EC++ (or EEC++) is the better choice.

MODES FOR EXCEPTIONS AND RTTI SUPPORT

Both exceptions and runtime type information result in increased code size simply by being included in your application. You might want to disable either or both of these features to avoid this increase:

- Support for runtime type information constructs can be disabled by using the compiler option --no_rtti
- Support for exceptions can be disabled by using the compiler option --no_exceptions

Even if support is enabled while compiling, the linker can avoid including the extra code and tables in the final application. If no part of your application actually throws an exception, the code and tables supporting the use of exceptions are not included in the application code image. Also, if dynamic runtime type information constructs (dynamic_cast/typeid) are not used with polymorphic types, the objects needed to support them are not included in the application code image. To control this behavior, use the linker options --no_exceptions, --force_exceptions, and --no_dynamic_rtti_elimination.

Disabling exception support

When you use the compiler option --no_exceptions, the following will generate a compiler error:

- throw expressions
- try-catch statements
- Exception specifications on function definitions.

In addition, the extra code and tables needed to handle destruction of objects with auto storage duration when an exception is propagated through a function will not be generated when the compiler option $--no_exceptions$ is used.

All functionality in system headers not directly involving exceptions is supported when the compiler option --no_exceptions is used.

The linker will produce an error if you try to link C++ modules compiled with exception support with modules compiled without exception support

For more information, see --no_exceptions, page 252.

Disabling RTTI support

When you use the compiler option --no_rtti, the following will generate a compiler error:

- The typeid operator
- The dynamic_cast operator.

Note: If --no_rtti is used but exception support is enabled, most RTTI support is still included in the compiler output object file because it is needed for exceptions to work.

For more information, see --no_rtti, page 254.

EXCEPTION HANDLING

Exception handling can be divided into three parts:

- Exception raise mechanisms—in C++ they are the throw and rethrow expressions.
- Exception catch mechanisms—in C++ they are the try-catch statements, the exception specifications for a function, and the implicit catch to prevent an exception leaking out from main.
- Information about currently active functions—if they have try-catch statements and the set of auto objects whose destructors need to be run if an exception is propagated through the function.

When an exception is raised, the function call stack is unwound, function by function, block by block. For each function or block, the destructors of auto objects that need destruction are run, and a check is made whether there is a catch handler for the exception. If there is, the execution will continue from that catch handler.

An application that mixes C++ code with assembler and C code, and that throws exceptions from one C++ function to another via assembler routines and C functions must use the linker option --exception_tables with the argument unwind.

The implementation of exceptions

Exceptions are implemented using a table method. For each function, the tables describe:

• How to unwind the function, that is, how to find its caller on the stack and restore registers that need restoring

- Which catch handlers that exist in the function
- Whether the function has an exception specification and which exceptions it allows to propagate
- The set of auto objects whose destructors must be run.

When an exception is raised, the runtime will proceed in two phases. The first phase will use the exception tables to search the stack for a function invocation containing a catch handler or exception specification that would cause stack unwinding to halt at that point. Once this point is found, the second phase is entered, doing the actual unwinding, and running the destructors of auto objects where that is needed.

The table method results in virtually no overhead in execution time or RAM usage when an exception is not actually thrown. It does incur a significant penalty in read-only memory usage for the tables and the extra code, and throwing and catching an exception is a relatively expensive operation.

The destruction of auto objects when the stack is being unwound as a result of an exception is implemented in code separated from the code that handles the normal operation of a function. This code, together with the code in catch handlers, is placed in a separate section (.exc.text) from the normal code (normally placed in .text). In some cases, for instance when there is fast and slow ROM memory, it can be advantageous to select on this difference when placing sections in the linker configuration file.

Enabling support for C++ and variants

In the compiler, the default language is C.

To compile files written in Standard C++, you must use the --c++ compiler option. See --c++, page 237.

To compile files written in Embedded C++, you must use the --ec++ compiler option. See --ec++, page 243.

To take advantage of *Extended* Embedded C++ features in your source code, you must use the --eec++ compiler option. See --eec++, page 243.



To enable EC++, or EEC++, or C++ in the IDE, choose **Project>Options>C/C++ Compiler>Language** and select the appropriate standard.

C++ and EC++ feature descriptions

When you write C++ source code for the IAR C/C++ Compiler for ARM, you must be aware of some benefits and some possible quirks when mixing C++ features—such as classes, and class members—with IAR language extensions, such as IAR-specific attributes.

USING IAR ATTRIBUTES WITH CLASSES

Static data members of C++ classes are treated the same way global variables are, and can have any applicable IAR type and object attribute.

Member functions are in general treated the same way free functions are, and can have any applicable IAR type and object attributes. Virtual member functions can only have attributes that are compatible with default function pointers, and constructors and destructors cannot have any such attributes.

The location operator @ and the #pragma location directive can be used on static data members and with all member functions.

Example

```
class MyClass
{
  public:
    // Locate a static variable in __memattr memory at address 60
    static __no_init int mI @ 60;
    // A static Thumb function
    static __thumb void F();
    // A Thumb function
    _thumb void G();
    // Interworking assumed
    virtual __arm void ArmH();
    // Locate a virtual function into SPECIAL
    virtual void M() const volatile @ "SPECIAL";
};
```

FUNCTION TYPES

A function type with extern "C" linkage is compatible with a function that has C++ linkage.

Example

```
extern "C"
{
 typedef void (*FpC)(void); // A C function typedef
3
typedef void (*FpCpp)(void);
                                // A C++ function typedef
FpC F1;
FpCpp F2;
void MyF(FpC);
void MyG()
{
 MyF(F1);
                                // Always works
 M_{VF}(F2);
                                // FpCpp is compatible with FpC
3
```

USING STATIC CLASS OBJECTS IN INTERRUPTS

If interrupt functions use static class objects that need to be constructed (using constructors) or destroyed (using destructors), your application will not work properly if the interrupt occurs before the objects are constructed, or, during or after the objects are destroyed.

To avoid this, make sure that these interrupts are not enabled until the static objects have been constructed, and are disabled when returning from main or calling exit. For information about system startup, see *System startup and termination*, page 113.

Function local static class objects are constructed the first time execution passes through their declaration, and are destroyed when returning from main or when calling exit.

USING NEW HANDLERS

To handle memory exhaustion, you can use the set_new_handler function.

New handlers in Standard C++

If you do not call set_new_handler, or call it with a NULL new handler, and operator new fails to allocate enough memory, operator new will throw std::bad_alloc if exceptions are enabled. If exceptions are not enabled, operator new will instead call abort.

If you call set_new_handler with a non-NULL new handler, the provided new handler will be called by operator new if the operator new fails to allocate enough memory. The new handler must then make more memory available and return, or abort execution in some manner. If exceptions are enabled, the new handler can also throw a std::bad_alloc exception. The nothrow variant of operator new will only return NULL in the presence of a new handler if exceptions are enabled and the new handler throws std::bad_alloc.

New handlers in Embedded C++

If you do not call set_new_handler, or call it with a NULL new handler, and operator new fails to allocate enough memory, it will call abort. The nothrow variant of the new operator will instead return NULL.

If you call set_new_handler with a non-NULL new handler, the provided new handler will be called by operator new if operator new fails to allocate memory. The new handler must then make more memory available and return, or abort execution in some manner. The nothrow variant of operator new will never return NULL in the presence of a new handler.

TEMPLATES

C++ and Extended EC++ supports templates according to the C++ standard, but not the export keyword. The implementation uses a two-phase lookup which means that the keyword typename must be inserted wherever needed. Furthermore, at each use of a template, the definitions of all possible templates must be visible. This means that the definitions of all templates must be in include files or in the actual source file.

DEBUG SUPPORT IN C-SPY

C-SPY® has debug support for some features in C++.

- C-SPY® has built-in display support for the STL containers. The logical structure of containers is presented in the watch views in a comprehensive way that is easy to understand and follow.
- You can make C-SPY stop at a throw statement or if a raised exception does not have any corresponding catch statement.

For more information about this, see the C-SPY® Debugging Guide for ARM.

EEC++ feature description

This section describes features that distinguish Extended EC++ from EC++.

TEMPLATES

The compiler supports templates with the syntax and semantics as defined by Standard C++. However, note that the STL (standard template library) delivered with the product is tailored for Extended EC++, see *Extended Embedded* C++, page 172.

VARIANTS OF CAST OPERATORS

In Extended EC++ these additional variants of C++ cast operators can be used:

const_cast<to>(from)
static_cast<to>(from)
reinterpret_cast<to>(from)

MUTABLE

The mutable attribute is supported in Extended EC++. A mutable symbol can be changed even though the whole class object is const.

NAMESPACE

The namespace feature is only supported in *Extended* EC++. This means that you can use namespaces to partition your code. Note, however, that the library itself is not placed in the std namespace.

THE STD NAMESPACE

The std namespace is not used in either standard EC++ or in Extended EC++. If you have code that refers to symbols in the std namespace, simply define std as nothing; for example:

#define std

You must make sure that identifiers in your application do not interfere with identifiers in the runtime library.

EC++ and C++ language extensions

When you use the compiler in any C++ mode and enable IAR language extensions, the following C++ language extensions are available in the compiler:

• In a friend declaration of a class, the class keyword can be omitted, for example:

• Constants of a scalar type can be defined within classes, for example:

According to the standard, initialized static data members should be used instead.

• In the declaration of a class member, a qualified name can be used, for example:

```
struct A
{
    int A::F(); // Possible when using IAR language extensions
    int G(); // According to the standard
};
```

• It is permitted to use an implicit type conversion between a pointer to a function with C linkage (extern "C") and a pointer to a function with C++ linkage (extern "C++"), for example:

```
extern "C" void F(); // Function with C linkage
void (*PF)() // PF points to a function with C++ linkage
= &F; // Implicit conversion of function pointer.
```

According to the standard, the pointer must be explicitly converted.

• If the second or third operands in a construction that contains the ? operator are string literals or wide string literals (which in C++ are constants), the operands can be implicitly converted to char * or wchar_t *, for example:

```
bool X;
```

- Default arguments can be specified for function parameters not only in the top-level function declaration, which is according to the standard, but also in typedef declarations, in pointer-to-function function declarations, and in pointer-to-member function declarations.
- In a function that contains a non-static local variable and a class that contains a non-evaluated expression (for example a sizeof expression), the expression can reference the non-static local variable. However, a warning is issued.
- An anonymous union can be introduced into a containing class by a typedef name. It is not necessary to first declare the union. For example:

```
typedef union
{
    int i,j;
} U; // U identifies a reusable anonymous union.
class A
{
public:
    U; // OK -- references to A::i and A::j are allowed.
};
```

In addition, this extension also permits *anonymous classes* and *anonymous structs*, as long as they have no C++ features (for example, no static data members or member functions, and no non-public members) and have no nested types other than other anonymous classes, structs, or unions. For example:

```
struct A
{
   struct
   {
    int i,j;
   }; // OK -- references to A::i and A::j are allowed.
};
```

• The friend class syntax allows nonclass types as well as class types expressed through a typedef without an elaborated type name. For example:

```
typedef struct S ST;
```

Note: If you use any of these constructions without first enabling language extensions, errors are issued.

Application-related considerations

This chapter discusses a selected range of application issues related to developing your embedded application.

Typically, this chapter highlights issues that are not specifically related to only the compiler or the linker.

Output format considerations

The linker produces an absolute executable image in the ELF/DWARF object file format.

You can use the IAR ELF Tool—ielftool— to convert an absolute ELF image to a format more suitable for loading directly to memory, or burning to a PROM or flash memory etc.

ielftool can produce these output formats:

- Plain binary
- Motorola S-records
- Intel hex.

Note: ielftool can also be used for other types of transformations, such as filling and calculating checksums in the absolute image.

The source code for ielftool is provided in the arm/src directory. For more information about ielftool, see *The IAR ELF Tool—ielftool*, page 440.

Stack considerations

The stack is used by functions to store variables and other information that is used locally by functions, as described in the chapter *Data storage*. It is a continuous block of memory pointed to by the processor stack pointer register SP.

The block used for holding the stack is called CSTACK. The system startup code initializes the stack pointer to the end of the stack.

STACK SIZE CONSIDERATIONS

The compiler uses the internal data stack for a variety of user application operations, and the required stack size depends heavily on the details of these operations. If the given stack size is too large, RAM will be wasted. If the given stack size is too small, two things can happen, depending on where in memory you located your stack. Both alternatives are likely to result in application failure. Either variable storage will be overwritten, leading to undefined behavior, or the stack will fall outside of the memory area, leading to an abnormal termination of your application. Because the second alternative is easier to detect, you should consider placing your stack so that it grows toward the end of the memory.

For more information about the stack size, see *Setting up the stack*, page 89, and *Saving stack space and RAM memory*, page 211.

STACK ALIGNMENT

The default cstartup code automatically initializes all stacks to an 8-byte aligned address.

For more information about aligning the stack, see *Calling convention*, page 150 and more specifically *Special registers*, page 152 and *Stack parameters and layout*, page 153.

EXCEPTION STACK

The ARM architecture supports five exception modes which are entered when different exceptions occur. Each exception mode has its own stack to avoid corrupting the System/User mode stack.

The table shows proposed stack names for the various exception stacks, but any name can be used:

Processor mode	Proposed stack section name	Description
Supervisor	SVC_STACK	Operation system stack.
IRQ	IRQ_STACK	Stack for general-purpose (IRQ) interrupt handlers.
FIQ	FIQ_STACK	Stack for high-speed (FIQ) interrupt handlers.
Undefined	UND_STACK	Stack for undefined instruction interrupts. Supports software emulation of hardware coprocessors and instruction set extensions.

Table 25: Exception stacks

Processor mode	Proposed stack section name	Description
Abort	ABT_STACK	Stack for instruction fetch and data
		access memory abort interrupt
		handlers.

Table 25: Exception stacks

For each processor mode where a stack is needed, a separate stack pointer must be initialized in your startup code, and section placement should be done in the linker configuration file. The IRQ and FIQ stacks are the only exception stacks which are preconfigured in the supplied cstartup.s and lnkarm.icf files, but other exception stacks can easily be added.

Cortex-M does not have individual exception stacks. By default, all exception stacks are placed in the CSTACK section.



To view any of these stacks in the Stack window available in the IDE, these preconfigured section names must be used instead of user-defined section names.

Heap considerations

The heap contains dynamic data allocated by use of the C function malloc (or one of its relatives) or the C++ operator new.

If your application uses dynamic memory allocation, you should be familiar with:

- · Linker sections used for the heap
- Allocating the heap size, see Setting up the heap, page 90.

The memory allocated to the heap is placed in the section HEAP, which is only included in the application if dynamic memory allocation is actually used.

HEAP SIZE AND STANDARD I/O



If you excluded FILE descriptors from the DLIB runtime environment, as in the normal configuration, there are no input and output buffers at all. Otherwise, as in the full configuration, be aware that the size of the input and output buffers is set to 512 bytes in the stdio library header file. If the heap is too small, I/O will not be buffered, which is considerably slower than when I/O is buffered. If you execute the application using the simulator driver of the IAR C-SPY® Debugger, you are not likely to notice the speed penalty, but it is quite noticeable when the application runs on an ARM core. If you use the standard I/O library, you should set the heap size to a value which accommodates the needs of the standard I/O buffer.

Interaction between the tools and your application

The linking process and the application can interact symbolically in four ways:

- Creating a symbol by using the ILINK command line option --define_symbol. ILINK will create a public absolute constant symbol that the application can use as a label, as a size, as setup for a debugger, etc.
- Creating an exported configuration symbol by using the command line option --config_def or the configuration directive define symbol, and exporting the symbol using the export symbol directive. ILINK will create a public absolute constant symbol that the application can use as a label, as a size, as setup for a debugger, etc.

One advantage of this symbol definition is that this symbol can also be used in expressions in the configuration file, for example to control the placement of sections into memory ranges.

- Using the compiler operators __section_begin, __section_end, or __section_size, or the assembler operators SFB, SFE, or SIZEOF on a named section or block. These operators provide access to the start address, end address, and size of a contiguous sequence of sections with the same name, or of a linker block specified in the linker configuration file.
- The command line option --entry informs ILINK about the start label of the application. It is used by ILINK as a root symbol and to inform the debugger where to start execution.

The following lines illustrate how to use these mechanisms. Add these options to your command line:

```
--define_symbol NrOfElements=10
--config_def HeapSize=1024
```

The linker configuration file can look like this:

```
define memory Mem with size = 4G;
define region ROM = Mem:[from 0x00000 size 0x10000];
define region RAM = Mem:[from 0x20000 size 0x10000];
/* Export of symbol */
export symbol HeapSize;
/* Setup a heap area witha size defined by an ILINK option */
define block MyHEAP with size = HeapSize, alignment = 8 {};
```

place in RAM { block MyHEAP };

Add these lines to your application source code:

```
#include <stdlib.h>
/* Use symbol defined by ILINK option to dynamically allocate an
array of elements with specified size. The value takes the form
of a label.
*/
extern int NrOfElements;
typedef char Elements;
Elements *GetElementArray()
{
 return malloc(sizeof(Elements) * (long) &NrOfElements);
}
/* Use a symbol defined by ILINK option, a symbol that in the
 * configuration file was made available to the application.
 * /
extern char HeapSize;
/* Declare the section that contains the heap. */
#pragma section = "MYHEAP"
char *MyHeap()
{
  /* First get start of statically allocated section, */
 char *p = __section_begin("MYHEAP");
  /* ...then we zero it, using the imported size. */
  for (int i = 0; i < (int) &HeapSize; ++i)
  {
   p[i] = 0;
 }
 return p;
}
```

Checksum calculation

The IAR ELF Tool—ielftool—fills specific ranges of memory with a pattern and then calculates a checksum for those ranges. The calculated checksum replaces the value of an existing symbol in the input ELF image. The application can then verify that the ranges did not change.

To use checksumming to verify the integrity of your application, you must:

- Reserve a place, with an associated name and size, for the checksum calculated by ielftool
- Choose a checksum algorithm, set up ielftool for it, and include source code for the algorithm in your application
- Decide what memory ranges to verify and set up both ielftool and the source code for it in your application source code.



To set up ielftool in the IDE, choose Project>Options>Linker>Checksum.

CALCULATING A CHECKSUM

In this example, a checksum is calculated for ROM memory at 0x8002 up to 0x8FFF and the 2-byte calculated checksum is placed at 0x8000.

Creating a place for the calculated checksum

You can create a place for the calculated checksum in two ways; by creating a global C/C++ or assembler constant symbol with a proper size, residing in a specific section (in this example .checksum), or by using the linker option --place_holder.

For example, to create a 2-byte space for the symbol __checksum in the section .checksum, with alignment 4:

--place_holder __checksum,2,.checksum,4

Note: The .checksum section will only be included in your application if the section appears to be needed. If the checksum is not needed by the application itself, you can use the linker option --keep=__checksum or the linker directive keep to force the section to be included.

To place the .checksum section, you must modify the linker configuration file. It can look like this (note the handling of the block CHECKSUM):

```
define memory Mem with size = 4G;
define region ROM_region = Mem: [from 0x8000 to 0x80000000 - 1];
define region RAM_region = Mem: [from 0x80000000 to 0x100000000 -2
1;
initialize by copy { rw };
do not initialize { section .noinit };
define block HEAP
                    with alignment = 8, size = 16M {};
define block CSTACK with alignment = 8, size = 16K {};
define block IRQ_STACK with alignment = 8, size = 16K {};
define block FIQ_STACK with alignment = 8, size = 16K {};
define block CHECKSUM
                      { ro section .checksum };
place at address Mem:0x0 { ro section .intvec};
place in ROM_region { ro, first block CHECKSUM };
place in RAM_region { rw, block HEAP, block CSTACK, block
                       IRO STACK, block FIO STACK };
```

Running ielftool

To calculate the checksum, run ielftool:

```
ielftool --fill=0x00;0x8000-0x8FFF
--checksum=__checksum:2,crc16;0x8000-0x8FFF sourceFile.out
destinationFile.out
```

To calculate a checksum you also must define a fill operation. In this example, the fill pattern 0×0 is used. The checksum algorithm used is crc16.

Note that ielftool needs an unstripped input ELF image. If you use the --strip linker option, remove it and use the --strip ielftool option instead.

ADDING A CHECKSUM FUNCTION TO YOUR SOURCE CODE

To check the value of the ielftool generated checksum, it must be compared with a checksum that your application calculated. This means that you must add a function for checksum calculation (that uses the same algorithm as ielftool) to your application source code. Your application must also include a call to this function.

A function for checksum calculation

This function—a slow variant but with small memory footprint—uses the crc16 algorithm:

```
unsigned short SlowCrc16(unsigned short sum,
                         unsigned char *p,
                         unsigned int len)
{
 while (len--)
 {
   int i;
   unsigned char byte = *(p++);
    for (i = 0; i < 8; ++i)
    {
     unsigned long oSum = sum;
     sum <<= 1;
     if (byte & 0x80)
       sum |= 1;
     if (oSum & 0x8000)
       sum ^= 0x1021;
     byte <<= 1;
   }
 }
 return sum;
}
```

You can find the source code for the checksum algorithms in the arm\src\linker directory of your product installation.

Checksum calculation

This code gives an example of how the checksum can be calculated:

```
/* The checksum calculated by ielftool
 * (note that it is located on address 0x8000)
*/
extern unsigned short const __checksum;
void TestChecksum()
{
 unsigned short calc = 0;
 unsigned char zeros[2] = \{0, 0\};
  /* Run the checksum algorithm */
  calc = SlowCrc16(0,
                  (unsigned char *) checksumStart,
                  (checksumEnd - checksumStart+1));
  /* Rotate out the answer */
 calc = SlowCrc16(calc, zeros, 2);
  /* Test the checksum */
 if (calc != __checksum)
  {
   abort(); /* Failure */
 }
}
```

THINGS TO REMEMBER

When calculating a checksum, you must remember that:

- The checksum must be calculated from the lowest to the highest address for every memory range
- Each memory range must be verified in the same order as defined
- It is OK to have several ranges for one checksum
- If several checksums are used, you should place them in sections with unique names and use unique symbol names
- If the a slow function variant is used, you must make a final call to the checksum calculation with as many bytes (with the value 0x00) as there are bytes in the checksum.

For more information, see also The IAR ELF Tool-ielftool, page 440.

C-SPY CONSIDERATIONS

By default, a symbol that you have allocated in memory by using the linker option --place_holder is considered by C-SPY to be of the type int. If the size of the checksum is different than the size of an int, you can change the display format of the checksum symbol to match its size.



In the C-SPY Watch window, select the symbol and choose **Show As** from the context menu. Choose the display format that matches the size of the checksum symbol.

Linker optimizations

VIRTUAL FUNCTION ELIMINATION

Virtual Function Elimination (VFE) is a linker optimization that removes unneeded virtual functions and dynamic runtime type information.

In order for Virtual Function Elimination to work, all relevant modules must provide information about virtual function table layout, which virtual functions are called, and for which classes dynamic runtime type information is needed. If one or more modules do not provide this information, a warning is generated by the linker and Virtual Function Elimination is not performed.

If you know that modules that lack such information do not perform any virtual function calls and do not define any virtual function tables, you can use the --vfe=forced linker option to enable Virtual Function Elimination anyway.

Currently, tools from IAR Systems and from RealView provide the information needed for Virtual Function Elimination in a way that the linker can use.

Note that you can disable Virtual Function Elimination entirely by using the --no_vfe linker option. In this case, no warning will be issued for modules that lack VFE information.

For more information, see --vfe, page 294 and --no_vfe, page 289.

AEABI compliance

The IAR build tools for ARM support the Embedded Application Binary Interface for ARM, AEABI, defined by ARM Limited. This interface is based on the Intel IA64 ABI interface. The advantage of adhering to AEABI is that any such module can be linked with any other AEABI-compliant module, even modules produced by tools provided by other vendors.

The IAR build tools for ARM support the following parts of the AEABI:

AAPCS	Procedure Call Standard for the ARM architecture
CPPABI	C++ ABI for the ARM architecture (EC++ parts only)
AAELF	ELF for the ARM architecture
AADWARF	DWARF for the ARM architecture
RTABI	Runtime ABI for the ARM architecture
CLIBABI	C library ABI for the ARM architecture

The IAR build tools only support a *bare metal* platform, that is a ROM-based system that lacks an explicit operating system.

Note that:

- The AEABI is specified for C89 only
- The IAR build tools only support using the default and C locales
- The AEABI does not specify C++ library compatibility
- Neither the size of an enum or of wchar_t is constant in the AEABI.

If AEABI compliance is enabled, almost all optimizations performed in the system header files are turned off, and certain preprocessor constants become real constant variables instead.

LINKING AEABI-COMPLIANT MODULES USING THE IAR ILINK LINKER

When building an application using the IAR ILINK Linker, the following types of modules can be combined:

- Modules produced using IAR build tools, both AEABI-compliant modules as well as modules that are not AEABI-compliant
- AEABI-compliant modules produced using build tools from another vendor.

Note: To link a module produced by a compiler from another vendor, extra support libraries from that vendor might be required.

The IAR ILINK Linker automatically chooses the appropriate standard C/C++ libraries to use based on attributes from the object files. Imported object files might not have all these attributes. Therefore, you might need to help ILINK choose the standard library by verifying one or more of the following details:

• The used CPU by specifying the --cpu linker option

- If full I/O is needed; make sure to link with a Full library configuration in the standard library
- Explicitly specify runtime library file(s), possibly in combination with the --no_library_search linker option.

When linking, you should also consider virtual function elimination, see *Virtual Function Elimination*, page 192.

LINKING AEABI-COMPLIANT MODULES USING A THIRD-PARTY LINKER

If you have a module produced using the IAR C/C++ Compiler and you plan to link that module using a linker from a different vendor, that module must be AEABI-compliant, see *Enabling AEABI compliance in the compiler*, page 194.

In addition, if that module uses any of the IAR-specific compiler extensions, you must make sure that those features are also supported by the tools from the other vendor. Note specifically:

- Support for the following extensions must be verified: #pragma pack, __no_init, __root, and __ramfunc
- The following extensions are harmless to use: #pragma location/@, __arm, __thumb, __swi, __irq, __fiq, and __nested.

ENABLING AEABI COMPLIANCE IN THE COMPILER

You can enable AEABI compliance in the compiler by setting the --aeabi option. In this case, you must also use the --guard_calls option.



In the IDE, use the **Project>Options>C/C++ Compiler>Extra Options** page to specify the --aeabi and --guard_calls options.



On the command line, use the options --aeabi and --guard_calls to enable AEABI support in the compiler.

Alternatively, to enable support for AEABI for a specific system header file, you must define the preprocessor symbol _AEABI_PORTABILITY_LEVEL to non-zero prior to including a system header file, and make sure that the symbol AEABI_PORTABLE is set to non-zero after the inclusion of the header file:

```
#define _AEABI_PORTABILITY_LEVEL 1
#undef _AEABI_PORTABLE
#include <header.h>
#ifndef _AEABI_PORTABLE
#error "header.h not AEABI compatible"
#endif
```

CMSIS integration

The arm\CMSIS subdirectory contains CMSIS (ARM Cortex Microcontroller Software Interface Standard) and CMSIS DSP header and library files, and documentation. For more information about CMSIS, see http://www.arm.com/cmsis.

The special header file inc\c\cmsis_iar.h is provided as a CMSIS adaptation of the current version of the IAR C/C++ Compiler.

CMSIS DSP LIBRARY

IAR Embedded Workbench comes with prebuilt CMSIS DSP libraries in the arm\CMSIS\Lib\IAR directory. The names of the library files are constructed in this way:

iar_cortexM<0 |3 |4><1 |b>[f]_math.a

where <0|3|4> selects the Cortex-M variant, <1|b> selects the byte order, and [f] indicates that the library is built for FPU (Cortex-M4 only).

CUSTOMIZING THE CMSIS DSP LIBRARY

The source code of the CMSIS DSP library is provided in the arm\CMSIS\DSP_Lib\Source directory. You can find an IAR Embedded Workbench project which is prepared for building a customized DSP library in the arm\CMSIS\DSP_Lib\Source\IAR directory.



BUILDING WITH CMSIS ON THE COMMAND LINE

This section contains examples of how to build your CMSIS-compatible application on the command line.

CMSIS only (that is without the DSP library)

iccarm -I \$EW_DIR\$\arm\CMSIS\Include

With the DSP library, for Cortex-M4, little-endian, and with FPU

```
iccarm --endian=little --cpu=Cortex-M4 --fpu=VFPv4_sp -I
$EW_DIR$\arm\CMSIS\Include -D ARM_MATH_CM4
```

ilinkarm \$EW_DIR\$\arm\CMSIS\Lib\IAR\iar_cortexM31_math.a



BUILDING WITH CMSIS IN IAR EMBEDDED WORKBENCH

Choose **Project>Options>General Options>Library Configuration** to enable CMSIS support.

When enabled, CMSIS include paths and the DSP library will automatically be used. For more information, see the *IDE Project Management and Building Guide for ARM*.

Efficient coding for embedded applications

For embedded systems, the size of the generated code and data is very important, because using smaller external memory or on-chip memory can significantly decrease the cost and power consumption of a system.

The topics discussed are:

- Selecting data types
- Controlling data and function placement in memory
- Controlling compiler optimizations
- Facilitating good code generation.

As a part of this, the chapter also demonstrates some of the more common mistakes and how to avoid them, and gives a catalog of good coding techniques.

Selecting data types

For efficient treatment of data, you should consider the data types used and the most efficient placement of the variables.

USING EFFICIENT DATA TYPES

The data types you use should be considered carefully, because this can have a large impact on code size and code speed.

- Use int or long instead of char or short whenever possible, to avoid sign extension or zero extension. In particular, loop indexes should always be int or long to minimize code generation. Also, in Thumb mode, accesses through the stack pointer (SP) is restricted to 32-bit data types, which further emphasizes the benefits of using one of these data types.
- Use unsigned data types, unless your application really requires signed values.
- Be aware of the costs of using 64-bit data types, such as double and long long.

- Bitfields and packed structures generate large and slow code.
- Using floating-point types on a microprocessor without a math co-processor is very inefficient, both in terms of code size and execution speed.
- Declaring a pointer to const data tells the calling function that the data pointed to will not change, which opens for better optimizations.

For information about representation of supported data types, pointers, and structures types, see the chapter *Data representation*.

FLOATING-POINT TYPES

Using floating-point types on a microprocessor without a math coprocessor is very inefficient, both in terms of code size and execution speed. Thus, you should consider replacing code that uses floating-point operations with code that uses integers, because these are more efficient.

The compiler supports two floating-point formats—32 and 64 bits. The 32-bit floating-point type float is more efficient in terms of code size and execution speed. However, the 64-bit format double supports higher precision and larger numbers.

In the compiler, the floating-point type float always uses the 32-bit format, and the type double always uses the 64-bit format.

Unless the application requires the extra precision that 64-bit floating-point numbers give, we recommend using 32-bit floating-point numbers instead.

By default, a *floating-point constant* in the source code is treated as being of the type double. This can cause innocent-looking expressions to be evaluated in double precision. In the example below, the calculation is performed using double precision, because one of the operands—the constant 1.0—is of the type double:

```
float Test(float a, float b)
{
    return a + 1.0 +b;
}
```

To treat a floating-point constant as a float rather than as a double, add the suffix f to it, for example:

```
float Test(float a, float b)
{
    return a + 1.0f + b;
}
```

For more information about floating-point types, see Floating-point types, page 304.

ALIGNMENT OF ELEMENTS IN A STRUCTURE

The ARM core requires that data in memory must be aligned. Each element in a structure must be aligned according to its specified type requirements. This means that the compiler might need to insert *pad bytes* to keep the alignment correct.

There are situations when this can be a problem:

- There are external demands; for example, network communication protocols are usually specified in terms of data types with no padding in between
- You need to save data memory.

For information about alignment requirements, see Alignment, page 297.

There are two ways to solve the problem:

- Use the #pragma pack directive or the __packed data type attribute for a tighter layout of the structure. The drawback is that each access to an unaligned element in the structure will use more code.
- Write your own customized functions for *packing* and *unpacking* structures. This is a more portable way, which will not produce any more code apart from your functions. The drawback is the need for two views on the structure data—packed and unpacked.

For more information about the #pragma pack directive, see pack, page 337.

ANONYMOUS STRUCTS AND UNIONS

When a structure or union is declared without a name, it becomes anonymous. The effect is that its members will only be seen in the surrounding scope.

Anonymous structures are part of the C++ language; however, they are not part of the C standard. In the IAR C/C++ Compiler for ARM they can be used in C if language extensions are enabled.



In the IDE, language extensions are enabled by default.



Use the -e compiler option to enable language extensions. See -e, page 242, for additional information.

Example

In this example, the members in the anonymous union can be accessed, in function F, without explicitly specifying the union name:

```
struct S
{
    char mTag;
    union
    {
        long mL;
        float mF;
    };
    St;
void F(void)
{
    St.mL = 5;
}
```

The member names must be unique in the surrounding scope. Having an anonymous struct or union at file scope, as a global, external, or static variable is also allowed. This could for instance be used for declaring I/O registers, as in this example:

```
___no__init volatile
union
{
 unsigned char IOPORT;
 struct
  {
    unsigned char way: 1;
   unsigned char out: 1;
 };
} @ 0x1000;
/* The variables are used here. */
void Test(void)
{
 IOPORT = 0;
 way = 1;
 out = 1;
}
```

This declares an I/O register byte IOPORT at address 0x1000. The I/O register has 2 bits declared, Way and Out. Note that both the inner structure and the outer union are anonymous.

Anonymous structures and unions are implemented in terms of objects named after the first field, with a prefix _A_ to place the name in the implementation part of the namespace. In this example, the anonymous union will be implemented through an object named _A_IOPORT.

Controlling data and function placement in memory

The compiler provides different mechanisms for controlling placement of functions and data objects in memory. To use memory efficiently, you should be familiar with these mechanisms to know which one is best suited for different situations. You can use:

• The @ operator and the #pragma location directive for absolute placement

Use the @ operator or the #pragmalocation directive to place individual global and static variables at absolute addresses.

This is useful for individual data objects that must be located at a fixed address to conform to external requirements, for example to populate interrupt vectors or other hardware tables. Note that it is not possible to use this notation for absolute placement of individual functions.

• The @ operator and the #pragma location directive for section placement

Use the @ operator or the #pragma location directive to place groups of functions or global and static variables in named sections, without having explicit control of each object. The sections can, for example, be placed in specific areas of memory, or initialized or copied in controlled ways using the section begin and end operators. This is also useful if you want an interface between separately linked units, for example an application project and a boot loader project. Use named sections when absolute control over the placement of individual variables is not needed, or not useful.

• The @ operator and the #pragma location directive for register placement

Use the @ operator or the #pragma location directive to place individual global and static variables in registers. The variables must be declared __no_init. This is useful for individual data objects that must be located in a specific register.

• The --section option

Use the --section option to place functions and/or data objects in named sections, which is useful, for example, if you want to direct them to different fast or slow memories. For more information, see *--section*, page 263.

At compile time, data and functions are placed in different sections, see *Modules and sections*, page 70. At link time, one of the most important functions of the linker is to assign load addresses to the various sections used by the application. All sections, except for the sections holding absolute located data, are automatically allocated to memory according to the specifications in the linker configuration file, see *Placing code and data*—the linker configuration file, page 73.

DATA PLACEMENT AT AN ABSOLUTE LOCATION

The @ operator, alternatively the #pragma location directive, can be used for placing global and static variables at absolute addresses.

To place a variable at an absolute address, the argument to the @ operator and the #pragma location directive should be a literal number, representing the actual address. The absolute location must fulfill the alignment requirement for the variable that should be located.

Note: All declarations of __no_init variables placed at an absolute address are *tentative definitions*. Tentatively defined variables are only kept in the output from the compiler if they are needed in the module being compiled. Such variables will be defined in all modules in which they are used, which will work as long as they are defined in the same way. The recommendation is to place all such declarations in header files that are included in all modules that use the variables.

Other variables placed at an absolute address use the normal distinction between declaration and definition. For these variables, you must provide the definition in only one module, normally with an initializer. Other modules can refer to the variable by using an extern declaration, with or without an explicit address.

Examples

In this example, a <u>__no_init</u> declared variable is placed at an absolute address. This is useful for interfacing between multiple processes, applications, etc:

__no_init volatile char alpha @ 0xFF2000;/* OK */

The next example contains two const declared objects. The first one is not initialized, and the second one is initialized to a specific value. Both objects are placed in ROM. The next example contains a const declared object which is not initialized. The object is placed in ROM. This is useful for configuration parameters, which are accessible from an external interface. Note that in the second case, the compiler is not obliged to actually read from the variable, because the value is known.

<pre>#pragma location=0xFF2002</pre>				
no init const int beta;	/* OK */			
const int gamma @ 0xFF2004 = 3;	/* OK */			

In the first case, the value is not initialized by the compiler; the value must be set by other means. The typical use is for configurations where the values are loaded to ROM separately, or for special function registers that are read-only.

These examples show incorrect usage:

```
int delta @ 0xFF2006; /* Error, not __no_init *//*
Error, neither */
/* "__no_init" nor "const".*/
__no_init int epsilon @ 0xFF2007; /* Error, misaligned. */
```

C++ considerations

In C++, module scoped const variables are static (module local), whereas in C they are global. This means that each module that declares a certain const variable will contain a separate variable with this name. If you link an application with several such modules all containing (via a header file), for instance, the declaration:

volatile const __no_init int x @ 0x100; /* Bad in C++ */

the linker will report that more than one variable is located at address 0x100.

To avoid this problem and make the process the same in C and C++, you should declare these variables extern, for example:

/* The extern keyword makes x public. */
extern volatile const __no_init int x @ 0x100;

Note: C++ static member variables can be placed at an absolute address just like any other static variable.

DATA AND FUNCTION PLACEMENT IN SECTIONS

The following methods can be used for placing data or functions in named sections other than default:

- The @ operator, alternatively the #pragma location directive, can be used for placing individual variables or individual functions in named s. The named section can either be a predefined section, or a user-defined section.
- The --section option can be used for placing variables and functions, which are parts of the whole compilation unit, in named sections.

C++ static member variables can be placed in named sections just like any other static variable.

If you use your own sections, in addition to the predefined sections, the sections must also be defined in the linker configuration file. **Note:** Take care when explicitly placing a variable or function in a predefined section other than the one used by default. This is useful in some situations, but incorrect placement can result in anything from error messages during compilation and linking to a malfunctioning application. Carefully consider the circumstances; there might be strict requirements on the declaration and use of the function or variable.

The location of the sections can be controlled from the linker configuration file.

For more information about sections, see the chapter Section reference.

Examples of placing variables in named sections

In the following examples, a data object is placed in a user-defined section.

```
__no_init int alpha @ "MY_NOINIT"; /* OK */
#pragma location="MY_CONSTANTS"
const int beta = 42; /* OK */
const int gamma @ "MY_CONSTANTS" = 17;/* OK */
int theta @ "MY_ZEROS"; /* OK */
int phi @ "MY_INITED" = 4711; /* OK */
```

The linker will in the zero and initialized cases arrange for the correct type of initialization for the variable. When placing a __no_init variable in a user-defined section, you must add a pattern that matches that section to your do not initialize directive in the linker configuration file. For initialized variables, you can disable the automatic initialization by using the initialize manually directive. **Examples of placing functions in named sections**

```
void f(void) @ "MY_FUNCTIONS";
void g(void) @ "MY_FUNCTIONS"
{
}
#pragma location="MY_FUNCTIONS"
void h(void);
```

DATA PLACEMENT IN REGISTERS

The @ operator, alternatively the #pragma location directive, can be used for placing global and static variables in a register.

To place a variable in a register, the argument to the @ operator and the #pragma location directive should be an identifier that corresponds to an ARM core register in the range R4-R11 (R9 cannot be specified in combination with the --rwpi command line option). A variable can be placed in a register only if it is declared as __no_init, has file scope, and its size is four bytes. A variable placed in a register does not have a memory address, so the address operator & cannot be used.

Within a module where a variable is placed in a register, the specified register will only be used for accessing that variable. The value of the variable is preserved across function calls to other modules because the registers R14-R11 are callee saved, and as such they are restored when execution returns. However, the value of a variable placed in a register is not always preserved as expected:

- In an exception handler or library callback routine (such as the comparator function passed to qsort) the value might not be preserved. The value will be preserved if the command line option --lock_regs is used for locking the register in all modules of the application, including library modules.
- In a fast interrupt handler, the value of a variable in R8-R11 is not preserved from outside the handler, because these registers are banked.
- The longjmp function and C++ exceptions might restore variables placed in registers to old values, unlike other variables with static storage duration which are not restored.

The linker does not prevent modules from placing different variables in the same register. Variables in different modules can be placed in the same register, and another module could use the register for other purposes.

Note: A variable placed in a register should be defined in an include file, to be included in every module that uses the variable. An unused definition in a module will cause the register to not be used in that module.

Controlling compiler optimizations

The compiler performs many transformations on your application to generate the best possible code. Examples of such transformations are storing values in registers instead of memory, removing superfluous code, reordering computations in a more efficient order, and replacing arithmetic operations by cheaper operations.

The linker should also be considered an integral part of the compilation system, because some optimizations are performed by the linker. For instance, all unused functions and variables are removed and not included in the final output.

SCOPE FOR PERFORMED OPTIMIZATIONS

You can decide whether optimizations should be performed on your whole application or on individual files. By default, the same types of optimizations are used for an entire project, but you should consider using different optimization settings for individual files. For example, put code that must execute very quickly into a separate file and compile it for minimal execution time, and the rest of the code for minimal code size. This will give a small program, which is still fast enough where it matters.

You can also exclude individual functions from the performed optimizations. The #pragma optimize directive allows you to either lower the optimization level, or specify another type of optimization to be performed. See *optimize*, page 335, for information about the pragma directive.

MULTI-FILE COMPILATION UNITS

In addition to applying different optimizations to different source files or even functions, you can also decide what a compilation unit consists of—one or several source code files.

By default, a compilation unit consists of one source file, but you can also use multi-file compilation to make several source files in a compilation unit. The advantage is that interprocedural optimizations such as inlining, and cross jump have more source code to work on. Ideally, the whole application should be compiled as one compilation unit. However, for large applications this is not practical because of resource restrictions on the host computer. For more information, see *--mfc*, page 250.

If the whole application is compiled as one compilation unit, it is very useful to make the compiler also discard unused public functions and variables before the interprocedural optimizations are performed. Doing this limits the scope of the optimizations to functions and variables that are actually used. For more information, see *--discard unused publics*, page 241.

OPTIMIZATION LEVELS

The compiler supports different levels of optimizations. This table lists optimizations that are typically performed on each level:

Optimization level	Description
None (Best debug support)	Variables live through their entire scope
Low	Same as above but variables only live for as long as they are needed, not necessarily through their entire scope
Medium	Same as above, and: Live-dead analysis and optimization Dead code elimination Redundant label elimination Redundant branch elimination Code hoisting Peephole optimization Some register content analysis and optimization Static clustering Common subexpression elimination
High (Balanced)	Same as above, and: Instruction scheduling Cross jumping Advanced register content analysis and optimization Loop unrolling Function inlining Code motion Type-based alias analysis

Table 26: Compiler optimization levels

Note: Some of the performed optimizations can be individually enabled or disabled. For more information about these, see *Fine-tuning enabled transformations*, page 208.

A high level of optimization might result in increased compile time, and will most likely also make debugging more difficult, because it is less clear how the generated code relates to the source code. For example, at the low, medium, and high optimization levels, variables do not live through their entire scope, which means processor registers used for storing variables can be reused immediately after they were last used. Due to this, the C-SPY Watch window might not be able to display the value of the variable throughout its scope. At any time, if you experience difficulties when debugging your code, try lowering the optimization level.

SPEED VERSUS SIZE

At the high optimization level, the compiler balances between size and speed optimizations. However, it is possible to fine-tune the optimizations explicitly for either size or speed. They only differ in what thresholds that are used; speed will trade size for speed, whereas size will trade speed for size. Note that one optimization sometimes enables other optimizations to be performed, and an application might in some cases become smaller even when optimizing for speed rather than size.

If you use the optimization level High speed, the --no_size_constraints compiler option relaxes the normal restrictions for code size expansion and enables more aggressive optimizations.

FINE-TUNING ENABLED TRANSFORMATIONS

At each optimization level you can disable some of the transformations individually. To disable a transformation, use either the appropriate option, for instance the command line option --no_inline, alternatively its equivalent in the IDE **Function inlining**, or the #pragma optimize directive. These transformations can be disabled individually:

- Common subexpression elimination
- Loop unrolling
- Function inlining
- Code motion
- Type-based alias analysis
- Static clustering
- Instruction scheduling.

Common subexpression elimination

Redundant re-evaluation of common subexpressions is by default eliminated at optimization levels Medium and High. This optimization normally reduces both code size and execution time. However, the resulting code might be difficult to debug.

Note: This option has no effect at optimization levels None and Low.

For more information about the command line option, see --no cse, page 252.

Loop unrolling

Loop unrolling means that the code body of a loop, whose number of iterations can be determined at compile time, is duplicated. Loop unrolling reduces the loop overhead by amortizing it over several iterations.

This optimization is most efficient for smaller loops, where the loop overhead can be a substantial part of the total loop body.

Loop unrolling, which can be performed at optimization level High, normally reduces execution time, but increases code size. The resulting code might also be difficult to debug.

The compiler heuristically decides which loops to unroll. Only relatively small loops where the loop overhead reduction is noticeable will be unrolled. Different heuristics are used when optimizing for speed, size, or when balancing between size and speed.

Note: This option has no effect at optimization levels None, Low, and Medium.

To disable loop unrolling, use the command line option --no_unroll, see --*no_unroll*, page 257.

Function inlining

Function inlining means that a function, whose definition is known at compile time, is integrated into the body of its caller to eliminate the overhead of the call. This optimization normally reduces execution time, but might increase the code size.

For more information, see Inlining functions, page 65.

Code motion

Evaluation of loop-invariant expressions and common subexpressions are moved to avoid redundant re-evaluation. This optimization, which is performed at optimization level High, normally reduces code size and execution time. The resulting code might however be difficult to debug.

Note: This option has no effect at optimization levels None, and Low.

For more information about the command line option, see --no_code_motion, page 251.

Type-based alias analysis

When two or more pointers reference the same memory location, these pointers are said to be *aliases* for each other. The existence of aliases makes optimization more difficult because it is not necessarily known at compile time whether a particular value is being changed.

Type-based alias analysis optimization assumes that all accesses to an object are performed using its declared type or as a char type. This assumption lets the compiler detect whether pointers can reference the same memory location or not.

Type-based alias analysis is performed at optimization level High. For application code conforming to standard C or C++ application code, this optimization can reduce code size and execution time. However, non-standard C or C++ code might result in the compiler producing code that leads to unexpected behavior. Therefore, it is possible to turn this optimization off.

Note: This option has no effect at optimization levels None, Low, and Medium.

For more information about the command line option, see --no_tbaa, page 256.

Example

```
short F(short *p1, long *p2)
{
    *p2 = 0;
    *p1 = 1;
    return *p2;
}
```

With type-based alias analysis, it is assumed that a write access to the short pointed to by p1 cannot affect the long value that p2 points to. Thus, it is known at compile time that this function returns 0. However, in non-standard-conforming C or C++ code these pointers could overlap each other by being part of the same union. If you use explicit casts, you can also force pointers of different pointer types to point to the same memory location.

Static clustering

When static clustering is enabled, static and global variables that are defined within the same module are arranged so that variables that are accessed in the same function are stored close to each other. This makes it possible for the compiler to use the same base pointer for several accesses.

Note: This option has no effect at optimization levels None and Low.

For more information about the command line option, see --no_clustering, page 251.

Instruction scheduling

The compiler features an instruction scheduler to increase the performance of the generated code. To achieve that goal, the scheduler rearranges the instructions to minimize the number of pipeline stalls emanating from resource conflicts within the microprocessor. Note that not all cores benefit from scheduling. The resulting code might be difficult to debug.

Note: This option has no effect at optimization levels None, Low and Medium.

For more information about the command line option, see --no_scheduling, page 255.

Facilitating good code generation

This section contains hints on how to help the compiler generate good code, for example:

- Using efficient addressing modes
- Helping the compiler optimize
- Generating more useful error message.

WRITING OPTIMIZATION-FRIENDLY SOURCE CODE

The following is a list of programming techniques that will, when followed, enable the compiler to better optimize the application.

- Local variables—auto variables and parameters—are preferred over static or global variables. The reason is that the optimizer must assume, for example, that called functions can modify non-local variables. When the life spans for local variables end, the previously occupied memory can then be reused. Globally declared variables will occupy data memory during the whole program execution.
- Avoid taking the address of local variables using the & operator. This is inefficient for two main reasons. First, the variable must be placed in memory, and thus cannot be placed in a processor register. This results in larger and slower code. Second, the optimizer can no longer assume that the local variable is unaffected over function calls.
- Module-local variables—variables that are declared static—are preferred over global variables (non-static). Also avoid taking the address of frequently accessed static variables.
- The compiler is capable of inlining functions, see *Function inlining*, page 209. To maximize the effect of the inlining transformation, it is good practice to place the definitions of small functions called from more than one module in the header file rather than in the implementation file. Alternatively, you can use multi-file compilation. For more information, see *Multi-file compilation units*, page 206.
- Avoid using inline assembler without operands and clobbered resources. Instead, use SFRs or intrinsic functions if available. Otherwise, use inline assembler *with* operands and clobbered resources or write a separate module in assembler language. For more information, see *Mixing C and assembler*, page 137.

SAVING STACK SPACE AND RAM MEMORY

The following is a list of programming techniques that will, when followed, save memory and stack space:

• If stack space is limited, avoid long call chains and recursive functions.

• Avoid using large non-scalar types, such as structures, as parameters or return type. To save stack space, you should instead pass them as pointers or, in C++, as references.

FUNCTION PROTOTYPES

It is possible to declare and define functions using one of two different styles:

- Prototyped
- Kernighan & Ritchie C (K&R C)

Both styles are valid C, however it is strongly recommended to use the prototyped style, and provide a prototype declaration for each public function in a header that is included both in the compilation unit defining the function and in all compilation units using it.

The compiler will not perform type checking on parameters passed to functions declared using K&R style. Using prototype declarations will also result in more efficient code in some cases, as there is no need for type promotion for these functions.

To make the compiler require that all function definitions use the prototyped style, and that all public functions have been declared before being defined, use the **Project>Options>C/C++ Compiler>Language 1>Require prototypes** compiler option (--require_prototypes).

Prototyped style

In prototyped function declarations, the type for each parameter must be specified.

```
int Test(char, int); /* Declaration */
int Test(char ch, int i) /* Definition */
{
   return i + ch;
}
```

Kernighan & Ritchie style

In K&R style—pre-Standard C—it is not possible to declare a function prototyped. Instead, an empty parameter list is used in the function declaration. Also, the definition looks different.

For example:

```
int Test(); /* Declaration */
int Test(ch, i) /* Definition */
char ch;
int i;
{
  return i + ch;
}
```

INTEGER TYPES AND BIT NEGATION

In some situations, the rules for integer types and their conversion lead to possibly confusing behavior. Things to look out for are assignments or conditionals (test expressions) involving types with different size, and logical operations, especially bit negation. Here, *types* also includes types of constants.

In some cases there might be warnings (for example, for constant conditional or pointless comparison), in others just a different result than what is expected. Under certain circumstances the compiler might warn only at higher optimizations, for example, if the compiler relies on optimizations to identify some instances of constant conditionals. In this example an 8-bit character, a 32-bit integer, and two's complement is assumed:

```
void F1(unsigned char c1)
{
    if (c1 == ~0x80)
    ;
}
```

Here, the test is always false. On the right hand side, 0×80 is 0×00000080 , and $\sim 0 \times 00000080$ becomes $0 \times FFFFF7F$. On the left hand side, c1 is an 8-bit unsigned character, so it cannot be larger than 255. It also cannot be negative, which means that the integral promoted value can never have the topmost 24 bits set.

PROTECTING SIMULTANEOUSLY ACCESSED VARIABLES

Variables that are accessed asynchronously, for example by interrupt routines or by code executing in separate threads, must be properly marked and have adequate protection. The only exception to this is a variable that is always *read-only*.

To mark a variable properly, use the volatile keyword. This informs the compiler, among other things, that the variable can be changed from other threads. The compiler will then avoid optimizing on the variable (for example, keeping track of the variable in registers), will not delay writes to it, and be careful accessing the variable only the number of times given in the source code.

For more information about the volatile type qualifier and the rules for accessing volatile objects, see *Declaring objects volatile*, page 309.

ACCESSING SPECIAL FUNCTION REGISTERS

Specific header files for several ARM devices are included in the IAR product installation. The header files are named *iodevice*.h and define the processor-specific special function registers (SFRs).

Note: Each header file contains one section used by the compiler, and one section used by the assembler.

SFRs with bitfields are declared in the header file. This example is from ioks32c5000a.h:

```
___no__init volatile union
{
 unsigned short mwctl2;
 struct
 {
   unsigned short edr: 1;
   unsigned short edw: 1;
   unsigned short lee: 2;
   unsigned short lemd: 2;
   unsigned short lepl: 2;
 } mwctl2bit;
} @ 0x1000;
/* By including the appropriate include file in your code,
* it is possible to access either the whole register or any
 * individual bit (or bitfields) from C code as follows.
 */
void Test()
{
 /* Whole register access */
 mwctl2 = 0x1234;
 /* Bitfield accesses */
 mwctl2bit.edw = 1;
 mwctl2bit.lepl = 3;
}
```

You can also use the header files as templates when you create new header files for other ARM devices.

PASSING VALUES BETWEEN C AND ASSEMBLER OBJECTS

The following example shows how you in your C source code can use inline assembler to set and get values from a special purpose register:

```
static unsigned long get_APSR( void )
{
    unsigned long value;
    asm volatile( "MRS %0, APSR" : "=r"(value) );
    return value;
}
static void set_APSR( unsigned long value)
{
    asm volatile( "MSR APSR, %0" :: "r"(value) );
}
```

The general purpose register is used for getting and setting the value of the special purpose register APSR. The same method can be used also for accessing other special purpose registers and specific instructions.

To read more about inline assembler, see *Inline assembler*, page 139. An alternative to inline assembler can be to use an intrinsic function instead. See *Summary of intrinsic functions*, page 345.

NON-INITIALIZED VARIABLES

Normally, the runtime environment will initialize all global and static variables when the application is started.

The compiler supports the declaration of variables that will not be initialized, using the __no_init type modifier. They can be specified either as a keyword or using the #pragma object_attribute directive. The compiler places such variables in a separate section.

For __no_init, the const keyword implies that an object is read-only, rather than that the object is stored in read-only memory. It is not possible to give a __no_init object an initial value.

Variables declared using the __no_init keyword could, for example, be large input buffers or mapped to special RAM that keeps its content even when the application is turned off.

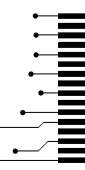
For more information, see __*no_init*, page 319. Note that to use this keyword, language extensions must be enabled; see *-e*, page 242. For more information, see also *object_attribute*, page 335.

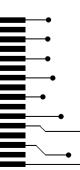
Facilitating good code generation

Part 2. Reference information

This part of the IAR C/C++ Development Guide for ARM contains these chapters:

- External interface details
- Compiler options
- Linker options
- Data representation
- Extended keywords
- Pragma directives
- Intrinsic functions
- The preprocessor
- Library functions
- The linker configuration file
- Section reference
- Stack usage control files
- IAR utilities
- Implementation-defined behavior for Standard C
- Implementation-defined behavior for C89.





External interface details

This chapter provides reference information about how the compiler and linker interact with their environment. The chapter briefly lists and describes the invocation syntax, methods for passing options to the tools, environment variables, the include file search procedure, and finally the different types of compiler and linker output.

Invocation syntax

You can use the compiler and linker either from the IDE or from the command line. See the *IDE Project Management and Building Guide for ARM* for information about using the build tools from the IDE.

COMPILER INVOCATION SYNTAX

The invocation syntax for the compiler is:

iccarm [options] [sourcefile] [options]

For example, when compiling the source file prog.c, use this command to generate an object file with debug information:

iccarm prog.c --debug

The source file can be a C or C++ file, typically with the filename extension c or cpp, respectively. If no filename extension is specified, the file to be compiled must have the extension c.

Generally, the order of options on the command line, both relative to each other and to the source filename, is not significant. There is, however, one exception: when you use the -I option, the directories are searched in the same order as they are specified on the command line.

If you run the compiler from the command line without any arguments, the compiler version number and all available options including brief descriptions are directed to stdout and displayed on the screen.

ILINK INVOCATION SYNTAX

The invocation syntax for ILINK is:

ilinkarm [arguments]

Each argument is either a command-line option, an object file, or a library.

For example, when linking the object file prog.o, use this command:

ilinkarm prog.o --config configfile

If no filename extension is specified for the linker configuration file, the configuration file must have the extension *icf*.

Generally, the order of arguments on the command line is not significant. There is, however, one exception: when you supply several libraries, the libraries are searched in the same order that they are specified on the command line. The default libraries are always searched last.

The output executable image will be placed in a file named a . out, unless the -o option is used.

If you run ILINK from the command line without any arguments, the ILINK version number and all available options including brief descriptions are directed to stdout and displayed on the screen.

PASSING OPTIONS

There are three different ways of passing options to the compiler and to ILINK:

• Directly from the command line

Specify the options on the command line after the iccarm or ilinkarm commands; see *Invocation syntax*, page 219.

• Via environment variables

The compiler and linker automatically append the value of the environment variables to every command line; see *Environment variables*, page 220.

• Via a text file, using the -f option; see -f, page 245.

For general guidelines for the option syntax, an options summary, and a detailed description of each option, see the *Compiler options* chapter.

ENVIRONMENT VARIABLES

These environment variables can be used with the compiler:

Environment variable	Description
----------------------	-------------

C_INCLUDE	Specifies directories to search for include files; for example:
	C_INCLUDE=c:\program files\iar systems\embedded
	<pre>workbench 6.n\arm\inc;c:\headers</pre>
QCCARM	Specifies command line options; for example: $\ensuremath{\texttt{QCCARM}=-lA}$
	asm.lst

Table 27: Compiler environment variables

This environment variable can be used with ILINK:

Environment variable	Description
ILINKARM_CMD_LINE	Specifies command line options; for example:
	ILINKARM_CMD_LINE=config full.icf
	silent

Table 28: ILINK environment variables

Include file search procedure

This is a detailed description of the compiler's #include file search procedure:

- If the name of the #include file is an absolute path specified in angle brackets or double quotes, that file is opened.
- If the compiler encounters the name of an #include file in angle brackets, such as: #include <stdio.h>

it searches these directories for the file to include:

- 1 The directories specified with the -I option, in the order that they were specified, see -*I*, page 247.
- 2 The directories specified using the C_INCLUDE environment variable, if any; see *Environment variables*, page 220.
- 3 The automatically set up library system include directories. See --*dlib_config*, page 242.
- If the compiler encounters the name of an #include file in double quotes, for example:

#include "vars.h"

it searches the directory of the source file in which the #include statement occurs, and then performs the same sequence as for angle-bracketed filenames.

If there are nested #include files, the compiler starts searching the directory of the file that was last included, iterating upwards for each included file, searching the source file directory last. For example:

```
src.c in directory dir\src
    #include "src.h"
    ...
src.h in directory dir\include
    #include "config.h"
    ...
```

When dir\exe is the current directory, use this command for compilation:

iccarm ..\src\src.c -I..\include -I..\debugconfig

Then the following directories are searched in the order listed below for the file config.h, which in this example is located in the dir\debugconfig directory:

dir\include	Current file is src.h.
dir\src	File including current file (src.c).
dir\include	As specified with the first -I option.
dir\debugconfig	As specified with the second $-I$ option.

Use angle brackets for standard header files, like stdio.h, and double quotes for files that are part of your application.

For information about the syntax for including header files, see *Overview of the preprocessor*, page 381.

Compiler output

The compiler can produce the following output:

• A linkable object file

The object files produced by the compiler use the industry-standard format ELF. By default, the object file has the filename extension \circ .

Optional list files

Various kinds of list files can be specified using the compiler option -1, see -*l*, page 248. By default, these files will have the filename extension lst.

Optional preprocessor output files

A preprocessor output file is produced when you use the --preprocess option; by default, the file will have the filename extension i.

• Diagnostic messages

Diagnostic messages are directed to the standard error stream and displayed on the screen, and printed in an optional list file. For more information about diagnostic messages, see *Diagnostics*, page 224.

Error return codes

These codes provide status information to the operating system which can be tested in a batch file, see *Error return codes*, page 223.

• Size information

Information about the generated amount of bytes for functions and data for each memory is directed to the standard output stream and displayed on the screen. Some of the bytes might be reported as *shared*.

Shared objects are functions or data objects that are shared between modules. If any of these occur in more than one module, only one copy is retained. For example, in some cases inline functions are not inlined, which means that they are marked as shared, because only one instance of each function will be included in the final application. This mechanism is sometimes also used for compiler-generated code or data not directly associated with a particular function or variable, and when only one instance is required in the final application.

ERROR RETURN CODES

The compiler and linker return status information to the operating system that can be tested in a batch file.

These command line error codes are supported:

Code	Description
0	Compilation or linking successful, but there might have been warnings.
I	Warnings were produced and the optionwarnings_affect_exit_code was used.
2	Errors occurred.
3	Fatal errors occurred, making the tool abort.
4	Internal errors occurred, making the tool abort.
Table 20	· Ennon nations and as

Table 29: Error return codes

ILINK output

ILINK can produce the following output:

• An absolute executable image

The final output produced by the IAR ILINK Linker is an absolute object file containing the executable image that can be put into an EPROM, downloaded to a hardware emulator, or executed on your PC using the IAR C-SPY Debugger Simulator. By default, the file has the filename extension out. The output format is always in ELF, which optionally includes debug information in the DWARF format.

• Optional logging information

During operation, ILINK logs its decisions on stdout, and optionally to a file. For example, if a library is searched, whether a required symbol is found in a library module, or whether a module will be part of the output. Timing information for each ILINK subsystem is also logged.

Optional map files

A linker map file—containing summaries of linkage, runtime attributes, memory, and placement, as well as an entry list— can be generated by the ILINK option

--map, see --map, page 285. By default, the map file has the filename extension map.

Diagnostic messages

Diagnostic messages are directed to stderr and displayed on the screen, as well as printed in the optional map file. For more information about diagnostic messages, see *Diagnostics*, page 224.

• Error return codes

ILINK returns status information to the operating system which can be tested in a batch file, see *Error return codes*, page 223.

• Size information about used memory and amount of time

Information about the generated amount of bytes for functions and data for each memory is directed to stdout and displayed on the screen.

Diagnostics

This section describes the format of the diagnostic messages and explains how diagnostic messages are divided into different levels of severity.

MESSAGE FORMAT FOR THE COMPILER

All diagnostic messages are issued as complete, self-explanatory messages. A typical diagnostic message from the compiler is produced in the form:

filename,linenumber level[tag]: message

with these elements:

filename	The name of the source file in which the issue was encountered
linenumber	The line number at which the compiler detected the issue
level	The level of seriousness of the issue
tag	A unique tag that identifies the diagnostic message
message	An explanation, possibly several lines long

Diagnostic messages are displayed on the screen, as well as printed in the optional list file.

Use the option --diagnostics_tables to list all possible compiler diagnostic messages.

MESSAGE FORMAT FOR THE LINKER

All diagnostic messages are issued as complete, self-explanatory messages. A typical diagnostic message from ILINK is produced in the form:

level[tag]: message

with these elements:

level	The level of seriousness of the issue
tag	A unique tag that identifies the diagnostic message
message	An explanation, possibly several lines long

Diagnostic messages are displayed on the screen, as well as printed in the optional map file.

Use the option --diagnostics_tables to list all possible linker diagnostic messages.

SEVERITY LEVELS

The diagnostic messages are divided into different levels of severity:

Remark

A diagnostic message that is produced when the compiler or linker finds a construct that can possibly lead to erroneous behavior in the generated code. Remarks are by default not issued, but can be enabled, see *--remarks*, page 262.

Warning

A diagnostic message that is produced when the compiler or linker finds a potential problem which is of concern, but which does not prevent completion of the compilation or linking. Warnings can be disabled by use of the command line option --no_warnings, see page 258.

Error

A diagnostic message that is produced when the compiler or linker finds a serious error. An error will produce a non-zero exit code.

Fatal error

A diagnostic message that is produced when the compiler finds a condition that not only prevents code generation, but which makes further processing pointless. After the message is issued, compilation terminates. A fatal error will produce a non-zero exit code.

SETTING THE SEVERITY LEVEL

The diagnostic messages can be suppressed or the severity level can be changed for all diagnostics messages, except for fatal errors and some of the regular errors.

See *Summary of compiler options*, page 229, for information about the compiler options that are available for setting severity levels.

For the compiler see also the chapter *Pragma directives*, for information about the pragma directives that are available for setting severity levels.

INTERNAL ERROR

An internal error is a diagnostic message that signals that there was a serious and unexpected failure due to a fault in the compiler or linker. It is produced using this form:

Internal error: message

where *message* is an explanatory message. If internal errors occur, they should be reported to your software distributor or IAR Systems Technical Support. Include enough information to reproduce the problem, typically:

- The product name
- The version number of the compiler or of ILINK, which can be seen in the header of the list or map files generated by the compiler or by ILINK, respectively
- Your license number
- The exact internal error message text
- The files involved of the application that generated the internal error
- A list of the options that were used when the internal error occurred.

Compiler options

This chapter describes the syntax of compiler options and the general syntax rules for specifying option parameters, and gives detailed reference information about each option.

Options syntax

Compiler options are parameters you can specify to change the default behavior of the compiler. You can specify options from the command line—which is described in more detail in this section—and from within the IDE.



See the *IDE Project Management and Building Guide for ARM* for information about the compiler options available in the IDE and how to set them.

TYPES OF OPTIONS

There are two *types of names* for command line options, *short* names and *long* names. Some options have both.

- A short option name consists of one character, and it can have parameters. You specify it with a single dash, for example -e
- A long option name consists of one or several words joined by underscores, and it can have parameters. You specify it with double dashes, for example --char_is_signed.

For information about the different methods for passing options, see *Passing options*, page 220.

RULES FOR SPECIFYING PARAMETERS

There are some general syntax rules for specifying option parameters. First, the rules depending on whether the parameter is *optional* or *mandatory*, and whether the option has a short or a long name, are described. Then, the rules for specifying filenames and directories are listed. Finally, the remaining rules are listed.

Rules for optional parameters

For options with a short name and an optional parameter, any parameter should be specified without a preceding space, for example:

-0 or -0h

For options with a long name and an optional parameter, any parameter should be specified with a preceding equal sign (=), for example:

--misrac2004=n

Rules for mandatory parameters

For options with a short name and a mandatory parameter, the parameter can be specified either with or without a preceding space, for example:

-I..\src or -I ..\src\

For options with a long name and a mandatory parameter, the parameter can be specified either with a preceding equal sign (=) or with a preceding space, for example:

--diagnostics_tables=MyDiagnostics.lst

or

--diagnostics_tables MyDiagnostics.lst

Rules for options with both optional and mandatory parameters

For options taking both optional and mandatory parameters, the rules for specifying the parameters are:

- For short options, optional parameters are specified without a preceding space
- For long options, optional parameters are specified with a preceding equal sign (=)
- For short and long options, mandatory parameters are specified with a preceding space.

For example, a short option with an optional parameter followed by a mandatory parameter:

-lA MyList.lst

For example, a long option with an optional parameter followed by a mandatory parameter:

--preprocess=n PreprocOutput.lst

Rules for specifying a filename or directory as parameters

These rules apply for options taking a filename or directory as parameters:

• Options that take a filename as a parameter can optionally take a file path. The path can be relative or absolute. For example, to generate a listing to the file List.lst in the directory ...listings:

iccarm prog.c -1 ...listings\List.lst

• For options that take a filename as the destination for output, the parameter can be specified as a path without a specified filename. The compiler stores the output in that directory, in a file with an extension according to the option. The filename will be the same as the name of the compiled source file, unless a different name was specified with the option -0, in which case that name is used. For example:

```
iccarm prog.c -l ..\listings\
```

The produced list file will have the default name ... \listings\prog.lst

• The current directory is specified with a period (.). For example:

iccarm prog.c -1 .

- / can be used instead of \setminus as the directory delimiter.
- By specifying –, input files and output files can be redirected to the standard input and output stream, respectively. For example:

iccarm prog.c -1 -

Additional rules

These rules also apply:

• When an option takes a parameter, the parameter cannot start with a dash (-) followed by another character. Instead, you can prefix the parameter with two dashes; this example will create a list file called -r:

iccarm prog.c -1 ---r

• For options that accept multiple arguments of the same type, the arguments can be provided as a comma-separated list (without a space), for example:

--diag_warning=Be0001,Be0002

Alternatively, the option can be repeated for each argument, for example:

```
--diag_warning=Be0001
--diag_warning=Be0002
```

Summary of compiler options

This table summarizes the compiler command line options:

Command line option	Description
aapcs	Specifies the calling convention
aeabi	Enables AEABI-compliant code generation
align_sp_on_irq	Generates code to align SP on entry toirq functions
arm	Sets the default function mode to ARM

Table 30: Compiler options summary

Command line option	Description
c89	Specifies the C89 dialect
char_is_signed	Treats char as signed
char_is_unsigned	Treats char as unsigned
cpu	Specifies a processor variant
C++	Specifies Standard C++
-D	Defines preprocessor symbols
debug	Generates debug information
dependencies	Lists file dependencies
diag_error	Treats these as errors
diag_remark	Treats these as remarks
diag_suppress	Suppresses these diagnostics
diag_warning	Treats these as warnings
diagnostics_tables	Lists all diagnostic messages
discard_unused_publics	Discards unused public symbols
dlib_config	Uses the system include files for the DLIB library and determines which configuration of the library to use
-e	Enables language extensions
ec++	Specifies Embedded C++
eec++	Specifies Extended Embedded C++
enable_hardware_workaround	Enables a specific hardware workaround
enable_multibytes	Enables support for multibyte characters in source files
endian	Specifies the byte order of the generated code and data
enum_is_int	Sets the minimum size on enumeration types
error_limit	Specifies the allowed number of errors before compilation stops
-f	Extends the command line
fpu	Selects the type of floating-point unit
guard_calls	Enables guards for function static variable initialization
header_context	Lists all referred source files and header files

Command line option	Description
-I	Specifies include file path
interwork	Generates interworking code
-1	Creates a list file
legacy	Generates object code linkable with older tool chains
lock_regs	Prevents the compiler from using specified registers
macro_positions_in _diagnostics	Obtains positions inside macros in diagnostic messages
mfc	Enables multi-file compilation
misrac	Enables error messages specific to MISRA-C:1998. This option is a synonym ofmisrac1998 and is only available for backwards compatibility.
misrac1998	Enables error messages specific to MISRA-C:1998. See the IAR Embedded Workbench® MISRA C:1998 Reference Guide.
misrac2004	Enables error messages specific to MISRA-C:2004. See the IAR Embedded Workbench® MISRA C:2004 Reference Guide.
misrac_verbose	Enables verbose logging of MISRA C checking. See the IAR Embedded Workbench® MISRA C:1998 Reference Guide or the IAR Embedded Workbench® MISRA C:2004 Reference Guide.
no_clustering	Disables static clustering optimizations
no_code_motion	Disables code motion optimization
no_cse	Disables common subexpression elimination
no_const_align	Disables the alignment optimization for constants.
no_exceptions	Disables C++ exception support
no_fragments	Disables section fragment handling
no_inline	Disables function inlining
no_path_in_file_macros	Removes the path from the return value of the symbolsFILE andBASE_FILE
no_rtti	Disables C++ RTTI support
no_rw_dynamic_init	Disables runtime initialization of static C variables.
no_scheduling	Disables the instruction scheduler

Command line option	Description
no_size_constraints	Relaxes the normal restrictions for code size expansion when optimizing for speed.
no_static_destruction	Disables destruction of C++ static variables at program exit
no_system_include	Disables the automatic search for system include files
no_tbaa	Disables type-based alias analysis
no_typedefs_in_diagnostics	Disables the use of typedef names in diagnostics
no_unaligned_access	Avoids unaligned accesses
no_unroll	Disables loop unrolling
no_warnings	Disables all warnings
no_wrap_diagnostics	Disables wrapping of diagnostic messages
-0	Sets the optimization level
-0	Sets the object filename. Alias foroutput.
only_stdout	Uses standard output only
output	Sets the object filename
predef_macros	Lists the predefined symbols.
preinclude	Includes an include file before reading the source file
preprocess	Generates preprocessor output
public_equ	Defines a global named assembler label
-r	Generates debug information. Alias fordebug.
relaxed_fp	Relaxes the rules for optimizing floating-point expressions
remarks	Enables remarks
require_prototypes	Verifies that functions are declared before they are defined
ropi	Generates code that uses PC-relative references to address code and read-only data.
rwpi	Generates code that uses an offset from the static base register to address-writable data.
section	Changes a section name
separate_cluster_for_ initialized_variables	Separates initialized and non-initialized variables

Command line option	Description
silent	Sets silent operation
strict	Checks for strict compliance with Standard C/C++
system_include_dir	Specifies the path for system include files
thumb	Sets default function mode to Thumb
use_c++_inline	Uses C++ inline semantics in C99
use_unix_directory_	Uses / as directory separator in paths
separators	
vla	Enables C99 VLA support
warnings_affect_exit_code	Warnings affect exit code
warnings_are_errors	Warnings are treated as errors

Descriptions of compiler options

A

The following section gives detailed reference information about each compiler option.

Note that if you use the options page **Extra Options** to specify specific command line options, the IDE does not perform an instant check for consistency problems like conflicting options, duplication of options, or use of irrelevant options.

aapcs		
Syntax	aapcs={std	vfp}
Parameters	std	Processor registers are used for floating-point parameters and return values in function calls according to standard AAPCS. std is the default when the software FPU is selected.
	vfp	VFP registers are used for floating-point parameters and return values. The generated code is not compatible with AEABI code. vfp is the default when a VFP unit is used.
Description	1	n to specify the floating-point calling convention.

--aeabi

Syntax	aeabi
Description	Use this option to generate AEABI-compliant object code. Note that this option must be used together with theguard_calls option.
See also	AEABI compliance, page 192 andguard_calls, page 247.
	To set this option, use Project>Options>C/C++ Compiler>Extra Options .

--align_sp_on_irq

Syntax	align_sp_on_irq
Description	Use this option to align the stack pointer (SP) on entry to $__irq$ declared functions.
	This is especially useful for nested interrupts, where the interrupted code uses the same SP as the interrupt handler. This means that the stack might only have 4-byte alignment, instead of the 8-byte alignment required by AEABI (and some instructions generated by the compiler for some cores).
See also	<i>irq</i> , page 318.
	To set this option, use Project>Options>C/C++ Compiler>Extra Options .

arm		
Syntax	arm	
Description	Use this option to set default function mode to ARM. This setting must be the same for all files included in a program, unless they are interworking.	
	Note: This option has the same effect as thecpu_mode=arm option.	

Y

See also

--interwork, page 248 and __interwork, page 318.

Project>Options>C/C++ Compiler>Code>Processor mode>Arm

--c89

Syntax	c89
Description	Use this option to enable the C89 C dialect instead of Standard C.
	Note: This option is mandatory when the MISRA C checking is enabled.
See also	C language overview, page 161.
2	Project>Options>C/C++ Compiler>Language 1>C dialect>C89

--char_is_signed

Syntax		char_is_signed
Description		By default, the compiler interprets the plain char type as unsigned. Use this option to make the compiler interpret the plain char type as signed instead. This can be useful when you, for example, want to maintain compatibility with another compiler.
		Note: The runtime library is compiled without thechar_is_signed option and cannot be used with code that is compiled with this option.
	X	Project>Options>C/C++ Compiler>Language 2>Plain 'char' is

--char_is_unsigned

Syntax		char_is_unsigned
Description		Use this option to make the compiler interpret the plain char type as unsigned. This is the default interpretation of the plain char type.
	X	Project>Options>C/C++ Compiler>Language 2>Plain 'char' is

--cpu

Syntax

--cpu=*core*

Parameters

core

Specifies a specific processor variant

Description	Use this option to select the processor variant for which the code is to be generated. The default is ARM7TDMI. The following cores and processor macrocells are recognized:		
	ARM7TDMI	ARM10E	Cortex-M0+
	ARM7TDMI-S	ARM1020E	Cortex-M1
	ARM710T	ARM1022E	Cortex-Ms1*
	ARM720T	ARM1026EJ-S	Cortex-M3
	ARM740T	ARM1136J	Cortex-M4
	ARM7EJ-S	ARM1136J-S	Cortex-M4F
	ARM9TDMI	ARM1136JF	Cortex-R4
	ARM920T	ARM1136JF-S	Cortex-R4F
	ARM922T	ARM1176J	Cortex-R5
		(alias for Arm1176JZ)	
	ARM940T	ARM1176J-S	Cortex-R5F
		(alias for ARM1176JZ-S)	
	ARM9E	ARM1176JF	Cortex-R7
		(alias for ARM1176JZF)	
	ARM9E-S	ARM1176JF-S	Cortex-R7F
		(alias for ARM1176JZF-S))
	ARM926EJ-S	Cortex-A5	XScale
	ARM966E-S	Cortex-A7	XScale-IR7
	ARM968E-S	Cortex-A15	
	ARM946E-S	Cortex-M0	

* Cortex-M1 with Operating System extension.

See also

Processor variant, page 48



Project>Options>General Options>Target>Processor configuration

--cpu_mode

Syntax	cpu_mode={arm a	cpu_mode={arm a thumb t}	
Parameters	arm, a(default) thumb, t	Selects the arm mode as the default mode for functions Selects the thumb mode as the default mode for functions	
Description	*	ct the default mode for functions. This setting must be the same a program, unless they are interworking.	
See also		<i>interwork</i> , page 248 and <i>interwork</i> , page 318. Project>Options>General Options>Target>Processor mode	

--c++

Syntax	C++
Description	By default, the language supported by the compiler is C. If you use Standard C++, you must use this option to set the language the compiler uses to C++.
See also	ec++, page 243,eec++, page 243, and Using C++, page 171.,
X	Project>Options>C/C++ Compiler>Language 1>C++
<i>9</i>	and
	Project>Options>C/C++ Compiler>Language 1>C++ dialect>C++

-D

Syntax	-D symbol[=value]	
Parameters	symbol value	The name of the preprocessor symbol The value of the preprocessor symbol
Description		a preprocessor symbol. If no value is specified, 1 is used. This or more times on the command line.

The option -D has the same effect as a #define statement at the top of the source file: -Dsymbol is equivalent to: #define symbol 1 To get the equivalence of: #define FOO specify the = sign but nothing after, for example: -DFOO= Project>Options>C/C++ Compiler>Preprocessor>Defined symbols

debug, -r		
Syntax		debug -r
Description		Use thedebug or -r option to make the compiler include information in the object modules required by the IAR C-SPY® Debugger and other symbolic debuggers.
		Note: Including debug information will make the object files larger than otherwise.
	X	Project>Options>C/C++ Compiler>Output>Generate debug information

--dependencies

Syntax	dependencies	dependencies[=[i m]] {filename directory}	
Parameters	i (default) m	Lists only the names of files Lists in makefile style	
	For information about specifying a filename or a directory, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.		
Description	1	Use this option to make the compiler list the names of all source and header files opened for input into a file with the default filename extension \pm .	

 Example
 If --dependencies or --dependencies=i is used, the name of each opened input file, including the full path, if available, is output on a separate line. For example:

```
c:\iar\product\include\stdio.h
d:\myproject\include\foo.h
```

If --dependencies=m is used, the output is in makefile style. For each input file, one line containing a makefile dependency rule is produced. Each line consists of the name of the object file, a colon, a space, and the name of an input file. For example:

```
foo.o: c:\iar\product\include\stdio.h
foo.o: d:\myproject\include\foo.h
```

An example of using --dependencies with a popular make utility, such as gmake (GNU make):

Set up the rule for compiling files to be something like:

```
%.0 : %.C
    $(ICC) $(ICCFLAGS) $< --dependencies=m $*.d</pre>
```

That is, in addition to producing an object file, the command also produces a dependency file in makefile style (in this example, using the extension .d).

2 Include all the dependency files in the makefile using, for example:

```
-include $(sources:.c=.d)
```

Because of the dash (-) it works the first time, when the .d files do not yet exist.



This option is not available in the IDE.

diag_error		
Syntax	diag_error=	tag[,tag,]
Parameters	tag	The number of a diagnostic message, for example the message number Pe117
Description	violation of the C	reclassify certain diagnostic messages as errors. An error indicates a or C++ language rules, of such severity that object code will not be it code will be non-zero. This option may be used more than once on .
	Project>Options	>C/C++ Compiler>Diagnostics>Treat these as errors

diag_remark			
Syntax	diag_rem	ark=tag[,tag,]	
Parameters	tag	The number of a diagnostic message, for example the message number Pe177	
Description	least severe ty may cause str	Use this option to reclassify certain diagnostic messages as remarks. A remark is the least severe type of diagnostic message and indicates a source code construction that may cause strange behavior in the generated code. This option may be used more than once on the command line.	
	Note: By defa	ault, remarks are not displayed; use theremarks option to display them.	
é	Project>Opt	ions>C/C++ Compiler>Diagnostics>Treat these as remarks	

U	diag_	suppress
----------	-------	----------

Syntax		diag_suppress=tag[,tag,]	
Parameters		tag	The number of a diagnostic message, for example the message number Pe117
Description	×	Use this option to suppress certain diagnostic messages. These messages will not be displayed. This option may be used more than once on the command line. Project>Options>C/C++ Compiler>Diagnostics>Suppress these diagnostics	
diag_warning			
Syntax		diag_warning= <i>tag</i>	[,tag,]

Parameters

tag

The number of a diagnostic message, for example the message number ${\tt Pe826}$

Description Use this option to reclassify certain diagnostic messages as warnings. A warning indicates an error or omission that is of concern, but which will not cause the compiler to stop before compilation is completed. This option may be used more than once on the command line.



Project>Options>C/C++ CompilerLinker>Diagnostics>Treat these as warnings

--diagnostics_tables

Syntax		diagnostics_tables {filename directory}
Parameters		For information about specifying a filename or a directory, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.
Description		Use this option to list all possible diagnostic messages in a named file. This can be convenient, for example, if you have used a pragma directive to suppress or change the severity level of any diagnostic messages, but forgot to document why.
		This option cannot be given together with other options.
	X	This option is not available in the IDE.

--discard_unused_publics

Syntax	discard_unused_publics
Description	Use this option to discard unused public functions and variables when compiling with themfc compiler option.
	Note: Do not use this option only on parts of the application, as necessary symbols might be removed from the generated output.
See also	mfc, page 250 and Multi-file compilation units, page 206.
2	Project>Options>C/C++ Compiler>Discard unused publics

--dlib_config Syntax

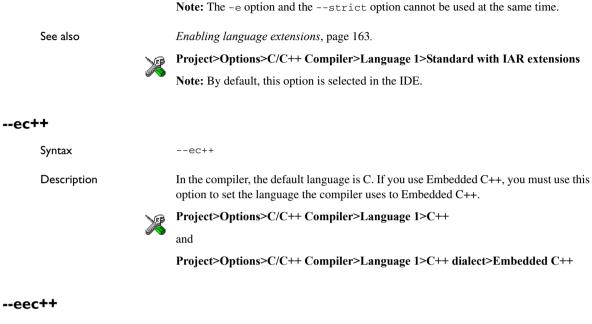
dlib_config	filename.h	config

Parameters		
	filename	A DLIB configuration header file. For information about specifying a filename, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.
	config	The default configuration file for the specified configuration will be used. Choose between:
		none, no configuration will be used
		normal, the normal library configuration will be used (default)
		full, the full library configuration will be used.
Description	explicit file or by specif that library configuratio	fy which library configuration to use, either by specifying an Yying a library configuration—in which case the default file for on will be used. Make sure that you specify a configuration that ry you are using. If you do not specify this option, the default e will be used.
	can find the library obje	raries are delivered with corresponding configuration files. You ext files and the library configuration files in the directory s and information about prebuilt runtime libraries, see <i>Using</i> 101.
	corresponding customiz	ustomized runtime library, you should also create a red library configuration file, which must be specified to the rmation, see <i>Building and using a customized library</i> , page 112.
X	To set related options, choose:	
	Project>Options>Gen	eral Options>Library Configuration
Syntax	-е	

-e

Description

In the command line version of the compiler, language extensions are disabled by default. If you use language extensions such as extended keywords and anonymous structs and unions in your source code, you must use this option to enable them.



Syntax		eec++
Description		In the compiler, the default language is C. If you take advantage of Extended Embedded C++ features like namespaces or the standard template library in your source code, you must use this option to set the language the compiler uses to Extended Embedded C++.
See also		Extended Embedded C++, page 172.
	X	Project>Options>C/C++ Compiler>Language 1>C++
	93	and
		Project>Options>C/C++ Compiler>Language 1>C++ dialect>Extended

Embedded C++

--enable_hardware_workaround

Syntax		enable_hardware_workaround=waid[,waid]	
Parameters		waid	The ID number of the workaround to enable. For a list of available workarounds to enable, see the release notes.
Description		Use this option to make the compiler generate a workaround for a specific hardware problem.	
See also		The release notes for the compiler for a list of available parameters.	
	X	To set this option, use F	Project>Options>C/C++ Compiler>Extra Options.

--enable_multibytes

Syntax		enable_multibytes
Description		By default, multibyte characters cannot be used in C or C++ source code. Use this option to make multibyte characters in the source code be interpreted according to the host computer's default setting for multibyte support.
		Multibyte characters are allowed in C and C++ style comments, in string literals, and in character constants. They are transferred untouched to the generated code.
	X	Project>Options>C/C++ Compiler>Language 2>Enable multibyte support

--endian

Syntax	endian={big b li	ttle 1}
Parameters	big, b little, l(default)	Specifies big-endian as the default byte order Specifies little-endian as the default byte order
Description	1 1	fy the byte order of the generated code and data. By default, the e in little-endian byte order.

 See also
 Byte order, page 49 and Byte order, page 298.

 Project>Options>General Options>Target>Endian mode

--enum_is_int

Syntax		enum_is_int
Description		Use this option to force the size of all enumeration types to be at least 4 bytes.
		Note: This option will not consider the fact that an enum type can be larger than an integer type.
See also		The enum type, page 299.
	X	To set this option, use Project>Options>C/C++ Compiler>Extra Options .

--error_limit

-f

Syntax	error_limit=n		
Parameters	п	The number of errors before the compiler stops the compilation. <i>n</i> must be a positive integer; 0 indicates no limit.	
Description	compiler stops the co	Use theerror_limit option to specify the number of errors allowed before the compiler stops the compilation. By default, 100 errors are allowed. This option is not available in the IDE.	
Syntax	-f filename		
Parameters	For information about directory as parameters	ut specifying a filename, see <i>Rules for specifying a filename or ters</i> , page 228.	

DescriptionDescriptions		Use this option to make the compiler read command line options from the named file, with the default filename extension xcl.		
		In the command file, you format the items exactly as if they were on the command line itself, except that you may use multiple lines, because the newline character acts just as a space or tab character.		
		Both C and C++ style comments are allowed in the file. Double quotes behave in the same way as in the Microsoft Windows command line environment.		



To set this option, use **Project>Options>C/C++ Compiler>Extra Options**.

--fpu

Syntax	fpu={VFPv2 VF	fpu={VFPv2 VFPv3 VFPv3_d16 VFPv4 VFPv4_sp VFP9-S none}	
Parameters	VFPv2	For a system that implements a VFP unit conforming to the architecture VFPv2.	
	VFPv3	For a system that implements a VFP unit conforming to the architecture VFPv3.	
	VFPv3_d16	For a system that implements a VFP unit conforming to the D16 variant of the architecture VFPv3.	
	VFPv4	For a system that implements a VFP unit conforming to the architecture VFPv4.	
	VFPv4_sp	For a system that implements a VFP unit conforming to the single-precision variant of the architecture VFPv4.	
	VFP9-S	VFP9-S is an implementation of the VFPv2 architecture that can be used with the ARM9E family of CPU cores. Selecting the VFP9-S coprocessor is therefore identical to selecting the VFPv2 architecture.	
	none (default)	The software floating-point library is used.	
Description	Floating Point (VFP	enerate code that performs floating-point operations using a Vector coprocessor. By selecting a VFP coprocessor, you will override the floating-point library for all supported floating-point operations.	

	See also		VFP and floating-point arithmetic, page 49.
		X	Project>Options>General Options>Target>FPU
gua	ard_calls		
	Syntax		guard_calls
	Description		Use this option to enable guards for function static variable initialization. This option should be used in a threaded C++ environment.
			Note: This option requires a threaded C++ environment, which is not supported in the IAR C/C++ Compiler for ARM.
	See also		Managing a multithreaded environment, page 130.
		X	To set this option, use Project>Options>C/C++ Compiler>Extra Options .

--header_context

-1

Syntax	header_con	ltext	
Description	that was include	Occasionally, to find the cause of a problem it is necessary to know which header file that was included from which source line. Use this option to list, for each diagnostic message, not only the source position of the problem, but also the entire include stack at that point.	
X	This option is no	ot available in the IDE.	
Syntax	-I path		
Parameters	path	The search path for #include files	
	pacii	The scale pair for #include files	
Description	1	to specify the search paths for #include files. This option can be used on the command line.	

Include file search procedure, page 221.

Project>Options>C/C++ Compiler>Preprocessor>Additional include directories

--interwork

See also

X

Syntax	interwork
Description	Use this option to generate interworking code.
	In code compiled with this option, functions will by default be of the type interwork. It will be possible to mix files compiled as arm and thumb (using thecpu_mode option) as long as they are all compiled with theinterwork option.
	Note: Source code compiled for an ARM architecture v5 or higher, or for AEABI compliance, is interworking by default.
See also	interwork, page 318.
X	Project>Options>General Options>Target>Generate interwork code

-1

Syntax	-l[a A b B c C D][N][H] {filename directory}	
Parameters	a (default)	Assembler list file
	А	Assembler list file with C or C++ source as comments
	b	Basic assembler list file. This file has the same contents as a list file produced with -la, except that no extra compiler-generated information (runtime model attributes, call frame information, frame size information) is included *
	В	Basic assembler list file. This file has the same contents as a list file produced with -1A, except that no extra compiler generated information (runtime model attributes, call frame information, frame size information) is included *
	С	C or C++ list file
	C (default)	C or C++ list file with assembler source as comments

		D	C or C++ list file with assembler source as comments, but without instruction offsets and hexadecimal byte values
		Ν	No diagnostics in file
		н	Include source lines from header files in output. Without this option, only source lines from the primary source file are included
		* This makes the list file by a human.	less useful as input to the assembler, but more useful for reading
		For information about sp filename or directory as	pecifying a filename or a directory, see <i>Rules for specifying a parameters</i> , page 228.
Description			ate an assembler or C/C++ listing to a file. Note that this option e times on the command line.
	VP	To set related options, choose:	
	I A A A A A A A A A A A A A A A A A A A	Project>Options>C/C+	++ Compiler>List
legacy			
legacy Syntax		legacy={RVCT3.0}	
•		legacy={RVCT3.0} RVCT3.0	Generates object code linkable with the linker in RVCT3.0. Use this mode together with theaeabi option to generate code that should be linked with the linker in RVCT3.0.
Syntax		RVCT3.0	Use this mode together with theaeabi option to generate
Syntax Parameters	×	RVCT3.0 Use this option to genera	Use this mode together with theaeabi option to generate code that should be linked with the linker in RVCT3.0.
Syntax Parameters	X	RVCT3.0 Use this option to genera	Use this mode together with theaeabi option to generate code that should be linked with the linker in RVCT3.0. ate code compatible with the specified toolchain.

Parameters

registers

A comma-separated list of register names and register intervals to be locked, in the range R4-R11.

Description	Use this option to prevent the compiler from generating code that uses the specified registers.
Example	lock_regs=R4 lock_regs=R8-R11 lock_regs=R4,R8-R11
	To set this option, use Project>Options>C/C++ Compiler>Extra Options.

--macro_positions_in_diagnostics

Syntax		macro_positions_in_diagnostics
Description		Use this option to obtain position references inside macros in diagnostic messages. This is useful for detecting incorrect source code constructs in macros.
	X	To set this option, use Project>Options>C/C++ Compiler>Extra Options .

--mfc

Syntax		mfc
Description		Use this option to enable <i>multi-file compilation</i> . This means that the compiler compiles one or several source files specified on the command line as one unit, which enhances interprocedural optimizations.
		Note: The compiler will generate one object file per input source code file, where the first object file contains all relevant data and the other ones are empty. If you want only the first file to be produced, use the -0 compiler option and specify a certain output file.
Example		<pre>iccarm myfile1.c myfile2.c myfile3.cmfc</pre>
See also		discard_unused_publics, page 241,output, -o, page 259, and Multi-file compilation units, page 206.
	X	Project>Options>C/C++ Compiler>Multi-file compilation

--no_clustering

	X	Project>Options>C/C++ Compiler>Optimizations>Enable transformations>Static clustering
See also		Static clustering, page 210.
Description		Use this option to disable static clustering optimizations. Note: This option has no effect at optimization levels below Medium.
Syntax		no_clustering

--no_code_motion

X	 Project>Options>C/C++ Compiler>Optimizations>Enable transformations>Code motion
See also	Code motion, page 209.
Description	Use this option to disable code motion optimizations. Note: This option has no effect at optimization levels below Medium.
Syntax	no_code_motion

--no_const_align

Syntax		no_const_align
Description		By default, the compiler uses alignment 4 for objects with a size of 4 bytes or more. Use this option to make the compiler align const objects based on the alignment of their type.
		For example, a string literal will get alignment 1, because it is an array with elements of the type const char which has alignment 1. Using this option might save ROM space, possibly at the expense of processing speed.
See also		Alignment, page 297
	X	To set this option, use Project>Options>C/C++ Compiler>Extra Options.

--no_cse

X	Project>Options>C/C++ Compiler>Optimizations>Enable transformations>Common subexpression elimination
See also	Common subexpression elimination, page 208.
Description	Use this option to disable common subexpression elimination. Note: This option has no effect at optimization levels below Medium.
Syntax	no_cse

--no_exceptions

Syntax		no_exceptions
Description		Use this option to disable exception support in the C++ language. Exception statements like throw and try-catch, and exception specifications on function definitions will generate an error message. Exception specifications on function declarations are ignored. The option is only valid when used together with the $c++$ compiler option. If exceptions are not used in your application, it is recommended to disable support for them by using this option, because exceptions cause a rather large increase in code size.
See also	X	Exception handling, page 174. Project>Options>C/C++ Compiler>Language 1>C++ and Project>Options>C/C++ Compiler>Language 1>C++ dialect>C++>With exceptions
6		

--no_fragments

Syntax	no_fragments
Description	Use this option to disable section fragment handling. Normally, the toolset uses IAR proprietary information for transferring section fragment information to the linker. The linker uses this information to remove unused code and data, and thus further minimize the size of the executable image.

 See also
 Keeping symbols and sections, page 89.

 To set this option, use Project>Options>C/C++ CompilerLinker>Extra Options

--no_inline

	X	Project>Options>C/C++ Compiler>Optimizations>Enable transformations>Function inlining
See also		Inlining functions, page 65.
Description		Use this option to disable function inlining.
Syntax		no_inline

--no_loop_align

Syntax		no_loop_align
Description		Use this option to disable the 4-byte alignment of labels in loops. This option is only useful in Thumb2 mode.
		In ARM/Thumb1 mode, this option is enabled but does not perform anything.
See also		Alignment, page 297
	X	To set this option, use Project>Options>C/C++ Compiler>Extra Options.

--no_mem_idioms

Syntax		no_mem_idioms
Description		Use this option to make the compiler not optimize code sequences that clear, set, or copy a memory region. These memory access patterns (idioms) can otherwise be aggressively optimized, in some cases using calls to the runtime library. In principle, the transformation can involve more than a library call.
	X	To set this option, use Project>Options>C/C++ Compiler>Extra Options.

--no_path_in_file_macros

Syntax		no_path_in_file_macros
Description		Use this option to exclude the path from the return value of the predefined preprocessor symbolsFILE andBASE_FILE
See also		Description of predefined preprocessor symbols, page 382.
	X	This option is not available in the IDE.
no_rtti		

Syntax --no_rtti

Description		Use this option to disable the runtime type information (RTTI) support in the C++ language. RTTI statements like dynamic_cast and typeid will generate an error message. This option is only valid when used together with thec++ compiler option.
See also		Using C++, page 171.
	X	Project>Options>C/C++ Compiler>Language 1>C++
	I.	and

Project>Options>C/C++ Compiler>Language 1>C++ dialect>C++>With RTTI

--no_rw_dynamic_init

Syntax	no_rw_dynamic_init
Description	Use this option to disable runtime initialization of static C variables.
	C source code that is compiled withropi orrwpi cannot have static pointer variables and constants initialized to addresses of objects that do not have a known address at link time. To solve this for writable static variables, the compiler generates code that performs the initialization at program startup (in the same way as dynamic initialization in C++).
See also	ropi, page 262 andrwpi, page 263
y	Project>Options>C/C++ Compiler>Code>No dynamic read/write/initialization

--no_scheduling

	X	Project>Options>C/C++ Compiler>Optimizations>Enable transformations>Instruction scheduling
See also		Instruction scheduling, page 210.
Description		Use this option to disable the instruction scheduler. Note: This option has no effect at optimization levels below High.
Syntax		no_scheduling

--no_size_constraints

Syntax	no_size_constraints
Description	Use this option to optimize for speed but relax the normal restrictions for code size expansion.
	Note: This option has no effect unless used with -Ohs.
See also	Speed versus size, page 208.
	Project>Options>C/C++ Compiler>Optimizations>Level>No size constraints

--no_static_destruction

Syntax		no_static_destruction
Description		Normally, the compiler emits code to destroy C++ static variables that require destruction at program exit. Sometimes, such destruction is not needed. Use this option to suppress the emission of such code.
See also	X	Setting up the atexit limit, page 90. To set this option, use Project>Options>C/C++ Compiler>Extra Options .

--no_system_include

Syntax	no_system_include
Description	By default, the compiler automatically locates the system include files. Use this option to disable the automatic search for system include files. In this case, you might need to set up the search path by using the -I compiler option.
See also	 dlib_config, page 242, andsystem_include_dir, page 265. Project>Options>C/C++ Compiler>Preprocessor>Ignore standard include directories

--no_tbaa

Syntax		no_tbaa
Description		Use this option to disable type-based alias analysis. Note: This option has no effect at optimization levels below High.
See also		Type-based alias analysis, page 209.
	X	Project>Options>C/C++ Compiler>Optimizations>Enable transformations>Type-based alias analysis

--no_typedefs_in_diagnostics

Syntax	no_typedefs_in_diagnostics
Description	Use this option to disable the use of typedef names in diagnostics. Normally, when a type is mentioned in a message from the compiler, most commonly in a diagnostic message of some kind, the typedef names that were used in the original declaration are used whenever they make the resulting text shorter.
Example	<pre>typedef int (*MyPtr)(char const *); MyPtr p = "foo"; will give an error message like this:</pre>
	Error[Pe144]: a value of type "char *" cannot be used to initialize an entity of type "MyPtr"

If the --no_typedefs_in_diagnostics option is used, the error message will be like this:

Error[Pe144]: a value of type "char *" cannot be used to initialize an entity of type "int (*)(char const *)"



To set this option, use **Project>Options>C/C++ Compiler>Extra Options**.

--no_unaligned_access

Syntax		no_unaligned_access
Description		Use this option to make the compiler avoid unaligned accesses. Data accesses are usually performed aligned for improved performance. However, some accesses, most notably when reading from or writing to packed data structures, might be unaligned. When using this option, all such accesses will be performed using a smaller data size to avoid any unaligned accesses. This option is only useful for ARMv6 architectures and higher.
See also	10	<i>Alignment</i> , page 297. To set this option, use Project>Options>C/C++ Compiler>Extra Options .
	X	

--no_unroll

	X	Project>Options>C/C++ Compiler>Optimizations>Enable transformations>Loop unrolling
See also		Loop unrolling, page 208.
Description		Use this option to disable loop unrolling. Note: This option has no effect at optimization levels below High.
Syntax		no_unroll

--no_warnings

Syntax

Description

By default, the compiler issues warning messages. Use this option to disable all warning messages.

This .

--no_warnings

This option is not available in the IDE.

--no_wrap_diagnostics

Syntax	no_wrap_diagnostics
Description	By default, long lines in diagnostic messages are broken into several lines to make the message easier to read. Use this option to disable line wrapping of diagnostic messages.
	This option is not available in the IDE.

-0

Syntax	-O[n l m h hs hz]	
Parameters	n	None* (Best debug support)
	1 (default)	Low*
	m	Medium
	h	High, balanced
	hs	High, favoring speed
	hz	High, favoring size
	*The most important di variables will live durin	fference between None and Low is that at None, all non-static g their entire scope.
Description	the code. If no optimiza	e optimization level to be used by the compiler when optimizing tion option is specified, the optimization level Low is used by ed without any parameter, the optimization level High balanced

A low level of optimization makes it relatively easy to follow the program flow in the debugger, and, conversely, a high level of optimization makes it relatively hard.

See also

Controlling compiler optimizations, page 205.

Project>Options>C/C++ Compiler>Optimizations

on	y_stdout		
	Syntax		only_stdout
	Description		Use this option to make the compiler use the standard output stream (stdout) also for messages that are normally directed to the error output stream (stderr).
		X	This option is not available in the IDE.
out	tput, -o		
	Syntax		output {filename directory} -o {filename directory}
	Parameters		For information about specifying a filename or a directory, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.
	Description		By default, the object code output produced by the compiler is located in a file with the same name as the source file, but with the extension o. Use this option to explicitly specify a different output filename for the object code output.
		X	This option is not available in the IDE.

--predef_macros

Syntax	predef_macros {filename directory}
Parameters	For information about specifying a filename, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.
Description	Use this option to list the predefined symbols. When using this option, make sure to also use the same options as for the rest of your project.

-

-

If a filename is specified, the compiler stores the output in that file. If a directory is specified, the compiler stores the output in that directory, in a file with the predef filename extension.

Note that this option requires that you specify a source file on the command line.

Generate #line directives

For information about specifying a filename or a directory, see Rules for specifying a

Project>Options>C/C++ Compiler>Preprocessor>Preprocessor output to file



1

This option is not available in the IDE.

pre	einclude		
	Syntax	preinclude includ	defile
	Parameters	For information about sp directory as parameters,	ecifying a filename, see <i>Rules for specifying a filename or</i> page 228.
	Description	read the source file. This	the compiler read the specified include file before it starts to is useful if you want to change something in the source code , for instance if you want to define a new symbol.
	X	Project>Options>C/C+	+ Compiler>Preprocessor>Preinclude file
pre	eprocess		
	Syntax	preprocess[=[c][1	n][1]] {filename directory}
	Parameters		
		С	Preserve comments
		n	Preprocess only

filename or directory as parameters, page 228.

Use this option to generate preprocessed output to a named file.

Description

--public_equ

	Syntax		public_equ symbol	l[=value]
	Parameters		symbol value	The name of the assembler symbol to be defined An optional value of the defined assembler symbol
	Description			to defining a label in assembler language using the EQU using the PUBLIC directive. This option can be used more than ne.
		X	This option is not availab	ble in the IDE.
rela	axed_fp			

Syntax	relaxed_fp
Description	Use this option to allow the compiler to relax the language rules and perform more aggressive optimization of floating-point expressions. This option improves performance for floating-point expressions that fulfill these conditions:
	• The expression consists of both single- and double-precision values
	• The double-precision values can be converted to single precision without loss of accuracy
	• The result of the expression is converted to single precision.
	Note that performing the calculation in single precision instead of double precision might cause a loss of accuracy.
Example	<pre>float F(float a, float b) { return a + b * 3.0; }</pre>
	The C standard states that 3.0 in this example has the type double and therefore the whole expression should be evaluated in double precision. However, when therelaxed_fp option is used, 3.0 will be converted to float and the whole expression can be evaluated in float precision.



To set related options, choose:

Project>Options>C/C++ Compiler>Language 2>Floating-point semantics

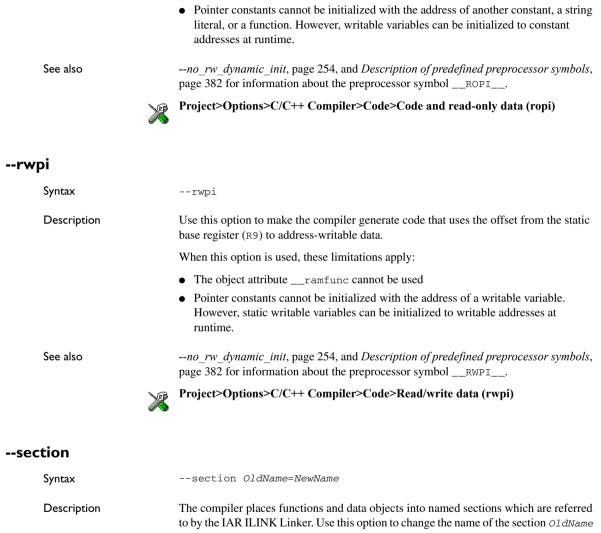
--remarks

Syntax		remarks
Description		The least severe diagnostic messages are called remarks. A remark indicates a source code construct that may cause strange behavior in the generated code. By default, the compiler does not generate remarks. Use this option to make the compiler generate remarks.
See also		Severity levels, page 225.
	X	Project>Options>C/C++ Compiler>Diagnostics>Enable remarks

--require_prototypes

Syntax	require_prototypes
Description	Use this option to force the compiler to verify that all functions have proper prototypes. Using this option means that code containing any of the following will generate an error:
	• A function call of a function with no declaration, or with a Kernighan & Ritchie C declaration
	• A function definition of a public function with no previous prototype declaration
	• An indirect function call through a function pointer with a type that does not include a prototype.
	Project>Options>C/C++ Compiler>Language 1>Require prototypes
ropi	
Syntax	ropi
Description	Use this option to make the compiler generate code that uses PC-relative references to address code and read-only data.
	When this option is used, these limitations apply:
	• C++ constructions cannot be used

• The object attribute __ramfunc cannot be used



to NewName.

This is useful if you want to place your code or data in different address ranges and you find the @ notation, alternatively the #pragma location directive, insufficient. Note that any changes to the section names require corresponding modifications in the linker configuration file.

Example	To place functions in the section MyText, use:
	section .text=MyText
See also	Controlling data and function placement in memory, page 201. Project>Options>C/C++ Compiler>Output>Code section name

--separate_cluster_for_initialized_variables

Syntax	separate_cluster_for_initialized_variables
Description	Use this option to separate initialized and non-initialized variables when using variable clustering. This might reduce the number of bytes in the ROM area which are needed for data initialization, but it might lead to larger code.
	This option can be useful if you want to have your own data initialization routine, but want the IAR tools to arrange for the zero-initialized variables.
See also	Manual initialization, page 91 and Initialize directive, page 408.
	To set this option, use Project>Options>C/C++ Compiler>Extra Options.

--silent

Syntax		silent
Description		By default, the compiler issues introductory messages and a final statistics report. Use this option to make the compiler operate without sending these messages to the standard output stream (normally the screen).
		This option does not affect the display of error and warning messages.
	X	This option is not available in the IDE.

--strict

Syntax	strict	
Description	By default, the compiler accepts a relaxed superse option to ensure that the source code of your appli Standard C and C++.	
	Note: The $-e$ option and the $strict$ option ca	nnot be used at the same time.
See also	Enabling language extensions, page 163.	
	Project>Options>C/C++ Compiler>Language	l>Language conformance>Strict

--system_include_dir

Syntax	system_include_d	system_include_dir path	
Parameters	path	The path to the system include files. For information about specifying a path, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.	
Description	to explicitly specify a di	By default, the compiler automatically locates the system include files. Use this option to explicitly specify a different path to the system include files. This might be useful if you have not installed IAR Embedded Workbench in the default location.	
See also		dlib_config, page 242, andno_system_include, page 256.	
	This option is not availa	ble in the IDE.	

--thumb

Syntax	thumb
Description	Use this option to set default function mode to Thumb. This setting must be the same for all files included in a program, unless they are interworking.
	Note: This option has the same effect as thecpu_mode=thumb option.

See also --interwork, page 248 and __interwork, page 318.
Project>Options>C/C++ Compiler>Code>Processor mode>Thumb
--use_c++_inline

Syntax	use_c++_inline	
Description	Standard C uses slightly different semantics for the inla Use this option if you want C++ semantics when you are	5
See also	Inlining functions, page 65	
	 Project>Options>C/C++ Compiler>Language 1>C d semantics 	ialect>C99>C++ inline

--use_unix_directory_separators

Syntax	use_unix_directory_separators
Description	Use this option to make DWARF debug information use / (instead of $\)$ as directory separators in file paths.
	This option can be useful if you have a debugger that requires directory separators in UNIX style.
é	To set this option, use Project>Options>C/C++ Compiler>Extra Options .

--vla

Syntax	vla
Description	Use this option to enable support for C99 variable length arrays. Such arrays are located on the heap. This option requires Standard C and cannot be used together with the $c89$ compiler option.
	Note: vla should not be used together with the longjmp library function, as that can lead to memory leakages.

See also

C language overview, page 161.

Project>Options>C/C++ Compiler>Language 1>C dialect>Allow VLA

--warnings_affect_exit_code

Syntax

--warnings_affect_exit_code

Description

By default, the exit code is not affected by warnings, because only errors produce a non-zero exit code. With this option, warnings will also generate a non-zero exit code.

This option is not available in the IDE.

--warnings_are_errors

Syntax	warnings_are_errors
Description	Use this option to make the compiler treat all warnings as errors. If the compiler encounters an error, no object code is generated. Warnings that have been changed into remarks are not treated as errors.
	Note: Any diagnostic messages that have been reclassified as warnings by the optiondiag_warning or the #pragma diag_warning directive will also be treated as errors whenwarnings_are_errors is used.
See also	diag_warning, page 240.
	Project>Options>C/C++ Compiler>Diagnostics>Treat all warnings as errors

Descriptions of compiler options

Linker options

This chapter gives detailed reference information about each linker option.

For general syntax rules, see Options syntax, page 227.

Summary of linker options

This table summarizes the linker options:

Command line option	Description
basic_heap	Uses a basic heap instead of an advanced heap
BE8	Uses the big-endian format BE8
BE32	Uses the big-endian format BE32
call_graph	Produces a call graph file in XML format
config	Specifies the linker configuration file to be used by the linker
config_def	Defines symbols for the configuration file
cpp_init_routine	Specifies a user-defined C++ dynamic initialization routine
cpu	Specifies a processor variant
debug_lib	Uses the C-SPY debug library
define_symbol	Defines symbols that can be used by the application
dependencies	Lists file dependencies
diag_error	Treats these message tags as errors
diag_remark	Treats these message tags as remarks
diag_suppress	Suppresses these diagnostic messages
diag_warning	Treats these message tags as warnings
diagnostics_tables	Lists all diagnostic messages
enable_hardware_workaround	Enables specified hardware workaround
enable_stack_usage	Enables stack usage analysis
entry	Treats the symbol as a root symbol and as the start of the application

Table 31: Linker options summary

Command line option	Description
error_limit	Specifies the allowed number of errors before linking stops
exception_tables	Generates exception tables for C code
export_builtin_config	Produces an ${\tt icf}$ file for the default configuration
extra_init	Specifies an extra initialization routine that will be called if it is defined.
-f	Extends the command line
force_exceptions	Always includes exception runtime code
force_output	Produces an output file even if errors occurred
image_input	Puts an image file in a section
inline	Inlines small routines
keep	Forces a symbol to be included in the application
log	Enables log output for selected topics
log_file	Directs the log to a file
mangled_names_in_messages	Adds mangled names in messages
map	Produces a map file
merge_duplicate_sections	Merges equivalent read-only sections
misrac	Enables error messages specific to MISRA-C: 1998. This option is a synonym tomisrac1998 and is only available for backwards compatibility.
misrac1998	Enables error messages specific to MISRA-C:1998. See the IAR Embedded Workbench® MISRA C:1998 Reference Guide.
misrac2004	Enables error messages specific to MISRA-C:2004. See the IAR Embedded Workbench® MISRA C:2004 Reference Guide.
misrac_verbose	Enables verbose logging of MISRA C checking. See the IAR Embedded Workbench® MISRA C:1998 Reference Guide and the IAR Embedded Workbench® MISRA C:2004 Reference Guide.
no_dynamic_rtti_elimination	Includes dynamic runtime type information even when it is not needed.
no_exceptions	Generates an error if exceptions are used
no_fragments	Disables section fragment handling

Table 31: Linker options summary (Continued)

Command line option	Description
no_library_search	Disables automatic runtime library search
no_locals	Removes local symbols from the ELF executable image.
no_range_reservations	Disables range reservations for absolute symbols
no_remove	Disables removal of unused sections
no_veneers	Disables generation of veneers
no_vfe	Disables Virtual Function Elimination
no_warnings	Disables generation of warnings
no_wrap_diagnostics	Does not wrap long lines in diagnostic messages
-0	Sets the object filename. Alias foroutput.
only_stdout	Uses standard output only
output	Sets the object filename
pi_veneers	Generates position independent veneers.
place_holder	Reserve a place in ROM to be filled by some other tool, for example a checksum calculated by ielftool.
redirect	Redirects a reference to a symbol to another symbol
remarks	Enables remarks
search	Specifies more directories to search for object and library files
semihosting	Links with debug interface
silent	Sets silent operation
stack_usage_control	Specifies a stack usage control file
strip	Removes debug information from the executable image
vfe	Controls Virtual Function Elimination
warnings_affect_exit_code	Warnings affects exit code
warnings_are_errors	Warnings are treated as errors

Table 31: Linker options summary (Continued)

Descriptions of linker options

The following section gives detailed reference information about each compiler and linker option.



Note that if you use the options page **Extra Options** to specify specific command line options, the IDE does not perform an instant check for consistency problems like conflicting options, duplication of options, or use of irrelevant options.

h	asic	h۵	an
v	asic	110	ap

Syntax

Description

--basic_heap

Use this option to use a basic heap instead of an advanced heap. The basic heap introduces less overhead and is suitable for example in applications that only allocate heap memory, but never calls free. For any other use of the heap, the advanced heap is usually more efficient.

To set this option, use **Project>Options>Linker>Extra Options**.

--BE8

Syntax	BE8
Description	Use this option to specify the Byte Invariant Addressing mode.
	This means that the linker reverses the byte order of the instructions, resulting in little-endian code and big-endian data. This is the default byte addressing mode for ARMv6 big-endian images. This is the only mode available for ARM v6M and ARM v7 with big-endian images.
	Byte Invariant Addressing mode is only available on ARM processors that support ARMv6, ARM v6M, and ARM v7.
See also	Byte order, page 49, Byte order, page 298,BE32, page 272, andendian, page 244.
X	Project>Options>General Options>Target>Endian mode

--BE32

Syntax

--BE32

Description Use this option to specify the legacy big-endian mode.

This produces big-endian code and data. This is the only byte-addressing mode for all big-endian images prior to ARMv6. This mode is also available for ARM v6 with big-endian, but not for ARM v6M or ARM v7.

See also

Byte order, page 49, Byte order, page 298, --BE8, page 272, and --endian, page 244.



--call_graph

Syntax	call_graph {filename directory}
Parameters	For information about specifying a filename or directory, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.
Description	Use this option to produce a call graph file. If no filename extension is specified, the extension cgx is used. This option can only be used once on the command line.
	Using this option enables stack usage analysis in the linker.
See also	Stack usage analysis, page 78
	Project>Options>Linker>Advanced>Call graph output (XML)

--config

Syntax		config filename
Parameters		For information about specifying a filename, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.
Description		Use this option to specify the configuration file to be used by the linker (the default filename extension is icf). If no configuration file is specified, a default configuration is used. This option can only be used once on the command line.
See also		The chapter The linker configuration file.
	X	Project>Options>Linker>Config>Linker configuration file

--config_def

Syntax	config_def symbo	config_def symbol[=constant_value]	
Parameters	symbol	symbolThe name of the symbol to be used in the configuration file.By default, the value 0 (zero) is used.	
	constant_value	The constant value of the configuration symbol.	
Description	file. This option has the	Use this option to define a constant configuration symbol to be used in the configuration file. This option has the same effect as the define symbol directive in the linker configuration file. This option can be used more that once on the command line.	
See also	define_symbol, page 2	define_symbol, page 275 and Interaction between ILINK and the application, page 94.	
2	Project>Options>Link	Project>Options>Linker>Config>Defined symbols for configuration file	

--cpp_init_routine

Syntax	cpp_init_routine routine	
Parameters	routine A user-defined C++ dynamic initialization routine.	
Description	When using the IAR C/C++ compiler and the standard library, C++ dynamic initialization is handled automatically. In other cases you might need to use this option.	
	If any sections with the section type INIT_ARRAY or PREINIT_ARRAY are included in your application, the C++ dynamic initialization routine is considered to be needed. By default, this routine is namediar_cstart_call_ctors and is called by the startup code in the standard library. Use this option if you are not using the standard library and require another routine to handle these section types.	
	To set this option, use Project>Options>Linker>Extra Options .	

--cpu

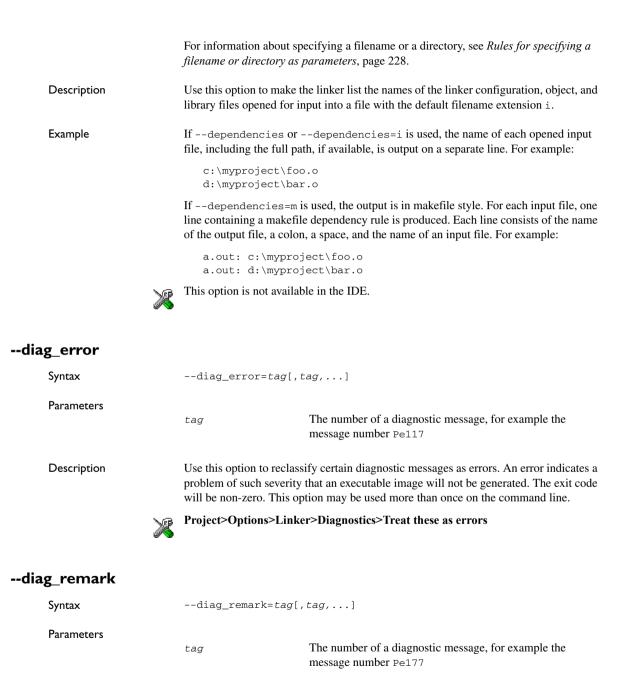
Syntax		cpu= <i>core</i>	
Parameters		core Specific	es a specific processor variant
Description		Use this option to select the processor variant to link your application for. The default is to use a processor or architecture compatible with the object file attributes.	
See also		cpu, page 235 for a list of recognized cores and processor macrocells.	
	X	Project>Options>General Options>Target>Processor configuration	

--define_symbol

Syntax	define_symbol sy	define_symbol symbol=constant_value	
Parameters	symbol	<i>symbol</i> The name of the constant symbol that can be used by the application.	
	constant_value	The constant value of the symbol.	
Description	Use this option to define a constant symbol, that is a label, that can be used by your application. This option can be used more than once on the command line. Note that his option is different from the define symbol directive.		
See also	config_def, page 274 and Interaction between ILINK and the application, page 94.		
) e	Project>Options>Linker>#define>Defined symbols		

--dependencies

Syntax	dependencies[=[i m]] {filename directory}	
Parameters	i (default) m	Lists only the names of files Lists in makefile style



Description Use this option to reclassify certain diagnostic messages as remarks. A remark is the least severe type of diagnostic message and indicates a construction that may cause strange behavior in the executable image. This option may be used more than once on the command line.

Note: By default, remarks are not displayed; use the --remarks option to display them.



diag_suppress		
Syntax	diag_sup	press=tag[,tag,]
Parameters	tag	The number of a diagnostic message, for example the message number Pe117
Description	1	on to suppress certain diagnostic messages. These messages will not be is option may be used more than once on the command line.
	Project>Opt	ions>Linker>Diagnostics>Suppress these diagnostics

-

diag_warning			
Syntax		<pre>diag_warning=tag[,tag,]</pre>	
Parameters	t	ag	The number of a diagnostic message, for example the message number Pe826
Description	in st	Use this option to reclassify certain diagnostic messages as warnings. A warning indicates an error or omission that is of concern, but which will not cause the linker to stop before linking is completed. This option may be used more than once on the command line.	
	<u>у</u> Р	roject>Options>C/	C++ CompilerLinker>Diagnostics>Treat these as warnings

--diagnostics_tables

Syntax	diagnostics_tables {filename directory}
Parameters	For information about specifying a filename or a directory, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.
Description	Use this option to list all possible diagnostic messages in a named file. This can be convenient, for example, if you have used a pragma directive to suppress or change the severity level of any diagnostic messages, but forgot to document why.
	This option cannot be given together with other options.
	This option is not available in the IDE.

--enable_hardware_workaround

Syntax	enable_hardware_workaround=waid[waid[]]	
Parameters	waid The ID number of the workaround that you want to enable. For a list of available workarounds, see the releanotes available in the Information Center.	ise
Description	Use this option to make the linker generate a workaround for a specific hardware problem.	
See also	The release notes for the linker for a list of available parameters.	
	To set this option, use Project>Options>Linker>Extra Options .	

--enable_stack_usage

Syntax	enable_stack_usage
Description	Use this option to enable stack usage analysis. If a linker map file is produced, a stack usage chapter is included in the map file.
	Note: If you use at least one of thestack_usage_control orcall_graph options, stack usage analysis is automatically enabled.

Project>Options>Linker>Advanced>Enable stack usage analysis --entry Syntax --entry symbol **Parameters** symbol The name of the symbol to be treated as a root symbol and start label Description Use this option to make a symbol be treated as a root symbol and the start label of the application. This is useful for loaders. If this option is not used, the default start symbol is __iar_program_start. A root symbol is kept whether or not it is referenced from the rest of the application, provided its module is included. A module in an object file is always included and a module part of a library is only included if needed. Project>Options>Linker>Library>Override default program entry --error limit --error_limit=n Syntax **Parameters** The number of errors before the linker stops linking. n must п be a positive integer; 0 indicates no limit. Description Use the --error_limit option to specify the number of errors allowed before the

Stack usage analysis, page 78



See also

This option is not available in the IDE.

linker stops the linking. By default, 100 errors are allowed.

--exception_tables

Syntax	exception_table	es={nocreate unwind cantunwind}
Parameters	nocreate (default)	Does not generate entries. Uses the least amount of memory, but the result is undefined if an exception is propagated through a function without exception information.
	unwind	Generates unwind entries that enable an exception to be correctly propagated through functions without exception information.
	cantunwind	Generates no-unwind entries so that any attempt to propagate an exception through the function will result in a call to terminate.
Description		ermine what the linker should do with functions that do not have but which do have correct call frame information.
	-	that C functions get correct call frame information. For functions anguage you need to use assembler directives to generate call
See also	Using C++, page 171	
1	To set this option, use	Project>Options>Linker>Extra Options.

--export_builtin_config

Syntax		export_builtin_config filename
Parameters		For information about specifying a filename, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.
Description		Exports the configuration used by default to a file.
	X	This option is not available in the IDE.

--extra_init

Syntax		extra_init routine	
Parameters		routine A user-define	ed initialization routine.
Description		Use this option to make the linker add an entry for the specified routine at the end of the initialization table. The routine will be called during system startup, after other initialization routines have been called and before main is called. Note that the routine must preserve the value passed to it in register R0. No entry is added if the routine is not defined.	
	X	To set this option, use Project>Option	s>Linker>Extra Options.

-f

Syntax		-f filename
Parameters		For information about specifying a filename, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.
Description		Use this option to make the linker read command line options from the named file, with the default filename extension xcl.
		In the command file, you format the items exactly as if they were on the command line itself, except that you may use multiple lines, because the newline character acts just as a space or tab character.
		Both C and C++ style comments are allowed in the file. Double quotes behave in the same way as in the Microsoft Windows command line environment.
	X	To set this option, use Project>Options>Linker>Extra Options.

--force_exceptions

Syntax	force_exceptions
Description	Use this option to make the linker include exception tables and exception code even when the linker heuristics indicate that exceptions are not used.

The linker considers exceptions to be used if there is a throw expression that is not a rethrow in the included code. If there is no such throw expression in the rest of the code, the linker arranges for operator new, dynamic_cast, and typeid to call abort instead of throwing an exception on failure. If you need to catch exceptions from these constructs but your code contains no other throws, you might need to use this option.

See also

Using C++, page 171.



Project>Options>Linker>Optimizations>C++ Exceptions>Allow>Always include

--force_output

Syntax	force_output		
Description	Use this option to produce an output executable image regardless of any linking errors.	Us	s.
	To set this option, use Project>Options>Linker>Extra Options	70 To	

--image_input

Syntax	image_input	<pre>filename [,symbol,[section[,alignment]]]</pre>
Parameters	filename	The pure binary file containing the raw image you want to link
	symbol	The symbol which the binary data can be referenced with.
	section	The section where the binary data will be placed; default is .text.
	alignment	The alignment of the section; default is 1.
Description	1	ink pure binary files in addition to the ordinary input files. The file's placed in the section, which means it can only contain pure binary

		Note: Just as for sections from object files, sections created by using theimage_input option are not included unless actually needed. You can either specify a symbol in the option and reference this symbol in your application (or by use of akeep option), or you can specify a section name and use the keep directive in a linker configuration file to ensure that the section is included.
Example		image_input bootstrap.abs, Bootstrap,CSTARTUPCODE,4 The contents of the pure binary file bootstrap.abs are placed in the section CSTARTUPCODE. The section where the contents are placed is 4-byte aligned and will only be included if your application (or the command line optionkeep) includes a reference to the symbol Bootstrap.
See also	X	keep, page 283. Project>Options>Linker>Input>Raw binary image

--inline

Syntax		inline
Description		Some routines are so small that they can fit in the space of the instruction that calls the routine. Use this option to make the linker replace the call of a routine with the body of the routine, where applicable.
	X	Project>Options>Linker>Optimizations>Inline small routines

--keep

Description	3 ,	er keeps a symbol only if it is needed by your application. Use this ymbol always be included in the final application.
·	option to make a s	ymbol always be included in the final application.
X	Project>Options>	Linker>Input>Keep symbols

--log

Syntax

Parameters

log	<pre>topic[,topic,]</pre>
topic	can be one of:

initialization	Lists copy batches and the compression selected for each batch.
libraries	Lists all decisions taken by the automatic library selector. This might include extra symbols needed (keep), redirections (redirect), as well as which runtime libraries that were selected.
modules	Lists each module that is selected for inclusion in the application, and which symbol that caused it to be included.
redirects	Lists redirected symbols.
sections	Lists each symbol and section fragment that is selected for inclusion in the application, and the dependence that caused it to be included.
unused_fragments	Lists those section fragments that were not included in the application.
veeners	Lists some veneer creation and usage statistics.
-	he linker log information to stdout. The log information can ng why an executable image became the way it is.

See also

Description

--log_file, page 284.

Project>Options>Linker>List>Generate log

--log_file

Syntax	log_file filename
Parameters	For information about specifying a filename, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.
Description	Use this option to direct the log output to the specified file.

See also

--log, page 284.

Project>Options>Linker>List>Generate log

--mangled_names_in_messages

Syntax

Description

--mangled names in messages

Use this option to produce both mangled and unmangled names for C/C++ symbols in messages. Mangling is a technique used for mapping a complex C name or a C++ name (for example, for overloading) into a simple name. For example, void h(int, char) becomes _Zlhic.

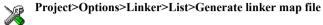


This option is not available in the IDE.

--map

--map {filename|directory} Syntax Description Use this option to produce a linker memory map file. The map file has the default filename extension map. The map file contains: • Linking summary in the map file header which lists the version of the linker, the current date and time, and the command line that was used. • Runtime attribute summary which lists runtime attributes. Placement summary which lists each section/block in address order, sorted by placement directives. Initialization table layout which lists the data ranges, packing methods, and compression ratios. • Module summary which lists contributions from each module to the image, sorted by directory and library. • Entry list which lists all public and some local symbols in alphabetical order, indicating which module they came from. • Some of the bytes might be reported as *shared*. Shared objects are functions or data objects that are shared between modules. If any of these occur in more than one module, only one copy is retained. For example, in some cases inline functions are not inlined, which means that they are marked as shared, because only one instance of each function will be included in the final application. This mechanism is sometimes also used for compiler-generated code or data not directly associated with a particular function or variable, and when only one instance is required in the final application.

This option can only be used once on the command line.



--merge_duplicate_sections

×	Project>Options>Linker>Optimizations>Merge duplicate sections
Description	Use this option to keep only one copy of equivalent read-only sections. Note that this can cause different functions or constants to have the same address, so an application that depends on the addresses being different will not work correctly with this option enabled.
Syntax	merge_duplicate_sections

--no_dynamic_rtti_elimination

Syntax		no_dynamic_rtti_elimination
Description		Use this option to make the linker include dynamic (polymorphic) runtime type information (RTTI) data in the application image even when the linker heuristics indicate that it is not needed.
		The linker considers dynamic runtime type information to be needed if there is a typeid or dynamic_cast expression for a polymorphic type in the included code. By default, if the linker detects no such expression, RTTI data will not be included just to make dynamic RTTI requests work.
		Note: A typeid expression for a <i>non</i> -polymorphic type results in a direct reference to a particular RTTI object and will not cause the linker to include any potentially unneeded objects.
See also		Using C++, page 171.
	X	To set this option, use Project>Options>Linker>Extra Options.

--no_exceptions

Syntax		no_exceptions
Description		Use this option to make the linker generate an error if there is a throw in the included code. This option is useful for making sure that your application does not use exceptions.
See also		Using C++, page 171.
	X	To set related options, choose: Project>Options>Linker>Optimizations>Allow C++ exceptions
		rejet options Limite optimizations finds of electrons

--no_fragments

Syntax		no_fragments
Description		Use this option to disable section fragment handling. Normally, the toolset uses IAR proprietary information for transferring section fragment information to the linker. The linker uses this information to remove unused code and data, and thus further minimize the size of the executable image. The effect of using this option in the compiler is smaller object size.
See also		Keeping symbols and sections, page 89.
	X	To set this option, use Project>Options>C/C++ CompilerLinker>Extra Options

--no_library_search

Syntax	no_library_search
Description	Use this option to disable the automatic runtime library search. This option turns off the automatic inclusion of the correct standard libraries. This is useful, for example, if the application needs a user-built standard library, etc.
	Note that the option disables all steps of the automatic library selection, some of which might need to be reproduced if you are using the standard libraries. Use thelog libraries linker option together with automatic library selection enabled to determine which the steps are.
2	Project>Options>Linker>Library>Automatic runtime library selection

--no_locals

<u> </u>			
	Syntax		no_locals
	Description		Use this option to remove local symbols from the ELF executable image.
			Note: This option does not remove any local symbols from the DWARF information in the executable image.
		X	Project>Options>Linker>Output

--no_range_reservations

ero size,
to place

--no_remove

Syntax	no_remove
Description	When this option is used, unused sections are not removed. In other words, each module that is included in the executable image contains all its original sections.
See also	<i>Keeping symbols and sections</i> , page 89. To set this option, use Project>Options>Linker>Extra Options .

no	_veneers		
	Syntax		no_veneers
	Description		Use this option to disable the insertion of veneers even though the executable image needs it. In this case, the linker will generate a relocation error for each reference that needs a veneer.
	See also		Veneers, page 95.
		X	To set this option, use Project>Options>Linker>Extra Options.
no	_vfe		
	Syntax		no_vfe
	Description		Use this option to disable the Virtual Function Elimination optimization. All virtual functions in all classes with at least one instance will be kept, and Runtime Type Information data will be kept for all polymorphic classes. Also, no warning message will be issued for modules that lack VFE information.
	See also		vfe, page 294 and Virtual Function Elimination, page 192.
		VP	To set related options, choose:
		93	Project>Options>Linker>Advanced>PerformC++ Virtual Function Elimination

--no_warnings

Syntax

--no_warnings

Description

By default, the linker issues warning messages. Use this option to disable all warning messages.



This option is not available in the IDE.

--no_wrap_diagnostics

Syntax	no_wrap_diagnostics
Description	By default, long lines in diagnostic messages are broken into several lines to make the message easier to read. Use this option to disable line wrapping of diagnostic messages.
X	This option is not available in the IDE.
only_stdout	
Syntax	only_stdout
Description	Use this option to make the linker use the standard output stream (stdout) also for messages that are normally directed to the error output stream (stderr).
X	This option is not available in the IDE.
output, -o	
Syntax	output {filename directory} -o {filename directory}
Parameters	For information about specifying a filename or a directory, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.
Description	By default, the object executable image produced by the linker is located in a file with the name a.out. Use this option to explicitly specify a different output filename, which by default will have the filename extension out.
X	Project>Options>Linker>Output>Output file

pi_	veneers
-	

Syntax

--pi_veneers

Description

Use this option to make the linker generate position-independent veneers. Note that this type of veneer is larger and slower than normal veneers.

See also

Veneers, page 95.

To set this option, use **Project>Options>Linker>Extra Options**.

--place_holder

Syntax	place_holder sy	place_holder symbol[,size[,section[,alignment]]]		
Parameters	symbol size section alignment	The name of the symbol to create Size in ROM; by default 4 bytes Section name to use; by default .text Alignment of section; by default 1		
Description	checksum calculated b with the specified nam	Use this option to reserve a place in ROM to be filled by some other tool, for example a checksum calculated by ielftool. Each use of this linker option results in a section with the specified name, size, and alignment. The symbol can be used by your application to refer to the section.		
	be included in your ap	Note: Like any other section, sections created by theplace_holder option will only be included in your application if the section appears to be needed. Thekeep linker option, or the keep linker directive can be used for forcing such section to be included.		
See also <i>IAR utilities</i> , page 437.				
	To set this option, use	Project>Options>Linker>Extra Options		

--redirect

Syntax	redirect from_:	redirect from_symbol=to_symbol		
Parameters	from_symbol to_symbol	The name of the source symbol The name of the destination symbol		
Description	Ĩ	Use this option to change a reference from one symbol to another symbol. To set this option, use Project>Options>Linker>Extra Options		

--remarks

Syntax		remarks
Description		The least severe diagnostic messages are called remarks. A remark indicates a source code construct that may cause strange behavior in the generated code. By default, the linker does not generate remarks. Use this option to make the linker generate remarks.
See also		Severity levels, page 225.
	X	Project>Options>Linker>Diagnostics>Enable remarks
search		
Syntax		search path
Parameters		pathA path to a directory where the linker should search for object and library files.
Description		Use this option to specify more directories for the linker to search for object and library files in.
		By default, the linker searches for object and library files only in the working directory. Each use of this option on the command line adds another search directory.
See also		The linking process, page 71.
	X	This option is not available in the IDE.
semihosting		

Syntax	semihosting[=iar_breakpoint]		
Parameters	iar_breakpoint	The IAR-specific mechanism can be used when debugging applications that use SWI/SVC extensively.	

Description		Use this option to include the debug interface—breakpoint mechanism—in the output image. If no parameter is specified, the SWI/SVC mechanism is included for ARM7/9/11, and the BKPT mechanism is included for Cortex-M.
See also		Application debug support, page 107.
	X	Project>Options>General Options>Library Configuration>Semihosted

--silent

Syntax		silent
Description		By default, the linker issues introductory messages and a final statistics report. Use this option to make the linker operate without sending these messages to the standard output stream (normally the screen).
		This option does not affect the display of error and warning messages.
	X	This option is not available in the IDE.

--skip_dynamic_initialization

Syntax		skip_dynamic_initialization
Description		Use this option to suppress dynamic initialization to be performed during system startup. Typically, this can be useful if you need to set up, for example, heap management for an RTOS before the initialization takes place.
		In this case you must add a call to the library function iar_dynamic_initialization in your application source code. Initialization will then take place at the time of the call to this function.
	X	To set this option, use Project>Options>Linker>Extra Options.

--stack_usage_control

Syntax	stack_usage_control filename
Parameters	For information about specifying a filename, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.

	Description		analysis, or provides mo use this option multiple t	y a stack usage control file. This file controls stack usage re stack usage information for modules or functions. You can imes to specify multiple stack usage control files. If no filename he extension suc is used.
			Using this option enable	s stack usage analysis in the linker.
	See also		Stack usage analysis, pa	ge 78
		X	Project>Options>Link	er>Advanced>Control file
str	·ip			
	Syntax		strip	
				ains the debug information from the input object files in the . Use this option to remove that information.
		¥.	To set related options, cl	noose:
		I S	Project>Options>Linker>Output>Include debug information in output	
vf€	2			
	Syntax		vfe[=forced]	
	Parameters		forced	Performs Virtual Function Elimination even if one or more modules lack the needed virtual function elimination information.
	Description		Use this option to enable the Virtual Function Elimination optimization. By default, this optimization is enabled.	
			Forcing the use of Virtual Function Elimination can be unsafe if some of the modules that lack the needed information perform virtual function calls or use dynamic Runtime Type Information.	
	See also		no_vfe, page 289 and Virtual Function Elimination, page 192.	
To set related options, choose:		noose:		
		I A	Project>Ontions>Link	er>Advanced>Perform C++ Virtual Function Flimination

Project>Options>Linker>Advanced>Perform C++ Virtual Function Elimination

--warnings_affect_exit_code

Syntax		warnings_affect_exit_code
Description		By default, the exit code is not affected by warnings, because only errors produce a non-zero exit code. With this option, warnings will also generate a non-zero exit code.
	X	This option is not available in the IDE.

--warnings_are_errors

Syntax		warnings_are_errors
Description		Use this option to make the linker treat all warnings as errors. If the linker encounters an error, no executable image is generated. Warnings that have been changed into remarks are not treated as errors.
		Note: Any diagnostic messages that have been reclassified as warnings by the optiondiag_warning will also be treated as errors whenwarnings_are_errors is used.
See also		diag_warning, page 277 anddiag_warning, page 277.
	X	Project>Options>Linker>Diagnostics>Treat all warnings as errors

Descriptions of linker options

Data representation

This chapter describes the data types, pointers, and structure types supported by the compiler.

See the chapter *Efficient coding for embedded applications* for information about which data types and pointers provide the most efficient code for your application.

Alignment

Every C data object has an alignment that controls how the object can be stored in memory. Should an object have an alignment of, for example, 4, it must be stored on an address that is divisible by 4.

The reason for the concept of alignment is that some processors have hardware limitations for how the memory can be accessed.

Assume that a processor can read 4 bytes of memory using one instruction, but only when the memory read is placed on an address divisible by 4. Then, 4-byte objects, such as long integers, will have alignment 4.

Another processor might only be able to read 2 bytes at a time; in that environment, the alignment for a 4-byte long integer might be 2.

A structure type will have the same alignment as the structure member with the most strict alignment. To decrease the alignment requirements on the structure and its members, use #pragma pack or the __packed data type attribute.

All data types must have a size that is a multiple of their alignment. Otherwise, only the first element of an array would be guaranteed to be placed in accordance with the alignment requirements. This means that the compiler might add pad bytes at the end of the structure. For more information about pad bytes, see *Packed structure types*, page 307.

Note that with the #pragma data_alignment directive you can increase the alignment demands on specific variables.

ALIGNMENT ON THE ARM CORE

The alignment of a data object controls how it can be stored in memory. The reason for using alignment is that the ARM core can access 4-byte objects more efficiently when the object is stored at an address divisible by 4.

Objects with alignment 4 must be stored at an address divisible by 4, while objects with alignment 2 must be stored at addresses divisible by 2.

The compiler ensures this by assigning an alignment to every data type, ensuring that the ARM core will be able to read the data.

For related information, see *--align_sp_on_irq*, page 234 and *--no_const_align*, page 251.

Byte order

The ARM core stores data in either little-endian or big-endian byte order. To specify the byte order, use the --endian compiler option; see *--endian*, page 244.

In the little-endian byte order, which is default, the *least* significant byte is stored at the lowest address in memory. The *most* significant byte is stored at the highest address.

In the big-endian byte order, the *most* significant byte is stored at the lowest address in memory. The *least* significant byte is stored at the highest address. If you use the big-endian byte order, it might be necessary to use the

#pragma bitfields=reversed directive to be compatible with code for other compilers and I/O register definitions of some devices, see *Bitfields*, page 300.

Note: There are two variants of the big-endian mode, BE8 and BE32, which you specify at link time. In BE8 data is big-endian and code is little-endian. In BE32 both data and code are big-endian. In architectures before v6, the BE32 endian mode is used, and after v6 the BE8 mode is used. In the v6 (ARM11) architecture, both big-endian modes are supported.

Basic data types

The compiler supports both all Standard C basic data types and some additional types.

INTEGER TYPES

This table gives the size and range of each integer data type:

Data type	Size	Range	Alignment
bool	8 bits	0 to I	I
char	8 bits	0 to 255	I
signed char	8 bits	-128 to 127	I
unsigned char	8 bits	0 to 255	I

Table 32: Integer types

Data type	Size	Range	Alignment
signed short	16 bits	-32768 to 32767	2
unsigned short	16 bits	0 to 65535	2
signed int	32 bits	-2 ³¹ to 2 ³¹ -1	4
unsigned int	32 bits	-2 ³¹ to 2 ³¹ -1	4
signed long	32 bits	-2 ³¹ to 2 ³¹ -1	4
unsigned long	32 bits	0 to 2 ³² -1	4
signed long long	64 bits	-2 ⁶³ to 2 ⁶³ -1	8
unsigned long long	64 bits	0 to 2 ⁶⁴ -1	8

Table 32: Integer types (Continued)

Signed variables are represented using the two's complement form.

Bool

The bool data type is supported by default in the C++ language. If you have enabled language extensions, the bool type can also be used in C source code if you include the file stdbool.h. This will also enable the boolean values false and true.

The enum type

The compiler will use the smallest type required to hold enum constants, preferring signed rather than unsigned.

When IAR Systems language extensions are enabled, and in C++, the enum constants and types can also be of the type long, unsigned long, long long, or unsigned long long.

To make the compiler use a larger type than it would automatically use, define an enum constant with a large enough value. For example:

For related information, see --enum is int, page 245.

The char type

The char type is by default unsigned in the compiler, but the --char_is_signed compiler option allows you to make it signed. Note, however, that the library is compiled with the char type as unsigned.

The wchar_t type

The wchar_t data type is an integer type whose range of values can represent distinct codes for all members of the largest extended character set specified among the supported locals.

The wchar_t data type is supported by default in the C++ language. To use the wchar_t type also in C source code, you must include the file stddef.h from the runtime library.

Bitfields

In Standard C, int, signed int, and unsigned int can be used as the base type for integer bitfields. In standard C++, and in C when language extensions are enabled in the compiler, any integer or enumeration type can be used as the base type. It is implementation defined whether a plain integer type (char, short, int, etc) results in a signed or unsigned bitfield.

In the IAR C/C++ Compiler for ARM, plain integer types are treated as unsigned.

Bitfields in expressions are treated as int if int can represent all values of the bitfield. Otherwise, they are treated as the bitfield base type.

Each bitfield is placed in the next container of its base type that has enough available bits to accommodate the bitfield. Within each container, the bitfield is placed in the first available byte or bytes, taking the byte order into account.

In addition, the compiler supports an alternative bitfield allocation strategy (disjoint types), where bitfield containers of different types are not allowed to overlap. Using this allocation strategy, each bitfield is placed in a new container if its type is different from that of the previous bitfield, or if the bitfield does not fit in the same container as the previous bitfield. Within each container, the bitfield is placed from the least significant bit to the most significant bit (disjoint types) or from the most significant bit to the least significant bit (reverse disjoint types). This allocation strategy will never use less space than the default allocation strategy (joined types), and can use significantly more space when mixing bitfield types.

Use the #pragma bitfield directive to choose which bitfield allocation strategy to use, see *bitfields*, page 327.

Example

Assume this example:

```
struct BitfieldExample {
    uint32_t a : 12;
    uint16_t b : 3;
    uint16_t c : 7;
    uint8_t d;
};
```

The example in the joined types bitfield allocation strategy

To place the first bitfield, a, the compiler allocates a 32-bit container at offset 0 and puts a into the first and second bytes of the container.

For the second bitfield, b, a 16-bit container is needed and because there are still four bits free at offset 0, the bitfield is placed there.

For the third bitfield, c, as there is now only one bit left in the first 16-bit container, a new container is allocated at offset 2, and c is placed in the first byte of this container.

The fourth member, d, can be placed in the next available full byte, which is the byte at offset 3.

In little-endian mode, each bitfield is allocated starting from the least significant free bit of its container to ensure that it is placed into bytes from left to right.

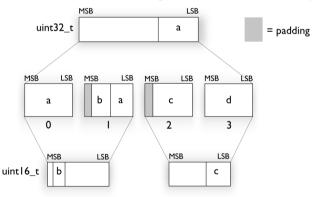


Figure 14: Layout of bitfield members for joined types in little-endian mode

In big-endian mode, each bitfield is allocated starting from the most significant free bit of its container to ensure that it is placed into bytes from left to right.

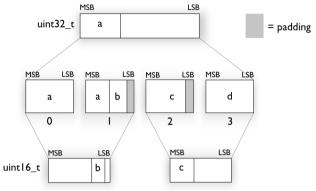


Figure 15: Layout of bitfield members for joined types in big-endian mode

The example in the disjoint types bitfield allocation strategy

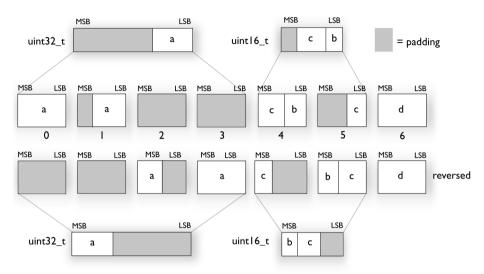
To place the first bitfield, a, the compiler allocates a 32-bit container at offset 0 and puts a into the least significant 12 bits of the container.

To place the second bitfield, b, a new container is allocated at offset 4, because the type of the bitfield is not the same as that of the previous one. b is placed into the least significant three bits of this container.

The third bitfield, c, has the same type as b and fits into the same container.

The fourth member, d, is allocated into the byte at offset 6. d cannot be placed into the same container as b and c because it is not a bitfield, it is not of the same type, and it would not fit.

When using reverse order (reverse disjoint types), each bitfield is instead placed starting from the most significant bit of its container.



This is the layout of bitfield_example in little-endian mode:

Figure 16: Layout of bitfield_example for disjoint types in little-endian mode

This is the layout of bitfield_example in big-endian mode:

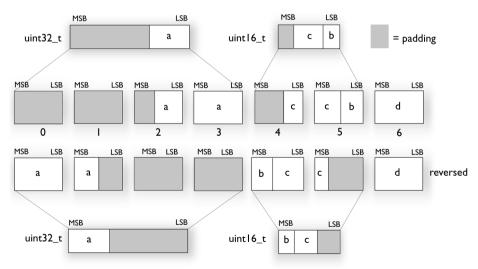


Figure 17: Layout of bitfield_example for disjoint types in big-endian mode

FLOATING-POINT TYPES

In the IAR C/C++ Compiler for ARM, floating-point values are represented in standard IEEE 754 format. The sizes for the different floating-point types are:

Туре	Size	Range (+/-)	Decimals	Exponent	Manissa	Alignment
float	32 bits	±1.18E-38 to ±3.40E+38	7	8 bits	23 bits	4
double	64 bits	±2.23E-308 to ±1.79E+308	15	II bits	52 bits	8
long double	64 bits	±2.23E-308 to ±1.79E+308	15	bits	52 bits	8

Table 33: Floating-point types

For Cortex-M0 and Cortex-M1, the compiler does not support subnormal numbers. All operations that should produce subnormal numbers will instead generate zero. For information about the representation of subnormal numbers for other cores, see *Representation of special floating-point numbers*, page 305.

Floating-point environment

Exception flags for floating-point values are supported for devices with a VFP unit, and they are defined in the fenv.h file. For devices without a VFP unit, the functions defined in the fenv.h file exists but have no functionality.

The feraiseexcept function does not raise an inexact floating-point exception when called with FE_OVERFLOW or FE_UNDERFLOW.

32-bit floating-point format

The representation of a 32-bit floating-point number as an integer is:

31	30	23 22	0
S	Exponent	Mantissa	

The exponent is 8 bits, and the mantissa is 23 bits.

The value of the number is:

(-1)^S * 2^(Exponent-127) * 1.Mantissa

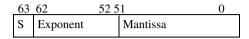
The range of the number is at least:

±1.18E-38 to ±3.39E+38

The precision of the float operators (+, -, *, and /) is approximately 7 decimal digits.

64-bit floating-point format

The representation of a 64-bit floating-point number as an integer is:



The exponent is 11 bits, and the mantissa is 52 bits.

The value of the number is:

(-1)^S * 2^(Exponent-1023) * 1.Mantissa

The range of the number is at least:

±2.23E-308 to ±1.79E+308

The precision of the float operators (+, -, *, and /) is approximately 15 decimal digits.

Representation of special floating-point numbers

This list describes the representation of special floating-point numbers:

- Zero is represented by zero mantissa and exponent. The sign bit signifies positive or negative zero.
- Infinity is represented by setting the exponent to the highest value and the mantissa to zero. The sign bit signifies positive or negative infinity.
- Not a number (NaN) is represented by setting the exponent to the highest value and the most significant bit in the mantissa to 1. The value of the sign bit is ignored.
- Subnormal numbers are used for representing values smaller than what can be represented by normal values. The drawback is that the precision will decrease with smaller values. The exponent is set to 0 to signify that the number is subnormal, even though the number is treated as if the exponent was 1. Unlike normal numbers, subnormal numbers do not have an implicit 1 as the most significant bit (the MSB) of the mantissa. The value of a subnormal number is:

(-1)^S * 2^(1-BIAS) * 0.Mantissa

where BIAS is 127 and 1023 for 32-bit and 64-bit floating-point values, respectively.

Pointer types

The compiler has two basic types of pointers: function pointers and data pointers.

FUNCTION POINTERS

The size of function pointers is always 32 bits, and the range is 0x0-0xFFFFFFFF.

When function pointer types are declared, attributes are inserted before the * sign, for example:

typedef void (__thumb __interwork * IntHandler) (void);

This can be rewritten using #pragma directives:

```
#pragma type_attribute=__thumb __interwork
typedef void IntHandler_function(void);
typedef IntHandler_function *IntHandler;
```

DATA POINTERS

There is one data pointer available. Its size is 32 bits and the range is 0x0-0xFFFFFFFF.

CASTING

Casts between pointers have these characteristics:

- Casting a *value* of an integer type to a pointer of a smaller type is performed by truncation
- Casting a value of an unsigned integer type to a pointer of a larger type is performed by zero extension
- Casting a value of a signed integer type to a pointer of a larger type is performed by sign extension
- Casting a pointer type to a smaller integer type is performed by truncation
- Casting a *pointer type* to a larger integer type is performed by zero extension
- Casting a *data pointer* to a function pointer and vice versa is illegal
- Casting a *function pointer* to an integer type gives an undefined result

size_t

size_t is the unsigned integer type required to hold the maximum size of an object. In the IAR C/C++ Compiler for ARM, the size of size_t is 32 bits.

ptrdiff_t

ptrdiff_t is the type of the signed integer required to hold the difference between two pointers to elements of the same array. In the IAR C/C++ Compiler for ARM, the size of ptrdiff_t is 32 bits.

intptr_t

intptr_t is a signed integer type large enough to contain a void *. In the IAR C/C++ Compiler for ARM, the size of intptr_t is 32 bits.

uintptr_t

uintptr_t is equivalent to intptr_t, with the exception that it is unsigned.

Structure types

The members of a struct are stored sequentially in the order in which they are declared: the first member has the lowest memory address.

ALIGNMENT

The struct and union types have the same alignment as the member with the highest alignment requirement. The size of a struct is also adjusted to allow arrays of aligned structure objects.

GENERAL LAYOUT

Members of a struct are always allocated in the order specified in the declaration. Each member is placed in the struct according to the specified alignment (offsets).

Example

```
struct First
{
    char c;
    short s;
} s;
```

This diagram shows the layout in memory:

с	pad	s		
0	1	2	3	

Figure 18: Structure layout

The alignment of the structure is 2 bytes, and a pad byte must be inserted to give short s the correct alignment.

PACKED STRUCTURE TYPES

The __packed data type attribute or the #pragma pack directive is used for relaxing the alignment requirements of the members of a structure. This changes the layout of the structure. The members are placed in the same order as when declared, but there might be less pad space between members.

Note that accessing an object that is not correctly aligned requires code that is both larger and slower. If such structure members are accessed many times, it is usually better to construct the correct values in a struct that is not packed, and access this struct instead.

Special care is also needed when creating and using pointers to misaligned members. For direct access to misaligned members in a packed struct, the compiler will emit the correct (but slower and larger) code when needed. However, when a misaligned member is accessed through a pointer to the member, the normal (smaller and faster) code is used. In the general case, this will not work, because the normal code might depend on the alignment being correct.

Example

This example declares a packed structure:

```
#pragma pack(1)
struct S
{
    char c;
    short s;
};
```

#pragma pack()

In this example, the structure S has this memory layout:

l	c	s	
	0		2

Figure 19: Packed structure layout

This example declares a new non-packed structure, S2, that contains the structure s declared in the previous example:

```
struct S2
{
   struct S s;
   long l;
};
```

S2 has this memory layout

с	s		pad		1		
0		2	3	4	5	6	7

Figure 20: Packed structure layout

The structure s will use the memory layout, size, and alignment described in the previous example. The alignment of the member 1 is 4, which means that alignment of the structure s_2 will become 4.

For more information, see Alignment of elements in a structure, page 199.

Type qualifiers

According to the C standard, volatile and const are type qualifiers.

DECLARING OBJECTS VOLATILE

By declaring an object volatile, the compiler is informed that the value of the object can change beyond the compiler's control. The compiler must also assume that any accesses can have side effects—thus all accesses to the volatile object must be preserved.

There are three main reasons for declaring an object volatile:

- Shared access; the object is shared between several tasks in a multitasking environment
- Trigger access; as for a memory-mapped SFR where the fact that an access occurs has an effect
- Modified access; where the contents of the object can change in ways not known to the compiler.

Definition of access to volatile objects

The C standard defines an abstract machine, which governs the behavior of accesses to volatile declared objects. In general and in accordance to the abstract machine:

- The compiler considers each read and write access to an object declared volatile as an access
- The unit for the access is either the entire object or, for accesses to an element in a composite object—such as an array, struct, class, or union—the element. For example:

```
char volatile a;
a = 5;  /* A write access */
a += 6;  /* First a read then a write access */
```

- An access to a bitfield is treated as an access to the underlying type
- Adding a const qualifier to a volatile object will make write accesses to the object impossible. However, the object will be placed in RAM as specified by the C standard.

However, these rules are not detailed enough to handle the hardware-related requirements. The rules specific to the IAR C/C++ Compiler for ARM are described below.

Rules for accesses

In the IAR C/C++ Compiler for ARM, accesses to volatile declared objects are subject to these rules:

- All accesses are preserved
- · All accesses are complete, that is, the whole object is accessed
- All accesses are performed in the same order as given in the abstract machine
- All accesses are atomic, that is, they cannot be interrupted.

The compiler adheres to these rules for accesses to all 8-, 16-, and 32-bit scalar types, except for accesses to unaligned 16- and 32-bit fields in packed structures.

For all combinations of object types not listed, only the rule that states that all accesses are preserved applies.

DECLARING OBJECTS VOLATILE AND CONST

If you declare a volatile object const, it will be write-protected but it will still be stored in RAM memory as the C standard specifies.

To store the object in read-only memory instead, but still make it possible to access it as a const volatile object, define the variable like this:

const volatile int x @ "FLASH";

The compiler will generate the read/write section FLASH. That section should be placed in ROM and is used for manually initializing the variables when the application starts up.

Thereafter, the initializers can be reflashed with other values at any time.

DECLARING OBJECTS CONST

The const type qualifier is used for indicating that a data object, accessed directly or via a pointer, is non-writable. A pointer to const declared data can point to both constant and non-constant objects. It is good programming practice to use const declared pointers whenever possible because this improves the compiler's possibilities to optimize the generated code and reduces the risk of application failure due to erroneously modified data.

Static and global objects declared const are allocated in ROM.

In C++, objects that require runtime initialization cannot be placed in ROM.

Data types in C++

In C++, all plain C data types are represented in the same way as described earlier in this chapter. However, if any Embedded C++ features are used for a type, no assumptions can be made concerning the data representation. This means, for example, that it is not supported to write assembler code that accesses class members.

Data types in C++

Extended keywords

This chapter describes the extended keywords that support specific features of the ARM core and the general syntax rules for the keywords. Finally the chapter gives a detailed description of each keyword.

General syntax rules for extended keywords

To understand the syntax rules for the extended keywords, it is important to be familiar with some related concepts.

The compiler provides a set of attributes that can be used on functions or data objects to support specific features of the ARM core. There are two types of attributes—*type attributes* and *object attributes*:

- Type attributes affect the *external functionality* of the data object or function
- Object attributes affect the *internal functionality* of the data object or function.

The syntax for the keywords differs slightly depending on whether it is a type attribute or an object attribute, and whether it is applied to a data object or a function.

Note: The extended keywords are only available when language extensions are enabled in the compiler.



In the IDE, language extensions are enabled by default.

>__

Use the -e compiler option to enable language extensions. See *-e*, page 242 for more information.

TYPE ATTRIBUTES

Type attributes define how a function is called, or how a data object is accessed. This means that if you use a type attribute, it must be specified both when a function or data object is defined and when it is declared.

You can either place the type attributes explicitly in your declarations, or use the pragma directive #pragma type_attribute.

General type attributes

These general type attributes are available:

• *Function type attributes* affect how the function should be called: __arm, __fiq, __interwork, __irq, __swi, __task, and __thumb

• Data type attributes: __big_endian, __little_endian, and __packed.

You can specify as many type attributes as required for each level of pointer indirection.

Syntax for type attributes used on data objects

In general, type attributes for data objects follow the same syntax as the type qualifiers const and volatile.

The following declaration assigns the__little_endian type attribute to the variables i and j; in other words, the variable i and j will be accessed with little endian byte order. However, note that an individual member of a struct or union cannot have a type attribute. The variables k and l behave in the same way:

```
__little_endian int i, j;
int __little_endian k, l;
```

Note that the attribute affects both identifiers.

This declaration of i and j is equivalent with the previous one:

```
#pragma type_attribute=__little_endian
int i, j;
```

The advantage of using pragma directives for specifying keywords is that it offers you a method to make sure that the source code is portable.

Syntax for type attributes on data pointers

The syntax for declaring pointers using type attributes follows the same syntax as the type qualifiers const and volatile:

intpacked * p;	A pointer to a packed integer.
int *packed p;	A packed pointer to an integer.
<pre>packed int * p;</pre>	A packed pointer to an integer.

Syntax for type attributes on functions

The syntax for using type attributes on functions differs slightly from the syntax of type attributes on data objects. For functions, the attribute must be placed either in front of the return type, or in parentheses, for example:

```
__irq __arm void my_handler(void);
or
void (__irq __arm my_handler)(void);
```

This declaration of my_handler is equivalent with the previous one:

```
#pragma type_attribute=__irq __arm
void my_handler(void);
```

OBJECT ATTRIBUTES

Normally, object attributes affect the internal functionality of functions and data objects, but not directly how the function is called or how the data is accessed. This means that an object attribute does not normally need to be present in the declaration of an object.

These object attributes are available:

- Object attributes that can be used for variables: __absolute, __no_init
- Object attributes that can be used for functions and variables: location, @, __root, and __weak
- Object attributes that can be used for functions: __intrinsic, __nested, __noreturn, __ramfunc, and __stackless.

You can specify as many object attributes as required for a specific function or data object.

For more information about location and @, see *Controlling data and function* placement in memory, page 201.

Syntax for object attributes

The object attribute must be placed in front of the type. For example, to place myarray in memory that is not initialized at startup:

```
__no_init int myarray[10];
```

The #pragma object_attribute directive can also be used. This declaration is equivalent to the previous one:

```
#pragma object_attribute=__no_init
int myarray[10];
```

Note: Object attributes cannot be used in combination with the typedef keyword.

Summary of extended keywords

This table summarizes the extended keywords:

Extended keyword	Description
absolute	Makes references to the object use absolute addressing

Table 34: Extended keywords summary

Extended keyword	Description
arm	Makes a function execute in ARM mode
big_endian	Declares a variable to use the big-endian byte order
fiq	Declares a fast interrupt function
interwork	Declares a function to be callable from both ARM and Thumb mode
intrinsic	Reserved for compiler internal use only
irq	Declares an interrupt function
little_endian	Declares a variable to use the little-endian byte order
nested	Allows an $__irq$ declared interrupt function to be nested, that is, interruptible by the same type of interrupt
no_init	Places a data object in non-volatile memory
noreturn	Informs the compiler that the function will not return
packed	Decreases data type alignment to I
pcrel	Used internally by the compiler for constant data when theropi compiler option is used
ramfunc	Makes a function execute in RAM
root	Ensures that a function or variable is included in the object code even if unused
sbrel	Used internally by the compiler for constant data when therwpi compiler option is used
stackless	Makes a function callable without a working stack
swi	Declares a software interrupt function
task	Relaxes the rules for preserving registers
thumb	Makes a function execute in Thumb mode
weak	Declares a symbol to be externally weakly linked

Table 34: Extended keywords summary (Continued)

Descriptions of extended keywords

These sections give detailed information about each extended keyword.

__absolute

Syntax

Follows the generic syntax rules for object attributes, see Object attributes, page 315.

Description	Theabsolute keyword makes references to the object use absolute addressing.
	The following limitations apply:
	• Only available when theropi orrwpi compiler option is used
	• Can only be used on external declarations.
Example	<pre>externabsolute char otherBuffer[100];</pre>
arm	
Syntax	Follows the generic syntax rules for type attributes that can be used on functions, see <i>Type attributes</i> , page 313.
Description	Thearm keyword makes a function execute in ARM mode. Anarm declared function can, unless it is also declaredinterwork, only be called from functions that also execute in ARM mode.
	A function declaredarm cannot be declaredthumb.
	Note: Non-interwork ARM functions cannot be called from Thumb mode.
Example	<pre>arm int func1(void);</pre>
See also	interwork, page 318.

__big_endian

Syntax	Follows the generic syntax rules for type attributes that can be used on data objects, see <i>Type attributes</i> , page 313.
Description	Thebig_endian keyword is used for accessing a variable that is stored in the big-endian byte order regardless of what byte order the rest of the application uses. Thebig_endian keyword is available when you compile for ARMv6 or higher.
	Note that this keyword cannot be used on pointers. Also, this attribute cannot be used on arrays.
Example	<pre>big_endian long my_variable;</pre>
See also	little_endian, page 319.

fiq	
Syntax	Follows the generic syntax rules for type attributes that can be used on functions, see <i>Type attributes</i> , page 313.
Description	Thefiq keyword declares a fast interrupt function. All interrupt functions must be compiled in ARM mode. A function declaredfiq does not accept parameters and does not have a return value. This keyword is not available when you compile for Cortex-M devices.
Example	<pre>fiqarm void interrupt_function(void);</pre>
interwork	
Syntax	Follows the generic syntax rules for type attributes that can be used on functions, see <i>Type attributes</i> , page 313.
Description	A function declaredinterwork can be called from functions executing in either ARM or Thumb mode.
	Note: By default, functions are interwork when theinterwork compiler option is used, and when thecpu option is used and it specifies a core where interwork is default.
Example	<pre>typedef void (thumbinterwork *IntHandler)(void);</pre>
intrinsic	
Description	Theintrinsic keyword is reserved for compiler internal use only.
irq	
Syntax	Follows the generic syntax rules for type attributes that can be used on functions, see <i>Type attributes</i> , page 313.
Description	Theirq keyword declares an interrupt function. All interrupt functions must be compiled in ARM mode. A function declaredirq does not accept parameters and does not have a return value. This keyword is not available when you compile for Cortex-M devices.
Example	<pre>irqarm void interrupt_function(void);</pre>

Syntax	Follows the generic syntax rules for type attributes that can be used on data objects <i>Type attributes</i> , page 313.
Description	Thelittle_endian keyword is used for accessing a variable that is stored in little-endian byte order regardless of what byte order the rest of the application useslittle_endian keyword is available when you compile for ARMv6 or higher
	Note that this keyword cannot be used on pointers. Also, this attribute cannot be use arrays.
Example	<pre>little_endian long my_variable;</pre>
See also	<i>big_endian</i> , page 317.
sted	
Syntax	Follows the generic syntax rules for object attributes, see Object attributes, page 3
Description	Thenested keyword modifies the enter and exit code of an interrupt function allow for nested interrupts. This allows interrupts to be enabled, which means new interrupts can be served inside an interrupt function, without overwriting the SPSR return address in R14. Nested interrupts are only supported forirq declared functions.
	Note: The <u>nested</u> keyword requires the processor mode to be in either User of System mode.
Example	<pre>irqnestedarm void interrupt_handler(void);</pre>
	Nested interrupts, page 61 andalign sp on irq, page 234.

--align_sp_on_irq, page 234

See also

Syntax	Follows the generic syntax rules for object attributes, see Object attributes, page 315.
Description	Use the <u>no_init</u> keyword to place a data object in non-volatile memory. This means that the initialization of the variable, for example at system startup, is suppressed.

Example	<pre>no_init int myarray[10];</pre>
See also	Do not initialize directive, page 410.
noreturn	
Syntax	Follows the generic syntax rules for object attributes, see Object attributes, page 315.
Description	Thenoreturn keyword can be used on a function to inform the compiler that the function will not return. If you use this keyword on such functions, the compiler can optimize more efficiently. Examples of functions that do not return are abort and exit.
Example	<pre>noreturn void terminate(void);</pre>
packed	
Syntax	Follows the generic syntax rules for type attributes that can be used on data, see <i>Type attributes</i> , page 313.
Description	Use thepacked keyword to decrease the data type alignment to 1packed can be used for two purposes:
	• When used with a struct or union type definition, the maximum alignment of members of that struct or union is set to 1, to eliminate any gaps between the members. The type of each members also receives thepacked type attribute.
	• When used with any other type, the resulting type is the same as the type without thepacked type attribute, but with an alignment of 1. Types that already have an alignment of 1 are not affected by thepacked type attribute.
	A normal pointer can be implicitly converted to a pointer topacked, but the reverse conversion requires a cast.
	Note: Accessing data types at other alignments than their natural alignment can result in code that is significantly larger and slower.
	Use eitherpacked or #pragma pack to relax the alignment restrictions for a type and the objects defined using that type. Mixingpacked and #pragma pack might lead to unexpected behavior.

Example	packed struct X {char ch; int i;}; /* No pad bytes */ void Foo(struct X * xp) /* No need forpacked here */ {
	<pre>int * p1 = &xp->i;/* Error:"int *">"intpacked *" */ intpacked * p2 = &xp->i;</pre>
See also	<i>pack</i> , page 336.

___ramfunc

Syntax	Follows the generic syntax rules for object attributes, see Object attributes, page 315.
Description	Theramfunc keyword makes a function execute in RAM. Two code sections will be created: one for the RAM execution (.textrw), and one for the ROM initialization (.textrw_init).
	If a function declaredramfunc tries to access ROM, the compiler will issue a warning. This behavior is intended to simplify the creation of <i>upgrade</i> routines, for instance, rewriting parts of flash memory. If this is not why you have declared the functionramfunc, you can safely ignore or disable these warnings.
	Functions declaredramfunc are by default stored in the section named .textrw.
Example	<pre>ramfunc int FlashPage(char * data, char * page);</pre>
See also	To read more aboutramfunc declared functions in relation to breakpoints, see the <i>C-SPY</i> ® <i>Debugging Guide for ARM</i> .

__root

Syntax	Follows the generic syntax rules for object attributes, see Object attributes, page 315.
Description	A function or variable with theroot attribute is kept whether or not it is referenced from the rest of the application, provided its module is included. Program modules are always included and library modules are only included if needed.
Example	<pre>root int myarray[10];</pre>
See also	For more information about root symbols and how they are kept, see <i>Keeping symbols and sections</i> , page 89.

stackless	
Syntax	Follows the generic syntax rules for object attributes, see Object attributes, page 315.
Description	Thestackless keyword declares a function that can be called without a working stack.
ſ	A function declaredstackless violates the calling convention in such a way that it is not possible to return from it. However, the compiler cannot reliably detect if the function returns and will not issue an error if it does.
Example	<pre>stackless void start_application(void);</pre>
swi	
Syntax	Follows the generic syntax rules for type attributes that can be used on functions, see <i>Type attributes</i> , page 313.
Description	Theswi declares a software interrupt function. It inserts an SVC (formerly SWI) instruction and the specified software interrupt number to make a proper function call. A function declaredswi accepts arguments and returns values. Theswi keyword makes the compiler generate the correct return sequence for a specific software interrupt function. Software interrupt functions follow the same calling convention regarding parameters and return values as an ordinary function, except for the stack usage.
	Theswi keyword also expects a software interrupt number which is specified with the #pragma swi_number=number directive. The swi_number is used as an argument to the generated assembler SVC instruction, and can be used by the SVC interrupt handler, for example SWI_Handler, to select one software interrupt function in a system containing several such functions. Note that the software interrupt number should only be specified in the function declaration—typically, in a header file that you include in the source code file that calls the interrupt function—not in the function definition.
	Note: All interrupt functions must be compiled in ARM mode, except for Cortex-M. Use either thearm keyword or the #pragma type_attribute=arm directive to alter the default behavior if needed.
Example	To declare your software interrupt function, typically in a header file, write for example like this:
	<pre>#pragma swi_number=0x23swi int swi0x23_function(int a, int b);</pre>

To call the function:

Somewhere in your application source code, you define your software interrupt function:

```
...
__swi __arm int the_actual_swi0x23_function(int a, int b)
{
    ...
    return 42;
}
```

Software interrupts, page 62 and Calling convention, page 150.

__task

See also

Syntax	Follows the generic syntax rules for type attributes that can be used on functions, see <i>Type attributes</i> , page 313.
Description	This keyword allows functions to relax the rules for preserving registers. Typically, the keyword is used on the start function for a task in an RTOS.
	By default, functions save the contents of used preserved registers on the stack upon entry, and restore them at exit. Functions that are declaredtask do not save all registers, and therefore require less stack space.
	Because a function declaredtask can corrupt registers that are needed by the calling function, you should only usetask on functions that do not return or call such a function from assembler code.
	The function main can be declaredtask, unless it is explicitly called from the application. In real-time applications with more than one task, the root function of each task can be declaredtask.
Example	<pre>task void my_handler(void);</pre>

thumb	
Syntax	Follows the generic syntax rules for type attributes that can be used on functions, see <i>Type attributes</i> , page 313.
Description	Thethumb keyword makes a function execute in Thumb mode. Unless the function

Description	Thethumb keyword makes a function execute in Thumb mode. Unless the function is also declaredinterwork, the function declaredthumb can only be called from functions that also execute in Thumb mode.
	A function declaredthumb cannot be declaredarm.
	Note: Non-interwork Thumb functions cannot be called from ARM mode.
Example	<pre>thumb int func2(void);</pre>
See also	<i>interwork</i> , page 318.

__weak

Syntax	Follows the generic syntax rules for object attributes, see <i>Object attributes</i> , page 315.
Description	Using theweak object attribute on an external declaration of a symbol makes all references to that symbol in the module weak.
	Using theweak object attribute on a public definition of a symbol makes that definition a weak definition.
	The linker will not include a module from a library solely to satisfy weak references to a symbol, nor will the lack of a definition for a weak reference result in an error. If no definition is included, the address of the object will be zero.
	When linking, a symbol can have any number of weak definitions, and at most one non-weak definition. If the symbol is needed, and there is a non-weak definition, this definition will be used. If there is no non-weak definition, one of the weak definitions will be used.
Example	externweak int foo; /* A weak reference. */
	<pre>weak void bar(void) /* A weak definition. */ { /* Increment foo if it was included. */ if (&foo != 0) ++foo; }</pre>

Pragma directives

This chapter describes the pragma directives of the compiler.

The #pragma directive is defined by Standard C and is a mechanism for using vendor-specific extensions in a controlled way to make sure that the source code is still portable.

The pragma directives control the behavior of the compiler, for example how it allocates memory for variables and functions, whether it allows extended keywords, and whether it outputs warning messages.

The pragma directives are always enabled in the compiler.

Summary of pragma directives

This table lists the pragma directives of the compiler that can be used either with the #pragma preprocessor directive or the _Pragma() preprocessor operator:

Pragma directive	Description
bitfields	Controls the order of bitfield members
calls	Lists possible called functions for indirect calls
call_graph_root	Specifies that the function is a call graph root
data_alignment	Gives a variable a higher (more strict) alignment
diag_default	Changes the severity level of diagnostic messages
diag_error	Changes the severity level of diagnostic messages
diag_remark	Changes the severity level of diagnostic messages
diag_suppress	Suppresses diagnostic messages
diag_warning	Changes the severity level of diagnostic messages
error	Signals an error while parsing
include_alias	Specifies an alias for an include file
inline	Controls inlining of a function
language	Controls the IAR Systems language extensions

Table 35: Pragma directives summary

Pragma directive	Description
location	Specifies the absolute address of a variable, places a variable in a register, or places groups of functions or variables in named sections
message	Prints a message
object_attribute	Changes the definition of a variable or a function
optimize	Specifies the type and level of an optimization
pack	Specifies the alignment of structures and union members
printf_args	Verifies that a function with a printf-style format string is called with the correct arguments
required	Ensures that a symbol that is needed by another symbol is included in the linked output
rtmodel	Adds a runtime model attribute to the module
scanf_args	Verifies that a function with a scanf-style format string is called with the correct arguments
section	Declares a section name to be used by intrinsic functions
segment	This directive is an alias for #pragma section
STDC CX_LIMITED_RANGE	Specifies whether the compiler can use normal complex mathematical formulas or not
STDC FENV_ACCESS	Specifies whether your source code accesses the floating-point environment or not.
STDC FP_CONTRACT	Specifies whether the compiler is allowed to contract floating-point expressions or not.
swi_number	Sets the interrupt number of a software interrupt function
type_attribute	Changes the declaration and definitions of a variable or function
weak	Makes a definition a weak definition, or creates a weak alias for a function or a variable

Table 35: Pragma directives summary (Continued)

Note: For portability reasons, the pragma directives alignment, baseaddr, codeseg, constseg, dataseg, function, memory, and warnings are recognized but will give a diagnostic message. It is important to be aware of this if you need to port existing code that contains any of those pragma directives. See also *Recognized pragma directives* (6.10.6), page 474.

Descriptions of pragma directives

This section gives detailed information about each pragma directive.

bitfields

Syntax	<pre>#pragma bitfields=disjo</pre>	int_types joined_types sed_disjoint_types reversed default}
Parameters	disjoint_types	Bitfield members are placed from the least significant bit to the most significant bit in the container type. Storage containers of bitfields with different base types will not overlap.
	joined_types	Bitfield members are placed depending on the byte order. Storage containers of bitfields will overlap other structure members. For more information, see <i>Bitfields</i> , page 300.
	reversed_disjoint_types	Bitfield members are placed from the most significant bit to the least significant bit in the container type. Storage containers of bitfields with different base types will not overlap.
	reversed	This is an alias for reversed_disjoint_types.
	default	Restores to default layout of bitfield members. The default behavior for the compiler is joined_types.
Description	Use this pragma directive to co	ontrol the layout of bitfield members.
Example	<pre>struct S { unsigned char error unsigned char size unsigned short code };</pre>	disjoint bitfield types. */ : 1; : 4;
See also	Bitfields, page 300.	

calls

Syntax	<pre>#pragma calls=function[, function]</pre>	
Parameters	function Any declared function	
Description	Use this pragma directive to list the functions that can be indirectly called in the following statement. This information can be used for stack usage analysis in the linker.	
	Note: For an accurate result, you must list all possible called functions.	
Example	<pre>void Fun1(), Fun2();</pre>	
	<pre>void Caller(void (*fp)(void)) { #pragma calls = Fun1, Fun2 (*fp)(); }</pre>	
See also	Stack usage analysis, page 78	

call_graph_root

Syntax	<pre>#pragma call_graph_root[=category]</pre>	
Parameters	category	A string that identifies an optional call graph root category
Description	Use this pragma directive to specify that, for stack usage analysis purposes, the immediately following function is a call graph root. You can also specify an optional category. The compiler will usually automatically assign a call graph root category to interrupt and task functions. If you use the #pragma call_graph_root directive on such a function you will override the default category. You can specify any string as a category.	
Example	#pragma call_graph	_root="interrupt"
See also	Stack usage analysis, pa	ge 78

data_alignment

	Syntax	<pre>#pragma data_alignment=expression</pre>	
	Parameters	expression	A constant which must be a power of two (1, 2, 4, etc.).
	Description	Use this pragma directive to give a variable a higher (more strict) alignment of the start address than it would otherwise have. This directive can be used on variables with static and automatic storage duration.	
			ive on variables with automatic storage duration, there is an d alignment for each function, determined by the calling
		data_alignment direct	of a variable is a multiple of its alignment. The ive only affects the alignment of the variable's start address, thus be used for creating situations where the size is not a t.
diag	_default		
	Syntax	<pre>#pragma diag_defaul</pre>	t=tag[,tag,]
	Parameters	tag	The number of a diagnostic message, for example the message number Pe177.
	Description	severity level defined on	e to change the severity level back to the default, or to the the command line by any of the optionsdiag_error, ug_suppress, ordiag_warnings, for the diagnostic the tags.
	See also	Diagnostics, page 224.	
diag	_error		
	Syntax	<pre>#pragma diag_error=</pre>	tag[,tag,]
	D		

tag

Parameters

The number of a diagnostic message, for example the message number Pe177.

	Description	Use this pragma directiv diagnostics.	ve to change the severity level to error for the specified
	See also	Diagnostics, page 224.	
diag	_remark		
	Syntax	#pragma diag_remar	k=tag[,tag,]
	Parameters	tag	The number of a diagnostic message, for example the message number Pe177.
	Description	Use this pragma directiv diagnostic messages.	te to change the severity level to remark for the specified
	See also	Diagnostics, page 224.	
diag	_suppress		
	Syntax	#pragma diag_suppr	ess=tag[,tag,]
	Parameters	tag	The number of a diagnostic message, for example the message number Pe117.
	Description	Use this pragma directiv	e to suppress the specified diagnostic messages.
	See also	Diagnostics, page 224.	
diag	_warning		
	Syntax	#pragma diag_warni	ng=tag[,tag,]
	Parameters	tag	The number of a diagnostic message, for example the message number Pe826.
	Description	Use this pragma directiv diagnostic messages.	ve to change the severity level to warning for the specified

See also

Diagnostics, page 224.

error

	Syntax	#pragma error <i>message</i>	
	Parameters	message	A string that represents the error message.
	Description	is different from the pre- directive can be included	e to cause an error message when it is parsed. This mechanism processor directive #error, because the #pragma error d in a preprocessor macro using the _Pragma form of the s an error if the macro is used.
	Example	<pre>#if FO0_AVAILABLE #define FO0 #else #define FO0 _Pragma("error\"Foo is not available\"") #endif</pre>	
		If FOO_AVAILABLE is zero, an error will be signaled if the FOO macro is used in actual source code.	
inclu	ıde_alias		
	Syntax	<pre>#pragma include_alias ("orig_header" , "subst_header") #pragma include_alias (<orig_header> , <subst_header>)</subst_header></orig_header></pre>	
	Parameters	orig_header	The name of a header file for which you want to create an alias.
		subst_header	The alias for the original header file.

Description Use this pragma directive to provide an alias for a header file. This is useful for substituting one header file with another, and for specifying an absolute path to a relative file.

This pragma directive must appear before the corresponding <code>#include</code> directives and <code>subst_header</code> must match its corresponding <code>#include</code> directive exactly.

		This example will substitute the relative file stdio.h with a counterpart located according to the specified path.	
	See also	Include file search proc	edure, page 221.
inlin	e		
	Syntax	<pre>#pragma inline[=fc</pre>	prced =never]
	Parameters	No parameter	Has the same effect as the inline keyword.
		forced	Disables the compiler's heuristics and forces inlining.
		never	Disables the compiler's heuristics and makes sure that the function will not be inlined.
	Description	Use #pragma inline to advise the compiler that the function defined immediately after the directive should be inlined according to C++ inline semantics.	
		Specifying #pragma inline=forced will always inline the defined function. If the compiler fails to inline the function for some reason, for example due to recursion, a warning message is emitted.	
			formed only on the High optimization level. Specifying acced will enable inlining of the function also on the Medium
	See also	Inlining functions, page	65
lang	uage		
	Syntax	<pre>#pragma language={extended default save restore}</pre>	
	Parameters		
		extended	Enables the IAR Systems language extensions from the first use of the pragma directive and onward.
		default	From the first use of the pragma directive and onward, restores the settings for the IAR Systems language extensions to whatever that was specified by compiler options.

	save restore	Saves and restores, respectively, the IAR Systems language extensions setting around a piece of source code. Each use of save must be followed by a matching restore
		in the same file without any intervening #include directive.
Description	Use this pragma direct	ive to control the use of language extensions.
Example	At the top of a file that	needs to be compiled with IAR Systems extensions enabled:
	#pragma language= /* The rest of th	
	extensions enabled, bu	rt of the source code that needs to be compiled with IAR Systems t where the state before the sequence cannot be assumed to be the by the compiler options in use:
	#pragma language= #pragma language= /* Part of source #pragma language=	extended code. */
See also	-e, page 242 andstrie	<i>ct</i> , page 265.

location

Syntax	<pre>#pragma location={address register NAME}</pre>	
Parameters	address	The absolute address of the global or static variable or function for which you want an absolute location.
	register	An identifier that corresponds to one of the ARM core registers R4-R11.
	NAME	A user-defined section name; cannot be a section name predefined for use by the compiler and linker.
Description	Use this pragma directi	ve to specify:
		bsolute address—of the global or static variable whose the pragma directive. The variables must be declared

	• An identifier specifying a register. The variable defined after the pragma directive is placed in the register. The variable must be declared asno_init and have file scope.
	A string specifying a section for placing either a variable or function whose declaration follows the pragma directive. Do not place variables that would normally be in different sections (for example, variables declared asno_init and variables declared as const) in the same named section.
Example	<pre>#pragma location=0xFFFF0400no_init volatile char PORT1; /* PORT1 is located at address</pre>
	<pre>#pragma location=R8no_init int TASK; /* TASK is placed in R8 */</pre>
	<pre>#pragma location="FLASH" char PORT2; /* PORT2 is located in section FLASH */</pre>
	<pre>/* A better way is to use a corresponding mechanism */ #define FLASH _Pragma("location=\"FLASH\"") /* */ FLASH int i; /* i is placed in the FLASH section */</pre>
See also	Controlling data and function placement in memory, page 201.

message

Syntax	<pre>#pragma message(message)</pre>	
Parameters	message	The message that you want to direct to the standard output stream.
Description	Use this pragma directive to make the compiler print a message to the standard output stream when the file is compiled.	
Example	#ifdef TESTING #pragma message("Testing") #endif	

object_attribute

Syntax	<pre>#pragma object_attribute=object_attribute[,object_attribute,]</pre>	
Parameters	For information about object attributes that can be used with this pragma directive, see <i>Object attributes</i> , page 315.	
Description	Use this pragma directive to declare a variable or a function with an object attribute. This directive affects the definition of the identifier that follows immediately after the directive. The object is modified, not its type. Unlike the directive <code>#pragma</code> <code>type_attribute</code> that specifies the storing and accessing of a variable or function, it is not necessary to specify an object attribute in declarations.	
Example	<pre>#pragma object_attribute=no_init char bar;</pre>	
See also	General syntax rules for extended keywords, page 313.	
optimize		
Syntax	<pre>#pragma optimize=[goal][level][no_optimization]</pre>	
Parameters		

goal	Choose between:
	size, optimizes for size
	balanced, optimizes balanced between speed and size
	speed, optimizes for speed.
	no_size_constraints, optimizes for speed, but relaxes the normal restrictions for code size expansion.
level	Specifies the level of optimization; choose between none, low, medium, or high.

	no_optimization	Disables one or several optimizations; choose between:
		no_code_motion, disables code motion
		no_cse, disables common subexpression elimination
		-
		no_inline, disables function inlining
		no_tbaa, disables type-based alias analysis
		no_unroll, disables loop unrolling
		no_scheduling, disables instruction scheduling.
Description	Use this pragma directive to decrease the optimization level, or to turn off some spec optimizations. This pragma directive only affects the function that follows immediat after the directive.	
	effect on the high opti possible to optimize for	balanced, speed, and no_size_constraints only have imization level and only one of them can be used as it is not or speed and size at the same time. It is also not possible to use mbedded in this pragma directive. Any such macro will not be pocessor.
		pragma optimize directive to specify an optimization level that nization level you specify using a compiler option, the pragma
Example	<pre>#pragma optimize= int SmallAndUsed0 { /* Do something }</pre>	ften()
	<pre>#pragma optimize= int BigAndSeldomU { /* Do something }</pre>	sed()
See also	Fine-tuning enabled tr	ansformations, page 208.

pack

Syntax	#pragma pack	<pre>#pragma pack(n) #pragma pack() #pragma pack({push pop}[,name] [,n])</pre>	
Parameters	п	Sets an optional structure alignment; one of: 1, 2, 4, 8, or 16	
	Empty list	Restores the structure alignment to default	
	push	Sets a temporary structure alignment	
	pop	Restores the structure alignment from a temporarily pushed alignment	
	name	An optional pushed or popped alignment label	
Description	Use this pragma members.	Use this pragma directive to specify the maximum alignment of struct and union members.	
		The #pragma pack directive affects declarations of structures following the pragma directive to the next #pragma pack or the end of the compilation unit.	
	Note: This can r of the structure.	Note: This can result in significantly larger and slower code when accessing members of the structure.	
	and the objects d	Use eitherpacked or #pragma pack to relax the alignment restrictions for a typ and the objects defined using that type. Mixingpacked and #pragma pack mig lead to unexpected behavior.	
See also	Structure types,	Structure types, page 307 andpacked, page 320.	
printf_args			
Syntax	#pragmapr	<pre>#pragmaprintf_args</pre>	
Description	that function, the	directive on a function with a printf-style format string. For any call to e compiler verifies that the argument to each conversion specifier (for yntactically correct.	

```
Example #pragma __printf_args
int printf(char const *,...);
void PrintNumbers(unsigned short x)
{
    printf("%d", x); /* Compiler checks that x is an integer */
}
```

required

Syntax	<pre>#pragma required=symbol</pre>		
Parameters	<i>symbol</i> Any statically linked function or variable.		
Description	Use this pragma directive to ensure that a symbol which is needed by a second symbol is included in the linked output. The directive must be placed immediately before the second symbol.		
	Use the directive if the requirement for a symbol is not otherwise visible in the application, for example if a variable is only referenced indirectly through the section it resides in.		
Example	<pre>const char copyright[] = "Copyright by me"; #pragma required=copyright int main() { /* Do something here. */ } Even if the copyright string is not used by the application, it will still be included by the</pre>		
	linker and available in the output.		

rtmodel

Syntax	<pre>#pragma rtmodel="key","value"</pre>	
Parameters	"key" "value"	A text string that specifies the runtime model attribute. A text string that specifies the value of the runtime model attribute. Using the special value * is equivalent to not defining the attribute at all.
		defining the attribute at an.

•

Description	Use this pragma directive to add a runtime model attribute to a module, which can be used by the linker to check consistency between modules.
	This pragma directive is useful for enforcing consistency between modules. All modules that are linked together and define the same runtime attribute key must have the same value for the corresponding key, or the special value *. It can, however, be useful to state explicitly that the module can handle any runtime model.
	A module can have several runtime model definitions.
	Note: The predefined compiler runtime model attributes start with a double underscore. To avoid confusion, this style must not be used in the user-defined attributes.
Example	<pre>#pragma rtmode1="I2C","ENABLED"</pre>
	The linker will generate an error if a module that contains this definition is linked with a module that does not have the corresponding runtime model attributes defined.
See also	Checking module consistency, page 136.
scanf_args	

Syntax #pragma __scanf_args Description Use this pragma directive on a function with a scanf-style format string. For any call to that function, the compiler verifies that the argument to each conversion specifier (for example %d) is syntactically correct. #pragma ___scanf_args Example int scanf(char const *,...); int GetNumber() { int nr; scanf("%d", &nr); /* Compiler checks that the argument is a pointer to an integer */ return nr; }

section

Syntax	<pre>#pragma section="NAME"</pre>		
	alias		
	#pragma segment="NAME"		
Parameters	NAME The name of the section.		
Description	Use this pragma directive to define a section name that can be used by the section operatorssection_begin,section_end, andsection_size. All section declarations for a specific section must have the same alignment.		
Example	<pre>#pragma section="MYSECTION"</pre>		
See also	<i>Dedicated section operators</i> , page 165. For more information about sections and segment parts, see the chapter <i>Linking your application</i> .		

STDC CX_LIMITED_RANGE

Syntax	<pre>#pragma STDC CX_LIMITED_RANGE {ON OFF DEFAULT}</pre>	
Parameters	ON	Normal complex mathematic formulas can be used.
	OFF	Normal complex mathematic formulas cannot be used.
	DEFAULT	Sets the default behavior, that is OFF.
Description	Use this pragma directive to specify that the compiler can use the normal complex mathematic formulas for * (multiplication), / (division), and abs. Note: This directive is required by Standard C. The directive is recognized but has no effect in the compiler.	

STDC FENV_ACCESS

Syntax	#pragma STDC FENV	_ACCESS {ON OFF DEFAULT}
Parameters	ON	Source code accesses the floating-point environment. Note that this argument is not supported by the compiler.
	OFF	Source code does not access the floating-point environment.
	DEFAULT	Sets the default behavior, that is OFF.
Description	Use this pragma directive to specify whether your source code accesses the floating-point environment or not.	
	Note: This directive is required by Standard C.	

STDC FP_CONTRACT

	Syntax	<pre>#pragma STDC FP_CONTRACT {ON OFF DEFAULT}</pre>	
	Parameters	ON The compiler is allowed to contract floating-poin expressions.	
		OFF	The compiler is not allowed to contract floating-point expressions. Note that this argument is not supported by the compiler.
		DEFAULT	Sets the default behavior, that is ON.
	Description	Use this pragma directive to specify whether the compiler is allowed to contract floating-point expressions or not. This directive is required by Standard C.	
	Example	<pre>#pragma STDC FP_CONTRACT=ON</pre>	
swi_	number		
	Syntax	<pre>#pragma swi_number=number</pre>	
	Parameters	number	The software interrupt number

Description	Use this pragma directive together with theswi extended keyword. It is used as an argument to the generated SWC assembler instruction, and is used for selecting one software interrupt function in a system containing several such functions.
Example	#pragma swi_number=17
See also	Software interrupts, page 62

type_attribute

Synt	ax	<pre>#pragma type_attribute=type_attribute[,type_attribute,]</pre>	
Para	meters	For information about type attributes that can be used with this pragma directive, see <i>Type attributes</i> , page 313.	
Desc	cription	Use this pragma directive to specify IAR-specific <i>type attributes</i> , which are not part of Standard C. Note however, that a given type attribute might not be applicable to all kind of objects.	
		This directive affects the declaration of the identifier, the next variable, or the next function that follows immediately after the pragma directive.	
Exan	nple	In this example, thumb-mode code is generated for the function foo :	
		<pre>#pragma type_attribute=thumb void foo(void) { }</pre>	
		This declaration, which uses extended keywords, is equivalent:	
		thumb void foo(void) { }	
See	also	The chapter <i>Extended keywords</i> for more information.	
weak			
Synt	ax	<pre>#pragma weak symbol1={symbol2}</pre>	

Parameters

symbol1

A function or variable with external linkage.

	symbol2	A defined function or variable.
Description	This pragma directive ca	in be used in one of two ways:
		n of a function or variable with external linkage a weak ak attribute can also be used for this purpose.
	• To create a weak alias alias for the same fur	s for another function or variable. You can make more than one action or variable.
Example	To make the definition o	f foo a weak definition, write:
	#pragma weak foo	
	To make NMI_Handler	a weak alias for Default_Handler, write:
	#pragma weak NMI_H	andler=Default_Handler
	If NMI_Handler is not on NMI_Handler will reference to the second	lefined elsewhere in the program, all references to to Default_Handler.
See also	weak, page 324.	

Descriptions of pragma directives

Intrinsic functions

This chapter gives reference information about the intrinsic functions, a predefined set of functions available in the compiler.

The intrinsic functions provide direct access to low-level processor operations and can be very useful in, for example, time-critical routines. The intrinsic functions compile into inline code, either as a single instruction or as a short sequence of instructions.

Summary of intrinsic functions

To use intrinsic functions in an application, include the header file intrinsics.h.

To use Neon intrinsic functions in an application, include the header file arm_neon.h. For more information, see *Intrinsic functions for Neon instructions*, page 351.

Note that the intrinsic function names start with double underscores, for example:

__disable_interrupt

This table summarizes the intrinsic functions:

Intrinsic function	Description
CLREX	Inserts a CLREX instruction
CLZ	Inserts a CLZ instruction
disable_fiq	Disables fast interrupt requests (fiq)
disable_interrupt	Disables interrupts
disable_irq	Disables interrupt requests (irq)
DMB	Inserts a DMB instruction
DSB	Inserts a DSB instruction
enable_fiq	Enables fast interrupt requests (fiq)
enable_interrupt	Enables interrupts
enable_irq	Enables interrupt requests (irq)
get_BASEPRI	Returns the value of the Cortex-M3/Cortex-M4 BASEPRI
	register
get_CONTROL	Returns the value of the Cortex-M ${\tt CONTROL}$ register

Table 36: Intrinsic functions summary

Intrinsic function	Description
get_CPSR	Returns the value of the ARM CPSR (Current Program Status Register)
get_FAULTMASK	Returns the value of the Cortex-M3/Cortex-M4 FAULTMASK register
get_FPSCR	Returns the value of FPSCR
get_interrupt_state	Returns the interrupt state
get_IPSR	Returns the value of the IPSR register
get_LR	Returns the value of the link register
get_MSP	Returns the value of the ${\tt MSP}$ register
get_PRIMASK	Returns the value of the Cortex-M ${\tt PRIMASK}$ register
get_PSP	Returns the value of the \mathtt{PSP} register
get_PSR	Returns the value of the PSR register
get_SB	Returns the value of the static base register
get_SP	Returns the value of the stack pointer register
ISB	Inserts an ISB instruction
LDC	Inserts the coprocessor load instruction ${\tt LDC}$
LDCL	Inserts the coprocessor load instruction ${\tt LDCL}$
LDC2	Inserts the coprocessor load instruction ${\tt LDC2}$
LDC2L	Inserts the coprocessor load instruction ${\tt LDC2L}$
LDC_noidx	Inserts the coprocessor load instruction $\ensuremath{\mathtt{LDC}}$
LDCL_noidx	Inserts the coprocessor load instruction ${\tt LDCL}$
LDC2_noidx	Inserts the coprocessor load instruction $\ensuremath{\mathtt{LDC2}}$
LDC2L_noidx	Inserts the coprocessor load instruction ${\tt LDC2L}$
LDREX	Inserts an LDREX instruction
LDREXB	Inserts an LDREXB instruction
LDREXD	Inserts an LDREXD instruction
LDREXH	Inserts an LDREXH instruction
MCR	Inserts the coprocessor write instruction $\ensuremath{\mathbb{MCR}}$
MCR2	Inserts the coprocessor write instruction $\ensuremath{\mathtt{MCR2}}$
MRC	Inserts the coprocessor read instruction $\ensuremath{\mathbb{MRC}}$
MRC2	Inserts the coprocessor read instruction ${\tt MRC2}$
no_operation	Inserts a NOP instruction

ntrinsic function	Description
РКНВТ	Inserts a PKHBT instruction
РКНТВ	Inserts a PKHTB instruction
PLD	Inserts the preload data instruction PLD
PLDW	Inserts the preload data instruction ${\tt PLDW}$
_PLI	Inserts a PLI instruction
QADD	Inserts a QADD instruction
QADD8	Inserts a QADD8 instruction
QADD16	Inserts a QADD16 instruction
QASX	Inserts a QASX instruction
QCFlag	Returns the value of the cumulative saturation flag of the FPSCR register
QDADD	Inserts a QDADD instruction
QDOUBLE	Inserts a QDOUBLE instruction
QDSUB	Inserts a QDSUB instruction
QFlag	Returns the $\ensuremath{\mathbb{Q}}$ flag that indicates if overflow/saturation has occurred
QSAX	Inserts a QSAX instruction
QSUB	Inserts a QSUB instruction
QSUB8	Inserts a QSUB8 instruction
QSUB16	Inserts a QSUB16 instruction
RBIT	Inserts an RBIT instruction
reset_Q_flag	Clears the ${\tt Q}$ flag that indicates if overflow/saturation has occurred
reset_QC_flag	Clears the value of the cumulative saturation flag QC of the FPSCR register
REV	Inserts an REV instruction
REV16	Inserts an REV16 instruction
REVSH	Inserts an REVSH instruction
SADD8	Inserts an SADD8 instruction
SADD16	Inserts an SADD16 instruction
SASX	Inserts an SASX instruction
SEL	Inserts an SEL instruction

Intrinsic function	Description
set_BASEPRI	Sets the value of the Cortex-M3/Cortex-M4 BASEPRI register
set_CONTROL	Sets the value of the Cortex-M $\ensuremath{CONTROL}$ register
set_CPSR	Sets the value of the ARM CPSR (Current Program Status Register)
set_FAULTMASK	Sets the value of the Cortex-M3/Cortex-M4 FAULTMASK register
set_FPSCR	Sets the value of the FPSCR register
set_interrupt_state	Restores the interrupt state
set_LR	Assigns a new address to the link register
set_MSP	Sets the value of the ${\tt MSP}$ register
set_PRIMASK	Sets the value of the Cortex-M PRIMASK register
set_PSP	Sets the value of the \mathtt{PSP} register
set_SB	Assigns a new address to the static base register
set_SP	Assigns a new address to the stack pointer register
SEV	Inserts an SEV instruction
SHADD8	Inserts an SHADD8 instruction
SHADD16	Inserts an SHADD16 instruction
SHASX	Inserts an SHASX instruction
SHSAX	Inserts an SHSAX instruction
SHSUB8	Inserts an SHSUB8 instruction
SHSUB16	Inserts an SHSUB16 instruction
SMLABB	Inserts an SMLABB instruction
SMLABT	Inserts an SMLABT instruction
SMLAD	Inserts an SMLAD instruction
SMLADX	Inserts an SMLADX instruction
SMLALBB	Inserts an SMLALBB instruction
SMLALBT	Inserts an SMLALBT instruction
SMLALD	Inserts an SMLALD instruction
SMLALDX	Inserts an SMLALDX instruction
SMLALTB	Inserts an SMLALTB instruction
SMLALTT	Inserts an SMLALTT instruction
SMLATB	Inserts an SMLATB instruction

Intrinsic function	Description
SMLATT	Inserts an SMLATT instruction
SMLAWB	Inserts an SMLAWB instruction
SMLAWT	Inserts an SMLAWT instruction
SMLSD	Inserts an SMLSD instruction
SMLSDX	Inserts an SMLSDX instruction
SMLSLD	Inserts an SMLSLD instruction
SMLSLDX	Inserts an SMLSLDX instruction
SMMLA	Inserts an SMMLA instruction
SMMLAR	Inserts an SMMLAR instruction
SMMLS	Inserts an SMMLS instruction
SMMLSR	Inserts an SMMLSR instruction
SMMUL	Inserts an SMMUL instruction
SMMULR	Inserts an SMMULR instruction
SMUAD	Inserts an SMUAD instruction
SMUADX	Inserts an SMUADX instruction
SMUL	Inserts a signed 16-bit multiplication
SMULBB	Inserts an SMULBB instruction
SMULBT	Inserts an SMULBT instruction
SMULTB	Inserts an SMULTB instruction
SMULTT	Inserts an SMULTT instruction
SMULWB	Inserts an SMULWB instruction
SMULWT	Inserts an SMULWT instruction
SMUSD	Inserts an SMUSD instruction
SMUSDX	Inserts an SMUSDX instruction
SSAT	Inserts an SSAT instruction
SSAT16	Inserts an SSAT16 instruction
SSAX	Inserts an SSAX instruction
SSUB8	Inserts an SSUB8 instruction
SSUB16	Inserts an SSUB16 instruction
STC	Inserts the coprocessor store instruction $\ensuremath{\operatorname{STC}}$
STCL	Inserts the coprocessor store instruction $\ensuremath{\operatorname{STCL}}$
STC2	Inserts the coprocessor store instruction STC2

Intrinsic function	Description
STC2L	Inserts the coprocessor store instruction ${\tt STC2L}$
STC_noidx	Inserts the coprocessor store instruction $\ensuremath{\operatorname{STC}}$
STCL_noidx	Inserts the coprocessor store instruction $\ensuremath{\mathtt{STCL}}$
STC2_noidx	Inserts the coprocessor store instruction STC2
STC2L_noidx	Inserts the coprocessor store instruction ${\tt STC2L}$
STREX	Inserts a STREX instruction
STREXB	Inserts a STREXB instruction
STREXD	Inserts a STREXD instruction
STREXH	Inserts a STREXH instruction
SWP	Inserts an SWP instruction
SWPB	Inserts an SWPB instruction
SXTAB	Inserts an SXTAB instruction
SXTAB16	Inserts an SXTAB16 instruction
SXTAH	Inserts an SXTAH instruction
SXTB16	Inserts an SXTB16 instruction
UADD8	Inserts a UADD8 instruction
UADD16	Inserts a UADD16 instruction
UASX	Inserts a UASX instruction
UHADD8	Inserts a UHADD8 instruction
UHADD16	Inserts a UHADD16 instruction
UHASX	Inserts a UHASX instruction
UHSAX	Inserts a UHSAX instruction
UHSUB8	Inserts a UHSUB8 instruction
UHSUB16	Inserts a UHSUB16 instruction
UMAAL	Inserts a UMAAL instruction
UQADD8	Inserts a UQADD8 instruction
UQADD16	Inserts a UQADD16 instruction
UQASX	Inserts a UQASX instruction
UQSAX	Inserts a UQSAX instruction
UQSUB8	Inserts a UQSUB8 instruction
UQSUB16	Inserts a UQSUB16 instruction
USAD8	Inserts a USAD8 instruction

Intrinsic function	Description
USADA8	Inserts a USAD8A8 instruction
USAT	Inserts a USAT instruction
USAT16	Inserts a USAT16 instruction
USAX	Inserts a USAX instruction
USUB8	Inserts a USUB8 instruction
USUB16	Inserts a USUB16 instruction
UXTAB	Inserts a UXTAB instruction
UXTAB16	Inserts a UXTAB16 instruction
UXTAH	Inserts a UXTAH instruction
UXTB16	Inserts a UXTB16 instruction
WFE	Inserts a WFE instruction
WFI	Inserts a ${\tt WFI}$ instruction
YIELD	Inserts a YIELD instruction

INTRINSIC FUNCTIONS FOR NEON INSTRUCTIONS

The Neon co-processor implements the Advanced SIMD instruction set extension, as defined by the ARM architecture. To use Neon intrinsic functions in an application, include the header file arm_neon.h. The functions use vector types that are named according to this pattern:

<type><size>x<number_of_lanes>_t

where:

- type is int, unsigned int, float, or poly
- size is 8, 16, 32, or 64
- number_of_lanes is 1, 2, 4, 8, or 16.

The total bit width of a vector type is *size* times *number_of_lanes*, and should fit in a D register (64 bits) or a Q register (128 bits).

For example:

__intrinsic float32x2_t vsub_f32(float32x2_t, float32x2_t);

The intrinsic function vsub_f32 inserts a VSUB.F32 instruction that operates on two 64-bit vectors (D registers), each with two elements (lanes) of 32-bit floating-point type.

Some functions use an array of vector types. As an example, the definition of an array type with four elements of type float32x2_t is:

```
typedef struct
{
  float32x2_t val[4];
}
float32x2x4_t;
```

Descriptions of intrinsic functions

This section gives reference information about each intrinsic function.

Syntax	<pre>voidCLREX(void);</pre>
Description	Inserts a CLREX instruction.
	This intrinsic function requires architecture ARMv6K or ARMv7 for ARM mode, and AVRv7 for Thumb mode.
CLZ	
Syntax	<pre>unsigned charCLZ(unsigned long);</pre>
Description	Inserts a CLZ instruction.
	This intrinsic function requires an ARMv5 architecture or higher for ARM mode, and ARMv6T2 or higher for Thumb mode.
disable_fiq	
Syntax	<pre>voiddisable_fiq(void);</pre>

Disables fast interrupt requests (fiq).

This intrinsic function can only be used in privileged mode and is not available for Cortex-M devices.

Description

__disable_interrupt

Syntax	<pre>voiddisable_interrupt(void);</pre>
Description	Disables interrupts. For Cortex-M devices, it raises the execution priority level to 0 by setting the priority mask bit, PRIMASK. For other devices, it disables interrupt requests (irq) and fast interrupt requests (fiq).
	This intrinsic function can only be used in privileged mode.

__disable_irq

Syntax	<pre>voiddisable_irq(void);</pre>
Description	Disables interrupt requests (irq).
	This intrinsic function can only be used in privileged mode and is not available for Cortex-M devices.

__DMB

Syntax	<pre>voidDMB(void);</pre>
Description	Inserts a DMB instruction. This intrinsic function requires an ARMv6M architecture, or an ARMv7 architecture or higher.

__DSB

Syntax	<pre>voidDSB(void);</pre>
Description	Inserts a DSB instruction. This intrinsic function requires an ARMv6M architecture, or an ARMv7 architecture or higher.

__enable_fiq

Syntax	<pre>voidenable_fig(void);</pre>
Description	Enables fast interrupt requests (fiq).
	This intrinsic function can only be used in privileged mode, and it is not available for Cortex-M devices.

__enable_interrupt

Syntax	<pre>voidenable_interrupt(void);</pre>
Description	Enables interrupts. For Cortex-M devices, it resets the execution priority level to default by clearing the priority mask bit, PRIMASK. For other devices, it enables interrupt requests (irq) and fast interrupt requests (fiq).
	This intrinsic function can only be used in privileged mode.
enable_irq	
Syntax	<pre>voidenable_irq(void);</pre>
Description	Enables interrupt requests (irq).
	This intrinsic function can only be used in privileged mode, and it is not available for Cortex-M devices.
get_BASEPRI	
Syntax	<pre>unsigned longget_BASEPRI(void);</pre>
Description	Returns the value of the BASEPRI register. This intrinsic function can only be used in privileged mode and it requires a Cortex-M3 or Cortex-M4 device.
get_CONTROL	

Syntax	<pre>unsigned longget_CONTROL(void);</pre>
Description	Returns the value of the CONTROL register. This intrinsic function can only be used in privileged mode and it requires a Cortex-M device.
get_CPSR	
Syntax	<pre>unsigned longget_CPSR(void);</pre>
Description	Returns the value of the ARM CPSR (Current Program Status Register). This intrinsic function can only be used in privileged mode, is not available for Cortex-M devices, and it requires ARM mode.

__get_FAULTMASK

Syntax	<pre>unsigned longget_FAULTMASK(void);</pre>
Description	Returns the value of the FAULTMASK register. This intrinsic function can only be used in privileged mode and it requires a Cortex-M3 or Cortex-M4 device.

__get_FPSCR

Syntax	<pre>unsigned longget_FPSCR(void);</pre>
Description	Returns the value of FPSCR (floating-point status and control register).
	This intrinsic function is only available for devices with a VFP coprocessor.

__get_interrupt_state

Syntax	<pre>istate_tget_interrupt_state(void);</pre>
Description	Returns the global interrupt state. The return value can be used as an argument to theset_interrupt_state intrinsic function, which will restore the interrupt state.
	This intrinsic function can only be used in privileged mode, and cannot be used when using theaeabi compiler option.
Example	#include "intrinsics.h"
	<pre>void CriticalFn() { istate_t s =get_interrupt_state(); disable_interrupt(); /* Do something here. */ set_interrupt_state(s);</pre>
	set_interrupt_state(s); }
	The advantage of using this sequence of code compared to using

I he advantage of using this sequence of code compared to using ___disable_interrupt and __enable_interrupt is that the code in this example will not enable any interrupts disabled before the call of __get_interrupt_state.

get_IPSR	
Syntax	<pre>unsigned longget_IPSR(void);</pre>
Description	Returns the value of the IPSR register (Interrupt Program Status Register). This intrinsic function can only be used in privileged mode, and is only available for Cortex-M devices.
get_LR	
Syntax	<pre>unsigned longget_LR(void);</pre>
Description	Returns the value of the link register (R14).
get_MSP	
Syntax	<pre>unsigned longget_MSP(void);</pre>
Description	Returns the value of the MSP register (Main Stack Pointer). This intrinsic function can only be used in privileged mode, and is only available for Cortex-M devices.
get_PRIMASK	
Syntax	<pre>unsigned longget_PRIMASK(void);</pre>
Description	Returns the value of the PRIMASK register. This intrinsic function can only be used in privileged mode and it requires a Cortex-M device.
get_PSP	
Syntax	<pre>unsigned longget_PSP(void);</pre>
Description	Returns the value of the PSP register (Process Stack Pointer). This intrinsic function can only be used in privileged mode, and is only available for Cortex-M devices.

get_PSR	
Syntax	<pre>unsigned longget_PSR(void);</pre>
Description	Returns the value of the PSR register (combined Program Status Register). This intrinsic function can only be used in privileged mode, and is only available for Cortex-M devices.
get_SB	
Syntax	<pre>unsigned longget_SB(void);</pre>
Description	Returns the value of the static base register (R9).
get_SP	
Syntax	<pre>unsigned longget_SP(void);</pre>
Description	Returns the value of the stack pointer register (R13).
ISB	
Syntax	<pre>voidISB(void);</pre>
Description	Inserts an ISB instruction. This intrinsic function requires an ARMv6M architecture, or an ARMv7 architecture or higher.
LDC	
LDCL	
LDC2	
LDC2L	
Syntax	<pre>voidLDCxxx(ul coproc,ul CRn,ul const *src);</pre>

Parameters		
	coproc	The coprocessor number 015.
	CRn	The coprocessor register to load.
	SIC	A pointer to the data to load.
Description	Inserts the coprocessor load instruction LDC—or one of its variants—which r a value will be loaded into a coprocessor register. The parameters <i>coproc</i> an be encoded in the instruction and must therefore be constants. The intrinsic functionsLDC andLDCL require architecture ARMv4 or h ARM mode, and ARMv6T2 or higher for Thumb mode.	
		nctionsLDC2 andLDC2L require architecture ARMv5 or higher for d ARMv6T2 or higher for Thumb mode.
LDC_noidx		
LDCL_noidx		

__LDC2_noidx

__LDC2L_noidx

Syntax	<pre>voidLDCxxx_noidx(ul coproc,ul CRn,ul const *src,ul option);</pre>	
Parameters	coproc	The coprocessor number 015.
	CRn	The coprocessor register to load.
	SIC	A pointer to the data to load.
	option	Additional coprocessor option 0255.
Description	Inserts the coprocessor load instruction LDC, or one of its variants. A value will be loaded into a coprocessor register. The parameters <i>coproc</i> , <i>CRn</i> , and <i>option</i> will be encoded in the instruction and must therefore be constants. The intrinsic functionsLDC_noidx andLDCL_noidx require architecture ARMv or higher for ARM mode, and ARMv6T2 or higher for Thumb mode.	

The intrinsic functions __LDC2_noidx and __LDC2L_noidx require architecture ARMv5 or higher for ARM mode, and ARMv6T2 or higher for Thumb mode.

L	D	R	E)	(

__LDREXB

__LDREXH

Syntax	<pre>unsigned longLDREX(unsigned long *); unsigned charLDREXB(unsigned char *); unsigned long longLDREXD(unsigned long long *); unsigned shortLDREXH(unsigned short *);</pre>		
Description	Inserts the specified instruction.		
	TheLDREX intrinsic function requires architecture ARMv6 or higher for ARM mode, and ARMv6T2 or higher for Thumb mode.		
	TheLDREXB and theLDREXH intrinsic functions require architecture ARMv6K or ARMv7 for ARM mode, and ARMv7 for Thumb mode.		
	TheLDREX intrinsic function requires architecture ARMv6K or ARMv7 for ARM mode, and ARMv7 but not ARMv7-M for Thumb mode.		

__MCR

__MCR2

Syntax	CRm,ul op voidMCR2(<pre>voidMCR(ul coproc,ul opcode_1,ul src,ul CRn,ul CRm,ul opcode_2); voidMCR2(ul coproc,ul opcode_1,ul src,ul CRn,ul CRm,ul opcode_2);</pre>	
Parameters	coproc	The coprocessor number 015.	
	opcode_1	Coprocessor-specific operation code.	

	src	The value to be written to the coprocessor.
	CRn	The coprocessor register to write to.
	CRm	Additional coprocessor register; set to zero if not used.
	opcode_2	Additional coprocessor-specific operation code; set to zero if not used.
Description	coprocessor register.	write instruction (MCR or MCR2). A value will be written to a The parameters <i>coproc</i> , <i>opcode_1</i> , <i>CRn</i> , <i>CRm</i> , and <i>opcode_2</i> will struction and must therefore be constants.
	The intrinsic functior Thumb mode.	nMCR requires either ARM mode, or an ARMv6T2 or higher for
		nMCR2 requires an ARMv5T architecture or higher for ARM or higher for Thumb mode.
MRC		
MRC2		
Syntax	unsigned longMRC(ul coproc,ul opcode_1,ul CRn,ul CRm,ul opcode_2); unsigned longMRC2(ul coproc,ul opcode_1,ul CRn,ul CRm,ul opcode_2);	
Parameters		
	coproc	The coprocessor number 015.
	opcode_1	Coprocessor-specific operation code.
	CRn	The coprocessor register to write to.
	CRm	Additional coprocessor register; set to zero if not used.
	opcode_2	Additional coprocessor-specific operation code; set to zero if not used.
Description	iption Inserts a coprocessor read instruction (MRC or MRC2). Returns the value of the specified coprocessor register. The parameters <i>coproc</i> , <i>opcode_1</i> , <i>CRn</i> , <i>CRm</i> , and <i>opcode_2</i> will be encoded in the instruction and must therefore be constants. The intrinsic functionMRC requires either ARM mode, or an ARMv6T2 or higher for Thumb mode.	

The intrinsic function __MRC2 requires an ARMv5T architecture or higher for ARM mode, or ARMv6T2 or higher for Thumb mode.

__no_operation

Syntax	voidno_opera	ation(void);
Description	Inserts a NOP instruc	ction.
РКНВТ		
Syntax	unsigned long _ long <i>count</i>);	_PKHBT(unsigned long x , unsigned long y , unsigned
Parameters	x	First operand.
	У	Second operand, optionally shifted left.
	count	Shift count 0–31, where 0 means no shift.
Description	Inserts a PKHBT instruction, with an optionally shifted operand (LSL) for count in the range 1–31.	
		on requires an ARM v6 architecture or higher for ARM mode, and 7-R, or ARMv7E-M for Thumb mode.
РКНТВ		
Syntax	unsigned long _ long <i>count</i>);	_PKHTB(unsigned long x , unsigned long y , unsigned
Parameters		
	x	First operand.
	У	Second operand, optionally shifted right (arithmetic shift).
	count	Shift count 0–32, where 0 means no shift.
Description	Inserts a PKHTB inst range 1–32.	ruction, with an optionally shifted operand (ASR) for count in the

This intrinsic function requires an ARM v6 architecture or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARM v7E-M for Thumb mode.

__PLD

__PLDW

Syntax	<pre>voidPLD(void const *); voidPLDW(void const *);</pre>
Description	Inserts a preload data instruction (PLD or PLDW).
	The intrinsic functionPLD requires an ARMv7 architecturePLDW requires an ARMv7 architecture with MP extensions (for example Cortex-A5).

__PLI

Syntax	<pre>voidPLI(void const *);</pre>
Description	Inserts a PLI instruction.
	This intrinsic function requires an ARM v7 architecture.

__QADD

__QDADD

_QDSUB

__QSUB

Syntax	<pre>signed longQxxx(signed long, signed long);</pre>
Description	Inserts the specified instruction.
	These intrinsic functions require architecture ARMv5E or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.

QADD8	
QADD16	
QASX	
QSAX	
QSUB8	
QSUB16	
Syntax	unsigned longQxxx(unsigned long, unsigned long);
Description	Inserts the specified instruction.
	These intrinsic functions require architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.
QCFlag	
Syntax	<pre>unsigned longQCFlag(void);</pre>
Description	Returns the value of the cumulative saturation flag QC of the FPSCR register (Floating-point Status and Control Register). This intrinsic function is only available for devices with Neon (Advanced SIMD).
QDOUBLE	
Syntax	<pre>signed longQDOUBLE(signed long);</pre>
Description	Inserts an instruction <code>QADD Rd</code> , <code>Rs</code> , <code>Rs</code> for a source register <code>Rs</code> , and a destination register <code>Rd</code> .
	This intrinsic function requires architecture ARMv5E or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.

Syntax	<pre>intQFlag(void);</pre>
Description	Returns the Q flag that indicates if overflow/saturation has occurred.
	This intrinsic function requires architecture ARMv5E or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.

__RBIT

Syntax	<pre>unsigned longRBIT(unsigned long);</pre>
Description	Inserts an RBIT instruction, which reverses the bit order in a 32-bit register.
	This intrinsic function requires architecture ARMv6T2 or higher.

__reset_Q_flag

Syntax	<pre>voidreset_Q_flag(void);</pre>
Description	Clears the Q flag that indicates if overflow/saturation has occurred.
	This intrinsic function requires an ARM v5E architecture or higher for ARM mode, and ARM v7A, ARM v7R, or ARM v7E-M for Thumb mode.

__reset_QC_flag

Syntax	<pre>voidreset_QC_flag(void);</pre>
Description	Clears the value of the cumulative saturation flag QC of the FPSCR register (Floating-point Status and Control Register). This intrinsic function is only available for devices with Neon (Advanced SIMD).

__REV

__REV16

__REVSH

Syntax	<pre>unsigned longREV(unsigned long); unsigned longREV16(unsigned long); signed longREVSH(short);</pre>
Description	Inserts the specified instruction.
	These intrinsic functions require architecture ARMv6 or higher.

__SADD8

__SADD16

__SASX

__SSAX

__SSUB8

__SSUB16

Syntax	<pre>unsigned longSxxx(unsigned long, unsigned long);</pre>
Description	Inserts the specified instruction.
	These intrinsic functions require architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.

__SEL

Syntax	unsigned longSEL(unsigned long, unsigned long);
Description	Inserts an SEL instruction.

This intrinsic function requires architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.

__set_BASEPRI

Syntax	<pre>voidset_BASEPRI(unsigned long);</pre>
Description	Sets the value of the BASEPRI register. This intrinsic function can only be used in privileged mode and it requires a Cortex-M3 or Cortex-M4 device.

__set_CONTROL

Syntax	<pre>voidset_CONTROL(unsigned long);</pre>
Description	Sets the value of the CONTROL register. This intrinsic function can only be used in privileged mode and it requires a Cortex-M device.

___set_CPSR

Syntax	<pre>voidset_CPSR(unsigned long);</pre>
Description	Sets the value of the ARM CPSR (Current Program Status Register). Only the control field is changed (bits 0-7). This intrinsic function can only be used in privileged mode, is not available for Cortex-M devices, and it requires ARM mode.

__set_FAULTMASK

Syntax	<pre>voidset_FAULTMASK(unsigned long);</pre>
Description	Sets the value of the FAULTMASK register. This intrinsic function can only be used in privileged mode and it requires a Cortex-M3 or Cortex-M4 device.
set_FPSCR	
Syntax	<pre>voidset_FPSCR(unsigned long);</pre>
Description	Sets the value of FPSCR (floating-point status and control register)

This intrinsic function is only available for devices with a VFP coprocessor.

__set_interrupt_state

Syntax	<pre>voidset_interrupt_state(istate_t);</pre>
Description	Restores the interrupt state to a value previously returned by theget_interrupt_state function.
	For information about theistate_t type, seeget_interrupt_state, page 355.

___set_LR

Syntax	<pre>voidset_LR(unsigned long);</pre>
Description	Assigns a new address to the link register (R14).
set_MSP	
Syntax	<pre>voidset_MSP(unsigned long);</pre>
Description	Sets the value of the MSP register (Main Stack Pointer). This intrinsic function can only be used in privileged mode, and is only available for Cortex-M devices.

__set_PRIMASK

Syntax	<pre>voidset_PRIMASK(unsigned long);</pre>
Description	Sets the value of the PRIMASK register. This intrinsic function can only be used in privileged mode and it requires a Cortex-M device.

__set_PSP

Syntax	<pre>voidset_PSP(unsigned long);</pre>
Description	Sets the value of the PSP register (Process Stack Pointer). This intrinsic function can only be used in privileged mode, and is only available for Cortex-M devices.

__set_SB

Syntax	<pre>voidset_SB(unsigned long);</pre>
Description	Assigns a new address to the static base register (R9).
set_SP	
Syntax	<pre>voidset_SP(unsigned long);</pre>
Description	Assigns a new address to the stack pointer register (R13).
SEV	
Syntax	<pre>voidSEV(void);</pre>
Description	Inserts an SEV instruction.
	This intrinsic function requires architecture ARMv7 for ARM mode, and ARMv6-M or ARMv7 for Thumb mode.
SHADD8	
SHADD16	
SHASX	
SHSAX	
SHSUB8	
SHSUB16	
Syntax	unsigned longSHxxx(unsigned long, unsigned long);
Description	Inserts the specified instruction.

	These intrinsic functions require architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.
SMLABB	
SMLABT	
SMLATB	
SMLATT	
SMLAWB	
SMLAWT	
Syntax	unsigned longSMLAxxx(unsigned long, unsigned long, unsigned long);
Description	Inserts the specified instruction.
	These intrinsic functions require architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.
SMLAD	
SMLADX	
SMLSD	
SMLSDX	
Syntax	unsigned longSMLxxx(unsigned long, unsigned long, unsigned long);

Description Inserts the specified instruction.

These intrinsic functions require architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.

SMLALBB	
SMLALBT	
SMLALTB	
SMLALTT	
Syntax	unsigned long long \SMLALxxx (unsigned long, unsigned long, unsigned long);
Description	Inserts the specified instruction.
	These intrinsic functions require architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.
SMLALD	
SMLALDX	
SMLSLD	
SMLSLDX	
Syntax	unsigned long longSMLxxx(unsigned long, unsigned long, unsigned long long);
Description	Inserts the specified instruction.
	These intrinsic functions require architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.

__SMMLA

__SMMLAR

__SMMLS

___SMMLSR

Syntax	unsigned longSMMLxxx(unsigned long, unsigned long, unsigned long);
Description	Inserts the specified instruction.
	These intrinsic functions require architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.

__SMMUL

__SMMULR

Syntax	unsigned longSMMULxxx(unsigned long, unsigned long);
Description	Inserts the specified instruction.
	These intrinsic functions require architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.

__SMUAD

__SMUADX

__SMUSD

___SMUSDX

Syntax	<pre>unsigned longSMUxxx(unsigned long, unsigned long);</pre>
Description	Inserts the specified instruction.

	These intrinsic functions require architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.
SMUL	
Syntax	signed longSMUL(signed short, signed short);
Description	Inserts a signed 16-bit multiplication.
	This intrinsic function requires architecture ARMv5-E or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.
SMULBB	
SMULBT	
SMULTB	
SMULTT	
SMULWB	
Syntax	<pre>unsigned longSMULxxx(unsigned long, unsigned long);</pre>
Description	Inserts the specified instruction.
	These intrinsic functions require architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.
SSAT	
Syntax	<pre>unsigned longSSAT(unsigned long, unsigned long);</pre>
Description	Inserts an SSAT instruction.

The compiler will incorporate a shift instruction into the operand when possible. For example, $__SSAT(x << 3, 11)$ compiles to SSAT Rd, #11, Rn, LSL #3, where the value of x has been placed in register Rn and the return value of $__SSAT$ will be placed in register Rd.

This intrinsic function requires architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7-M for Thumb mode.

_SSAT16

Syntax	<pre>unsigned longSSAT16(unsigned long, unsigned long);</pre>
Description	Inserts an SSAT16 instruction.
	This intrinsic function requires architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARM v7E-M for Thumb mode.

__STC

__STC2

__STC2L

Syntax	voidSTCxxx(_	_ul coproc,ul CRn,ul const *dst);
Parameters		
	coproc	The coprocessor number 015.
	CRn	The coprocessor register to load.
	dst	A pointer to the destination.
Description	Inserts the coprocessor store instruction STC—or one of its variants—which means that the value of the specified coprocessor register will be written to a memory location. The parameters <i>coproc</i> and <i>CRn</i> will be encoded in the instruction and must therefore be constants.	
		onsSTC andSTCL require architecture ARMv4 or higher for RM v6T2 or higher for Thumb mode.

The intrinsic functions $__$ STC2 and $__$ STC2L require architecture ARMv5 or higher for ARM mode, and ARMv6-T2 or higher for Thumb mode.

__STC_noidx

__STCL_noidx

__STC2_noidx

__STC2L_noidx

Syntax	<pre>voidSTCxxx_no option);</pre>	<pre>pidx(ul coproc,ul CRn,ul const *dst,ul</pre>
Parameters		
	coproc	The coprocessor number 015.
	CRn	The coprocessor register to load.
	dst	A pointer to the destination.
	option	Additional coprocessor option 0255.
Description	Inserts the coprocessor store instruction STC—or one of its variants—which means that the value of the specified coprocessor register will be written to a memory location. The parameters <i>coproc</i> , <i>CRn</i> , and <i>option</i> will be encoded in the instruction and must therefore be constants.	
		nsSTC_noidx andSTCL_noidx require architecture r ARM mode, and ARMv6-T2 or higher for Thumb mode.
		nsSTC2_noidx andSTC2L_noidx require architecture r ARM mode, and ARMv6-T2 or higher for Thumb mode.

__STREX

__STREXB

__STREXD

__STREXH

Syntax	<pre>unsigned longSTREX(unsigned long, unsigned long *); unsigned longSTREXB(unsigned char, unsigned char *); unsigned longSTREXD(unsigned long long, unsigned long long*); unsigned longSTREXH(unsigned short, unsigned short *);</pre>
Description	Inserts the specified instruction.
	TheSTREX intrinsic function requires architecture ARMv6 or higher for ARM mode, and ARMv6-T2 or higher for Thumb mode.
	TheSTREXB and theSTREXH intrinsic functions require architecture ARMv6K or ARMv7 for ARM mode, and ARMv7 for Thumb mode.
	TheSTREXD intrinsic function requires architecture ARMv6K or ARMv7 for ARM mode, and ARMv7 except for ARMv7-M for Thumb mode.
_SWP	
_SWPB	
Syntax	<pre>unsigned longSWP(unsigned long, unsigned long *); charSWPB(unsigned char, unsigned char *);</pre>

Description Inserts the specified instruction.

These intrinsic functions require ARM mode.

__SXTAB

__SXTAB16

__SXTAH

__SXTB16

Syntax	<pre>unsigned longSXTxxx(unsigned long, unsigned long);</pre>
Description	Inserts the specified instruction.
	These intrinsic functions require architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.

__UADD8

__UADD16

__UASX

__USAX

__USUB8

__USUB16

Syntax	unsigned longUxxx(unsigned long, unsigned long);	
Description	Inserts the specified instruction.	
	These intrinsic functions require architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.	

__UHADD16

__UHASX

__UHSAX

__UHSUB8

__UHSUB16

Syntax	<pre>unsigned longUHxxx(unsigned long, unsigned long);</pre>
Description	Inserts the specified instruction.
	These intrinsic functions require architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.

__UMAAL

Syntax	unsigned long longUMAAL(unsigned long, unsigned long, unsigned long, unsigned long);
Description	Inserts an UMAAL instruction.
	This intrinsic function requires architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.

__UQADD8

__UQADD16

__UQASX

__UQSAX

__UQSUB8

__UQSUB16

Syntax	unsigned longUQxxx(unsigned long, unsigned long);
Description	Inserts the specified instruction.
	These intrinsic functions require architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.

_USAD8

_USADA8

Syntax	unsigned longUSADxxx(unsigned long, unsigned long);
Description	Inserts the specified instruction.
	These intrinsic functions require architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.
USAT	
Syntax	<pre>unsigned longUSAT(unsigned long, unsigned long);</pre>
Description	Inserts a USAT instruction.

The compiler will incorporate a shift instruction into the operand when possible. For example, __USAT(x << 3,11) compiles to USAT Rd, #11, Rn, LSL #3, where the value of x has been placed in register Rn and the return value of __USAT will be placed in register Rd.

This intrinsic function requires architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7-M for Thumb mode.

_USAT16

Syntax	unsigned longUSAT16(unsigned long, unsigned long);
Description	Inserts a USAT16 instruction.
	This intrinsic function requires architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.

__UXTAB

__UXTABI6

__UXTAH

__UXTB16

Syntax	unsigned longUXTxxx(unsigned long, unsigned long);
Description	Inserts the specified instruction.
	These intrinsic functions require architecture ARMv6 or higher for ARM mode, and ARMv7-A, ARMv7-R, or ARMv7E-M for Thumb mode.

__WFI

__WFE

__YIELD

 Syntax
 void long __xxx(void);

 Description
 Inserts the specified instruction.

 These intrinsic functions require architecture ARMv7 for ARM mode, and ARMv6-M, or ARMv7 for Thumb mode.

The preprocessor

This chapter gives a brief overview of the preprocessor, including reference information about the different preprocessor directives, symbols, and other related information.

Overview of the preprocessor

The preprocessor of the IAR C/C++ Compiler for ARM adheres to Standard C. The compiler also makes these preprocessor-related features available to you:

• Predefined preprocessor symbols

These symbols allow you to inspect the compile-time environment, for example the time and date of compilation. For more information, see *Description of predefined preprocessor symbols*, page 382.

• User-defined preprocessor symbols defined using a compiler option

In addition to defining your own preprocessor symbols using the #define directive, you can also use the option -D, see -D, page 237.

Preprocessor extensions

There are several preprocessor extensions, for example many pragma directives; for more information, see the chapter *Pragma directives*. For information about the corresponding _Pragma operator and the other extensions related to the preprocessor, see *Descriptions of miscellaneous preprocessor extensions*, page 387.

Preprocessor output

Use the option --preprocess to direct preprocessor output to a named file, see --preprocess, page 260.

To specify a path for an include file, use forward slashes:

#include "mydirectory/myfile"

In source code, use forward slashes:

file = fopen("mydirectory/myfile","rt");

Note that backslashes can also be used. In this case, use one in include file paths and two in source code strings.

Description of predefined preprocessor symbols

This section lists and describes the preprocessor symbols.

AAPCS Description An integer that is set based on the --aapcs option. The symbol is set to 1 if the AAPCS base standard is the selected calling convention (--aapcs=std). The symbol is undefined for other calling conventions. See also --aapcs, page 233. AAPCS VFP Description An integer that is set based on the --aapcs option. The symbol is set to 1 if the VFP variant of AAPCS is the selected calling convention (--aapcs=vfp). The symbol is undefined for other calling conventions. See also --aapcs, page 233. ARM_ADVANCED_SIMD__ Description An integer that is set based on the --cpu option. The symbol is set to 1 if the selected processor architecture has the Advanced SIMD architecture extension. The symbol is undefined for other cores. See also --cpu, page 235. ARM MEDIA Description An integer that is set based on the --cpu option. The symbol is set to 1 if the selected processor architecture has the ARMv6 SIMD extensions for multimedia. The symbol is undefined for other cores. See also --cpu, page 235.

__ARM_PROFILE_M__

Description	An integer that is set based on thecpu option. The symbol is set to 1 if the selected processor architecture is a profile M core. The symbol is undefined for other cores.
See also	<i>cpu</i> , page 235.
ARMVFP	
Description	An integer that reflects thefpu option and is defined toARMVFPV2, ARMVFPV3, orARMVFPV4 These symbolic names can be used when testing theARMVFP symbol. If VFP code generation is disabled (default), the symbol will be undefined.
See also	<i>fpu</i> , page 246.
ARMVFP_D16	
-	

Description	An integer that is set based on the $-fpu$ option. The symbol is set to 1 if the selected FPU is a VFPv3 or VFPv4 unit with only 16 D registers. Otherwise, the symbol is undefined.
See also	<i>fpu</i> , page 246.

__ARMVFP_FP16__

Description	An integer that is set based on the $-fpu$ option. The symbol is set to 1 if the selected FPU only supports 16-bit floating-point numbers. Otherwise, the symbol is undefined.
See also	<i>fpu</i> , page 246.

__ARMVFP_SP__

Description	An integer that is set based on thefpu option. The symbol is set to 1 if the selected FPU only supports 16-bit single-precision. Otherwise, the symbol is undefined.
See also	<i>fpu</i> , page 246.



Description	A string that identifies the name of the base source file (that is, not the header file), being compiled.
See also	See also <i>FILE</i> , page 385, and <i>no_path_in_file_macros</i> , page 254.
BUILD_NUMBER	
Description	A unique integer that identifies the build number of the compiler currently in use.
CORE	
Description	An integer that identifies the processor architecture in use. The symbol reflects thecore option and is defined toARM4TM,ARM5,ARM5E,ARM6,ARM6M,ARM6M,ARM7M,ARM7EM,ARM7A, orARM7R These symbolic names can be used when testing theCORE symbol.
cplusplus	
Description	An integer which is defined when the compiler runs in any of the C++ modes, otherwise it is undefined. When defined, its value is 199711L. This symbol can be used with #ifdef to detect whether the compiler accepts C++ code. It is particularly useful when creating header files that are to be shared by C and C++ code.*
	This symbol is required by Standard C.
CPU_MODE	
Description	An integer that reflects the selected CPU mode and is defined to 1 for Thumb and 2 for ARM.
DATE	
Description	A string that identifies the date of compilation, which is returned in the form "Mmm dd $yyyy$ ", for example "Oct 30 2010"
	This symbol is required by Standard C.

DOUBLE	
Description	An integer that identifies the size of the data type double. The symbol is defined to 64.
embedded_cplu	splus
Description	An integer which is defined to 1 when the compiler runs in Embedded C++ or Extended Embedded C++ mode, otherwise the symbol is undefined. This symbol can be used with $\#ifdef$ to detect whether the compiler accepts C++ code. It is particularly useful when creating header files that are to be shared by C and C++ code.
	This symbol is required by Standard C.
FILE	
Description	A string that identifies the name of the file being compiled, which can be both the base source file and any included header file.
	This symbol is required by Standard C.
See also	See alsoBASE_FILE, page 384, andno_path_in_file_macros, page 254.
func	
Description	A predefined string identifier that is initialized with the name of the function in which the symbol is used. This is useful for assertions and other trace utilities. The symbol requires that language extensions are enabled.
	This symbol is required by Standard C.
See also	-e, page 242. See alsoPRETTY_FUNCTION, page 386.
FUNCTION	
Description	A predefined string identifier that is initialized with the name of the function in which the symbol is used. This is useful for assertions and other trace utilities. The symbol requires that language extensions are enabled.

See also -e, page 242. See also __PRETTY_FUNCTION__, page 386.

IAR_SYSTEMS_ICC			
Description	An integer that identifies the IAR compiler platform. The current value is 8. Note that the number could be higher in a future version of the product. This symbol can be tested with #ifdef to detect whether the code was compiled by a compiler from IAR Systems.		
ICCARM			
Description	An integer that is set to 1 when the code is compiled with the IAR C/C++ Compiler for ARM.		
LINE			
Description	An integer that identifies the current source line number of the file being compiled, which can be both the base source file and any included header file.		
	This symbol is required by Standard C.		
LITTLE_ENDIAN_	LITTLE_ENDIAN		
Description	An integer that reflects theendian option and is defined to 1 when the byte order is little-endian. The symbol is defined to 0 when the byte order is big-endian.		
PRETTY_FUNCTI	ON		
Description	A predefined string identifier that is initialized with the function name, including parameter types and return type, of the function in which the symbol is used, for example "void func(char)". This symbol is useful for assertions and other trace utilities. The symbol requires that language extensions are enabled.		
See also	<i>-e</i> , page 242. See also <i>func</i> , page 385.		
ROPI			
Description	An integer that is defined when theropi compiler option is used.		
See also	<i>ropi</i> , page 262.		

__RWPI__

Description	An integer that is defined when therwpi compiler option is used.
See also	<i>rwpi</i> , page 263.

__STDC__

Description	An integer that is set to 1, which means the compiler adheres to Standard C. This symbol
	can be tested with $\#ifdef$ to detect whether the compiler in use adheres to Standard C.*
	This symbol is required by Standard C.

__STDC_VERSION__

Description	An integer that identifies the version of the C standard in use. The symbol expands to 199901L, unless the $c89$ compiler option is used in which case the symbol expands to 199409L. This symbol does not apply in EC++ mode.
	This symbol is required by Standard C.
TIME	
Description	A string that identifies the time of compilation in the form "hh:mm:ss".
	This symbol is required by Standard C.
VER	
Description	An integer that identifies the version number of the IAR compiler in use. For example, version 5.11.3 is returned as 5011003.

Descriptions of miscellaneous preprocessor extensions

This section gives reference information about the preprocessor extensions that are available in addition to the predefined symbols, pragma directives, and Standard C directives.

NDEBUG

Description	This preprocessor symbol determines whether any assert macros you have written in your application shall be included or not in the built application.
	If this symbol is not defined, all assert macros are evaluated. If the symbol is defined, all assert macros are excluded from the compilation. In other words, if the symbol is:
	• defined , the assert code will <i>not</i> be included
	• not defined, the assert code will be included
	This means that if you write any assert code and build your application, you should define this symbol to exclude the assert code from the final application.
	Note that the assert macro is defined in the assert.h standard include file.
See also	Assert, page 130.
	In the IDE, the NDEBUG symbol is automatically defined if you build your application in the Release build configuration.
#warning message	
C	

Syntax	#warning message
Description	where <i>message</i> can be any string. Use this preprocessor directive to produce messages. Typically, this is useful for
Description	assertions and other trace utilities, similar to the way the Standard C #error directive is used. This directive is not recognized when thestrict compiler option is used.

Library functions

This chapter gives an introduction to the C and C++ library functions. It also lists the header files used for accessing library definitions.

For detailed reference information about the library functions, see the online help system.

Library overview

The compiler comes with the IAR DLIB Library, a complete library, compliant with Standard C and C++. This library also supports floating-point numbers in IEEE 754 format and it can be configured to include different levels of support for locale, file descriptors, multibyte characters, etc.

For more information about customization, see the chapter *The DLIB runtime environment*.

For detailed information about the library functions, see the online documentation supplied with the product. There is also keyword reference information for the DLIB library functions. To obtain reference information for a function, select the function name in the editor window and press F1.

For more information about library functions, see the chapter *Implementation-defined* behavior for Standard C in this guide.

HEADER FILES

Your application program gains access to library definitions through header files, which it incorporates using the #include directive. The definitions are divided into several different header files, each covering a particular functional area, letting you include just those that are required.

It is essential to include the appropriate header file before making any reference to its definitions. Failure to do so can cause the call to fail during execution, or generate error or warning messages at compile time or link time.

LIBRARY OBJECT FILES

Most of the library definitions can be used without modification, that is, directly from the library object files that are supplied with the product. For information about how to set up a runtime library, see *Setting up the runtime environment*, page 100. The linker will include only those routines that are required—directly or indirectly—by your application.

ALTERNATIVE MORE ACCURATE LIBRARY FUNCTIONS

The default implementation of cos, sin, tan, and pow is designed to be fast and small. As an alternative, there are versions designed to provide better accuracy. They are named __iar_xxx_accuratef for float variants of the functions and __iar_xxx_accuratel for long double variants of the functions, and where xxx is cos, sin, etc.

To use any of these more accurate versions, use the --redirect linker option.

REENTRANCY

A function that can be simultaneously invoked in the main application and in any number of interrupts is reentrant. A library function that uses statically allocated data is therefore not reentrant.

Most parts of the DLIB library are reentrant, but the following functions and parts are not reentrant because they need static data:

- Heap functions—malloc, free, realloc, calloc, and the C++ operators new and delete
- Locale functions—localeconv, setlocale
- Multibyte functions—mbrlen, mbrtowc, mbsrtowc, mbtowc, wcrtomb, wcsrtomb, wctomb
- Rand functions—rand, srand
- Time functions—asctime, localtime, gmtime, mktime
- The miscellaneous functions atexit, strerror, strtok
- Functions that use files or the heap in some way. This includes printf, sprintf, scanf, sscanf, getchar, and putchar.

Functions that can set errno are not reentrant, because an errno value resulting from one of these functions can be destroyed by a subsequent use of the function before it is read. This applies to math and string conversion functions, among others.

Remedies for this are:

• Do not use non-reentrant functions in interrupt service routines

• Guard calls to a non-reentrant function by a mutex, or a secure region, etc.

THE LONGJMP FUNCTION

A longjmp is in effect a jump to a previously defined setjmp. Any variable length arrays or C++ objects residing on the stack during stack unwinding will not be destroyed. This can lead to resource leaks or incorrect application behavior.

IAR DLIB Library

The IAR DLIB Library provides most of the important C and C++ library definitions that apply to embedded systems. These are of the following types:

- Adherence to a free-standing implementation of Standard C. The library supports most of the hosted functionality, but you must implement some of its base functionality. For additional information, see the chapter *Implementation-defined behavior for Standard C* in this guide.
- Standard C library definitions, for user programs.
- C++ library definitions, for user programs.
- CSTARTUP, the module containing the start-up code, see the chapter *The DLIB runtime environment* in this guide.
- Runtime support libraries; for example low-level floating-point routines.
- Intrinsic functions, allowing low-level use of ARM features. See the chapter *Intrinsic functions* for more information.

In addition, the IAR DLIB Library includes some added C functionality, see *Added C functionality*, page 395.

C HEADER FILES

This section lists the header files specific to the DLIB library C definitions. Header files may additionally contain target-specific definitions; these are documented in the chapter *Using C*.

This table lists the C header files:

Header file	Usage
assert.h	Enforcing assertions when functions execute
complex.h	Computing common complex mathematical functions
ctype.h	Classifying characters
errno.h	Testing error codes reported by library functions

Table 37: Traditional Standard C header files-DLIB

Header file	Usage
fenv.h	Floating-point exception flags
float.h	Testing floating-point type properties
inttypes.h	Defining formatters for all types defined in ${\tt stdint.h}$
iso646.h	Using Amendment I—iso646.h standard header
limits.h	Testing integer type properties
locale.h	Adapting to different cultural conventions
math.h	Computing common mathematical functions
setjmp.h	Executing non-local goto statements
signal.h	Controlling various exceptional conditions
stdarg.h	Accessing a varying number of arguments
stdbool.h	Adds support for the bool data type in C.
stddef.h	Defining several useful types and macros
stdint.h	Providing integer characteristics
stdio.h	Performing input and output
stdlib.h	Performing a variety of operations
string.h	Manipulating several kinds of strings
tgmath.h	Type-generic mathematical functions
time.h	Converting between various time and date formats
uchar.h	Unicode functionality (IAR extension to Standard C)
wchar.h	Support for wide characters
wctype.h	Classifying wide characters

Table 37: Traditional Standard C header files—DLIB (Continued)

C++ HEADER FILES

This section lists the C++ header files:

• The C++ library header files

The header files that constitute the Standard C++ and the Embedded C++ library.

• The C++ standard template library (STL) header files

The header files that constitute STL for the Standard C++ and the Extended Embedded C++ library.

• The C++ C header files

The C++ header files that provide the resources from the C library.

The C++ library header files

This table lists the header files that can be used in C++ as well as in Embedded C++:

Header file	Usage
complex	Defining a class that supports complex arithmetic
exception	Defining several functions that control exception handling; only usable in C++
fstream	Defining several I/O stream classes that manipulate external files
iomanip	Declaring several I/O stream manipulators that take an argument
ios	Defining the class that serves as the base for many I/O streams classes
iosfwd	Declaring several I/O stream classes before they are necessarily defined
iostream	Declaring the I/O stream objects that manipulate the standard streams
istream	Defining the class that performs extractions
limits	Defining numerical values; only usable in C++
locale	Adapting to different cultural conventions; only usable in C++
new	Declaring several functions that allocate and free storage
ostream	Defining the class that performs insertions
sstream	Defining several I/O stream classes that manipulate string containers
stdexcept	Defining several classes useful for reporting exceptions; only usable in C++
streambuf	Defining classes that buffer I/O stream operations
string	Defining a class that implements a string container
strstream	Defining several I/O stream classes that manipulate in-memory character sequences
typeinfo	Defining type information support; only usable in C++

Table 38: C++ header files

The C++ standard template library (STL) header files

The following table lists the standard template library (STL) header files that can be used in C++ as well as in Extended Embedded C++:

Header file	Description
algorithm	Defines several common operations on sequences
bitset	Defining a container with fixed-sized sequences of bits; only usable in C++
deque	A deque sequence container
functional	Defines several function objects

Table 39: Standard template library header files

Header file	Description
hash_map	A map associative container, based on a hash algorithm
hash_set	A set associative container, based on a hash algorithm
iterator	Defines common iterators, and operations on iterators
list	A doubly-linked list sequence container
map	A map associative container
memory	Defines facilities for managing memory
numeric	Performs generalized numeric operations on sequences
queue	A queue sequence container
set	A set associative container
slist	A singly-linked list sequence container
stack	A stack sequence container
utility	Defines several utility components
valarray	Defining varying length array container; only usable in C++
vector	A vector sequence container

Table 39: Standard template library header files (Continued)

Using Standard C libraries in C++

The C++ library works in conjunction with some of the header files from the Standard C library, sometimes with small alterations. The header files come in two forms—new and traditional—for example, cassert and assert.h.

This table shows the new header files:

Header file	Usage
cassert	Enforcing assertions when functions execute
ccomplex	Computing common complex mathematical functions
cctype	Classifying characters
cerrno	Testing error codes reported by library functions
cfenv	Floating-point exception flags
cfloat	Testing floating-point type properties
cinttypes	Defining formatters for all types defined in stdint.h
ciso646	Using Amendment I—iso646.h standard header
climits	Testing integer type properties
clocale	Adapting to different cultural conventions

Table 40: New Standard C header files-DLIB

Header file	Usage
cmath	Computing common mathematical functions
csetjmp	Executing non-local goto statements
csignal	Controlling various exceptional conditions
cstdarg	Accessing a varying number of arguments
cstdbool	Adds support for the bool data type in C.
cstddef	Defining several useful types and macros
cstdint	Providing integer characteristics
cstdio	Performing input and output
cstdlib	Performing a variety of operations
cstring	Manipulating several kinds of strings
ctgmath	Type-generic mathematical functions
ctime	Converting between various time and date formats
cwchar	Support for wide characters
cwctype	Classifying wide characters

Table 40: New Standard C header files—DLIB (Continued)

LIBRARY FUNCTIONS AS INTRINSIC FUNCTIONS

Certain C library functions will under some circumstances be handled as intrinsic functions and will generate inline code instead of an ordinary function call, for example memcpy, memset, and strcat.

ADDED C FUNCTIONALITY

The IAR DLIB Library includes some added C functionality.

The following include files provide these features:

- fenv.h
- stdio.h
- stdlib.h
- string.h
- time.h

fenv.h

In fenv.h, trap handling support for floating-point numbers is defined with the functions fegettrapenable and fegettrapdisable.

stdio.h

These functions provide additional I/O functionality:

fdopen	Opens a file based on a low-level file descriptor.
fileno	Gets the low-level file descriptor from the file descriptor (FILE*).
gets	Corresponds to fgets on stdin.
getw	Gets a wchar_t character from stdin.
putw	Puts a wchar_t character to stdout.
ungetchar	Corresponds to ungetc on stdout.
write_array	Corresponds to fwrite on stdout.

string.h

These are the additional functions defined in string.h:

strdup	Duplicates a string on the heap.
strcasecmp	Compares strings case-insensitive.
strncasecmp	Compares strings case-insensitive and bounded.
strnlen	Bounded string length.

time.h

There are two interfaces for using time_t and the associated functions time, ctime, difftime, gmtime, localtime, and mktime:

- The 32-bit interface supports years from 1900 up to 2035 and uses a 32-bit integer for time_t. The type and function have names like __time32_t, __time32, etc. This variant is mainly available for backwards compatibility.
- The 64-bit interface supports years from -9999 up to 9999 and uses a signed long long for time_t. The type and function have names like __time64_t, __time64, etc.

The interfaces are defined in the system header file time.h.

An application can use either interface, and even mix them by explicitly using the 32or 64-bit variants. By default, the library and the header redirect time_t, time etc. to the 32-bit variants. However, to explicitly redirect them to their 64-bit variants, define _DLIB_TIME_USES_64 in front of the inclusion of time.h or ctime. See also, Time, page 127.

SYMBOLS USED INTERNALLY BY THE LIBRARY

The following symbols are used by the library, which means that they are visible in library source files, etc:

__assignment_by_bitwise_copy_allowed

This symbol determines properties for class objects.

___code, __data

These symbols are used as memory attributes internally by the compiler, and they might have to be used as arguments in certain templates.

__constrange()

Determines the allowed range for a parameter to an intrinsic function and that the parameter must be of type const.

__construction_by_bitwise_copy_allowed

This symbol determines properties for class objects.

__has_constructor, __has_destructor

These symbols determine properties for class objects and they function like the sizeof operator. The symbols are true when a class, base class, or member (recursively) has a user-defined constructor or destructor, respectively.

__memory_of

Determines the class memory. A class memory determines which memory a class object can reside in. This symbol can only occur in class definitions as a class memory.

Note: The symbols are reserved and should only be used by the library.

Use the compiler option --predef_macros to determine the value for any predefined symbols.

IAR DLIB Library

The linker configuration file

This chapter describes the purpose of the linker configuration file and describes its contents.

To read this chapter you must be familiar with the concept of sections, see *Modules and sections*, page 70.

Overview

To link and locate an application in memory according to your requirements, ILINK needs information about how to handle sections and how to place them into the available memory regions. In other words, ILINK needs a *configuration*, passed to it by means of the *linker configuration file*.

This file consists of a sequence of directives and typically, provides facilities for:

• Defining available addressable memories

giving the linker information about the maximum size of possible addresses and defining the available physical memory, as well as dealing with memories that can be addressed in different ways.

• Defining the regions of the available memories that are populated with ROM or RAM

giving the start and end address for each region.

Section groups

dealing with how to group sections into blocks and overlays depending on the section requirements.

• Defining how to handle initialization of the application

giving information about which sections that are to be initialized, and how that initialization should be made.

Memory allocation

defining where—in what memory region—each set of sections should be placed.

• Using symbols, expressions, and numbers

expressing addresses and sizes, etc, in the other configuration directives. The symbols can also be used in the application itself.

• Structural configuration

meaning that you can include or exclude directives depending on a condition, and to split the configuration file into several different files.

Comments can be written either as C comments (/*...*/) or as C++ comments (//...).

Defining memories and regions

ILINK needs information about the available memory spaces, or more specifically it needs information about:

• The maximum size of possible addressable memories

The define memory directive defines a *memory space* with a given size, which is the maximum possible amount of addressable memory, not necessarily physically available. See *Define memory directive*, page 400.

• Available physical memory

The define region directive defines a region in the available memories in which specific sections of application code and sections of application data can be placed. See *Define region directive*, page 401.

A region consists of one or several memory ranges. A range is a continuous sequence of bytes in a memory and several ranges can be expressed by using region expressions. See *Region expression*, page 403.

Define memory directive

Syntax	define memory [nam	<pre>define memory [name] with size = size_expr [,unit-size];</pre>		
	where unit-size is one of:			
		unitbitsize = <i>bitsize_expr</i> unitbytesize = <i>bytesize_expr</i>		
	and where <i>expr</i> is an expression, see <i>Expressions</i> , page 419.			
Parameters	size_expr	Specifies how many <i>units</i> the memory space contains; always counted from address zero.		
	bitsize_expr	Specifies how many bits each unit contains.		
	bytesize_expr	Specifies how many bytes each unit contains. Each byte contains 8 bits.		

Description	The define memory directive defines a <i>memory space</i> with a given size, which is the maximum possible amount of addressable memory, not necessarily physically available. This sets the limits for the possible addresses to be used in the linker configuration file. For many microcontrollers, one memory space is sufficient. However, some microcontrollers require two or more. For example, a Harvard architecture usually requires two different memory spaces, one for code and one for data. If only one memory is defined, the memory name is optional. If no <i>unit-size</i> is given, the unit contains 8 bits.
Example	/* Declare the memory space Mem of four Gigabytes */ define memory Mem with size = 4G;

Define region directive

Syntax	<pre>define region name = region-expr;</pre>		
	where <i>region-expr</i> is a region expression, see also <i>Regions</i> , page 401.		
Parameters	name	The name of the region.	
Description	The define region directive defines a region in which specific sections of code and sections of data can be placed. A region consists of one or several memory ranges, where each memory range consists of a continuous sequence of bytes in a specific memory. Several ranges can be combined by using region expressions. Note that those ranges do not need to be consecutive or even in the same memory.		
Example	0x10000 in memory Mer	yte code region ROM located at address n */ :[from 0x10000 size 0x10000];	

Regions

A *region* is s a set of non-overlapping memory ranges. A *region expression* is built up out of *region literals* and set operations (union, intersection, and difference) on regions.

Region literal

Syntax	[memory-name:][from expr { to expr size expr }		
	[repeat expr [displacement expr]]]		
	where <i>expr</i> is an expression, see <i>Expressions</i> , page 419.		
Parameters	memory-name	The name of the memory space in which the region literal will be located. If there is only one memory, the name is optional.	
	from <i>expr</i>	expr is the start address of the memory range (inclusive).	
	to <i>expr</i>	expr is the end address of the memory range (inclusive).	
	size <i>expr</i>	<i>expr</i> is the size of the memory range.	
	repeat <i>expr</i>	<i>expr</i> defines several ranges in the same memory for the region literal.	
	displacement <i>expr</i>	<i>expr</i> is the displacement from the previous range start in the repeat sequence. Default displacement is the same value as the range size.	
Description	A region literal consists of one memory range. When you define a range, the memory it resides in, a start address, and a size must be specified. The range size can be stated explicitly by specifying a size, or implicitly by specifying the final address of the range. The final address is included in the range and a zero-sized range will only contain an address. A range can span over the address zero and such a range can even be expressed by unsigned values, because it is known where the memory wraps. The repeat parameter will create a region literal that contains several ranges, one for each repeat. This is useful for <i>banked</i> or <i>far</i> regions.		

```
Example
                      /* The 5-byte size range spans over the address zero */
                      Mem: [from -2 to 2]
                      /* The 512-byte size range spans over zero, in a 64-Kbyte memory
                      */
                      Mem:[from 0xFF00 to 0xFF]
                      /* Defining several ranges in the same memory, a repeating
                         literal */
                      Mem:[from 0 size 0x100 repeat 3 displacement 0x1000]
                      /* Resulting in a region containing:
                         Mem:[from 0 size 0x100]
                         Mem:[from 0x1000 size 0x100]
                         Mem: [from 0x2000 size 0x100]
                      */
See also
                      Define region directive, page 401, and Region expression, page 403.
```

Region expression

Syntax	region-operand region-expr region-operand region-expr - region-operand region-expr & region-operand		
	where region-operand is one of:		
	<pre>(region-expr) region-name region-literal empty-region where region-name is a region, see Define region directive, page 401 where region-literal is a region literal, see Region literal, page 402</pre>		
	and where <i>empty-region</i> is an empty region, see <i>Empty region</i> , page 404.		
Description	Normally, a region consists of one memory range, which means a <i>region literal</i> is sufficient to express it. When a region contains several ranges, possibly in different memories, it is instead necessary to use a <i>region expression</i> to express it. Region expressions are actually set expressions on sets of memory ranges.		

	To create region expressions, three operators are available: union ($ $), intersection ($\&$), and difference (–). These operators work as in <i>set theory</i> . For example, if you have the sets A and B, then the result of the operators would be:	
	• A B: all elements in either set A or set B	
	 A & B: all elements in both set A and B A - B: all elements in set A but not in B. 	
	• A - B. an elements in set A but not in B.	
Example	/* Resulting in a range starting at 1000 and ending at 2FFF, in memory Mem */	
	Mem:[from 0x1000 to 0x1FFF] Mem:[from 0x1500 to 0x2FFF]	
	/* Resulting in a range starting at 1500 and ending at 1FFF, in memory Mem */	
	Mem:[from 0x1000 to 0x1FFF] & Mem:[from 0x1500 to 0x2FFF]	
	/* Resulting in a range starting at 1000 and ending at 14FF, in memory Mem */	
	Mem:[from 0x1000 to 0x1FFF] - Mem:[from 0x1500 to 0x2FFF]	
	<pre>/* Resulting in two ranges. The first starting at 1000 and ending at 1FFF, the second starting at 2501 and ending at 2FFF. Both located in memory Mem */</pre>	
	Mem:[from 0x1000 to 0x2FFF] - Mem:[from 0x2000 to 0x24FF]	

Empty region

Syntax	[]	
Description	The empty region does not contain any memory ranges. If the empty region is used in a placement directive that actually is used for placing one or more sections, ILINK will issue an error.	
Example	<pre>define region Code = Mem:[from 0 size 0x10000]; if (Banked) { define region Bank = Mem:[from 0x8000 size 0x1000]; } else { define region Bank = []; } define region NonBanked = Code - Bank; /* Depending on the Banked symbol, the NonBanked region is either one range with 0x10000 bytes, or two ranges with 0x8000 and 0x7000 bytes, respectively. */</pre>	

See also

Region expression, page 403.

Section handling

Section handling describes how ILINK should handle the sections of the execution image, which means:

Placing sections in regions

The place at and place into directives place sets of sections with similar attributes into previously defined regions. See *Place at directive*, page 412 and *Place in directive*, page 413.

• Making sets of sections with special requirements

The block directive makes it possible to create empty sections with specific sizes and alignments, sequentially sorted sections of different types, etc.

The overlay directive makes it possible to create an area of memory that can contain several overlay images. See *Define block directive*, page 405, and *Define overlay directive*, page 407.

• Initializing the application

The directives initialize and do not initialize control how the application should be started. With these directives, the application can initialize global symbols at startup, and copy pieces of code. The initializers can be stored in several ways, for example they can be compressed. See *Initialize directive*, page 408 and *Do not initialize directive*, page 410.

Keeping removed sections

The keep directive retains sections even though they are not referred to by the rest of the application, which means it is equivalent to the *root* concept in the assembler and compiler. See *Keep directive*, page 411.

Define block directive

```
Syntax
```

```
define [movable] block name
  [ with param, param... ]
{
    extended-selectors
}
[except
  {
    section_selectors
  }];
```

where *param* can be one of:

```
size = expr
maximum size = expr
alignment = expr
fixed order
static base [name]
static base
```

and where the rest of the directive selects sections to include in the block, see *Section* selection, page 413.

Parameters		
	name	The name of the block to be defined.
	size	Customizes the size of the block. By default, the size of a block is the sum of its parts dependent of its contents.
	maximum size	Specifies an upper limit for the size of the block. An error is generated if the sections in the block do not fit.
	alignment	Specifies a minimum alignment for the block. If any section in the block has a higher alignment than the minimum alignment, the block will have that alignment.
	fixed order	Places sections in fixed order; if not specified, the order of the sections will be arbitrary.
	static base <i>name</i>	Specifies that the static base with the specified name will be at the start of the block when the application executes.
	static base	Specifies that the static base will be at the start of the block when the application executes.
Description	criptionThe block directive defines a named set of sections. By defining a block you can create empty blocks of bytes that can be used, for example as stacks or heaps. Another use for the directive is to group certain types of sections, consecutive or non-consecutive. A third example of use for the directive is to group sections into one memory area to access the start and end of that area from the application.movable blocks are for use with read-only and read-write position independence. Making blocks movable enables the linker to validate the application's use of addresses. Movable blocks are located in exactly the same way as other blocks, but the linker will check that the appropriate relocations are used when referring to symbols in movable blocks.	
Example	/* Create a 0x1000-byte block for the heap */ define block HEAP with size = 0x1000, alignment = 8 { };	

See also	Interaction between the tools and your application, page 186. See Define overlay
	directive, page 407 for an Accessing example.

Define overlay directive

Syntax

define overlay name [with param, param...]
{
 extended-selectors;
}
[except
{
 section_selectors
 }];

For information about extended selectors and except clauses, see *Section selection*, page 413.

Parameters		
	name	The name of the overlay.
	size	Customizes the size of the overlay. By default, the size of a overlay is the sum of its parts dependent of its contents.
	maximum size	Specifies an upper limit for the size of the overlay. An error is generated if the sections in the overlay do not fit.
	alignment	Specifies a minimum alignment for the overlay. If any section in the overlay has a higher alignment than the minimum alignment, the overlay will have that alignment.
	fixed order	Places sections in fixed order; if not specified, the order of the sections will be arbitrary.
Description	The overlay directive defines a named set of sections. In contrast to the block directive, the overlay directive can define the same name several times. Each definition will then be grouped in memory at the same place as all other definitions of the same name. This creates an <i>overlaid</i> memory area, which can be useful for an application that has several independent sub-applications. Place each sub-application image in ROM and reserve a RAM overlay area that can hold all sub-applications. To execute a sub-application, first copy it from ROM to the RAM overlay. Note that ILINK does not help you with managing interdependent overlay definitions, apart from generating a diagnostic message for any reference from one overlay to another overlay.	

The size of an overlay will be the same size as the largest definition of that overlay name and the alignment requirements will be the same as for the definition with the highest alignment requirements.
Note: Sections that were overlaid must be split into a RAM and a ROM part and you must take care of all the copying needed.
Manual initialization, page 91.
<pre>initialize { by copy manually } [with packing = algorithm] { section-selectors } [except { section_selectors</pre>

```
}];
```

where the rest of the directive selects sections to include in the block. See *Section* selection, page 413.

Parameters	ру сору	Splits the section into sections for initializers and initialized data, and handles the initialization at application startup automatically.
	manually	Splits the section into sections for initializers and initialized data. The initialization at application startup is not handled automatically.

	algorithm	Specifies how to handle the initializers. Choose between:
		none - Disables compression of the selected section contents. This is the default method for initialize manually.
		${\tt zeros}$ – Compresses consecutive bytes with the value zero.
		packbits - Compresses with the PackBits algorithm. This method generates good results for data with many consecutive bytes of the same value.
		1z77 - Compresses with the Lempel-Ziv-77 algorithm.This method handles a larger variety of inputs well, but has a slightly larger decompressor.
		bwt - Compresses with the Burrows-Wheeler algorithm. This method improves the packbits method by transforming blocks of data before they are compressed.
		l_{ZW} - Compresses with the Lempel-Ziv-Welch algorithm. This method uses a dictionary to store byte patterns in the data.
		auto - Similar to smallest, but ILINK chooses between none, packbits, and 1z77. This is the default method for initialize by copy.
		<pre>smallest - ILINK estimates the resulting size using each packing method (except for auto), and then chooses the packing method that produces the smallest estimated size. Note that the size of the decompressor is also included.</pre>
Description	initializers and another s initialization at startup s	ive splits the initialization section into one section holding the section holding the initialized data. You can choose whether the hould be handled automatically (initialize by copy) or dle it yourself (initialize manually).
	When you use the packing method auto (default for initialize by copy) or smallest, ILINK will automatically choose an appropriate packing algorithm for the initializers. To override this, specify a different packing method. Thelog initialization option shows how ILINK decided which packing algorithm to use.	
	The decompressors for h	mpressed, a decompressor is automatically added to the image. owt and lzw use significantly more execution time and RAM for the other methods. Approximately 9 Kbytes of stack space 5 Kbytes for lzw.

	When initializers are compressed, the exact size of the compressed initializers is unknown until the exact content of the uncompressed data is known. If this data contains other addresses, and some of these addresses are dependent on the size of the compressed initializers, the linker fails with error Lp017. To avoid this, place compressed initializers last, or in a memory region together with sections whose addresses do not need to be known.
	Unless initialize manually is used, ILINK will arrange for initialization to occur during system startup by including an initialization table. Startup code calls an initialization routine that reads this table and performs the necessary initializations.
	Zero-initialized sections are not affected by the initialize directive.
	The initialize directive is normally used for initialized variables, but can be used for copying any sections, for example copying executable code from slow ROM to fast RAM, or for overlays. For another example, see <i>Define overlay directive</i> , page 407.
	Sections that are needed for initialization are not affected by the initialize by copy directive. This includes thelow_level_init function and anything it references.
	Anything reachable from the program entry label is considered <i>needed for initialization</i> unless reached via a section fragment with a label starting withiar_init\$\$done. Thelog sections option, in addition to logging the marking of section fragments to be included in the application, also logs the process of determining which sections are needed for initialization.
Example	<pre>/* Copy all read-write sections automatically from ROM to RAM at program start */ initialize by copy { rw }; place in RAM { rw }; place in ROM { ro };</pre>
See also	Initialization at system startup, page 76, and Do not initialize directive, page 410.

Do not initialize directive

```
Syntax
```

```
do not initialize
{
   section-selectors
}
[except
   {
   section-selectors
}];
```

	For information about extended selectors and except clauses, see <i>Section selection</i> , page 413.
Description	The do not initialize directive specifies the sections that should not be initialized by the system startup code. The directive can only be used on zeroinit sections.
	The compiler keywordno_init places variables into sections that must be handled by a do not initialize directive.
Example	<pre>/* Do not initialize read-write sections whose name ends with _noinit at program start */ do not initialize { rw section .*_noinit }; place in RAM { rw section .*_noinit };</pre>
See also	Initialization at system startup, page 76, and Initialize directive, page 408.

Keep directive

Syntax	keep
	{
	section-selectors
	}
	[except
	{
	section-selectors
	} 1 ;
	For information about extended selectors and except clauses, see <i>Section selection</i> , page 413.
Description	The keep directive specifies that all selected sections should be kept in the executable image, even if there are no references to the sections.
Example	<pre>keep { section .keep* } except {section .keep};</pre>

Place at directive

Syntax	<pre>["name":] place at { address [memory:] expr start of region_expr end of region_expr }</pre>	
	{ extended-selectors }	
	<pre>[except { section-selectors }</pre>	
	} 1 ;	
Parameters	memory: expr	A specific address in a specific memory. The address must be available in the supplied memory defined by the define memory directive. The memory specifier is optional if there is only one memory.
	start of <i>region_expr</i>	A region expression that results in a single-internal region. The start of the interval is used.
	end of <i>region_expr</i>	A region expression that results in a single-internal region. The end of the interval is used.
Description	The place at directive places sections and blocks either at a specific address or, at the beginning or the end of a region. The same address cannot be used for two different place at directives. It is also not possible to use an empty region in a place at directive. If placed in a region, the sections and blocks will be placed before ny other sections or blocks placed in the same region with a place in directive.	
	The name, if specified, is used in the ma	ap file and in some log messages.
Example	code_region */	<pre>.startup at the beginning of the [{ readonly section .startup };</pre>
See also	Place in directive, page 413.	

Place in directive

Syntax	<pre>["name":] place in region-expr { extended-selectors } [except{ section-selectors }]; where region-expr is a region expression, see also Regions, page 401. and where the rest of the directive selects sections to include in the block. See Section selection, page 413.</pre>
Description	The place in directive places sections and blocks in a specific region. The sections and blocks will be placed in the region in an arbitrary order.To specify a specific order, use the block directive. The region can have several ranges.The name, if specified, is used in the map file and in some log messages.
Example	<pre>/* Place the read-only sections in the code_region */ "ROM": place in ROM { readonly };</pre>
See also	Place at directive, page 412.

Section selection

The purpose of *section selection* is to specify—by means of *section selectors* and *except clauses*—the sections that an ILINK directive should be applied to. All sections that match one or more of the section selectors will be selected, and none of the sections selectors in the except clause, if any. Each section selector can match sections on section attributes, section name, and object or library name.

Some directives provide functionality that requires more detailed selection capabilities, for example directives that can be applied on both sections and blocks. In this case, the *extended-selectors* are used.

Section-selectors

Syntax

{ [section-selector][, section-selector...] }

	where section-selector is:		
	[section-attribute][section-type][section sectionname] [object {module filename }]		
	where section-attri	bute is:	
	[ro [code data	a] rw [code data] zi]	
	and where ro, rw, and zi also can be readonly, readwrite, and zeroinit, respectively.		
	And section-type is:		
	[preinit_array	init_array]	
Parameters			
	ro Or readonly	Read-only sections.	
	rw Of readwrite	Read/write sections.	
	zi or zeroinit	The name of the block.	
	code	Sections that contain code.	
	data	Sections that contain data.	
	preinit_array	Sections of the ELF section type $SHT_PREINIT_ARRAY$.	
	init_array	Sections of the ELF section type SHT_INIT_ARRAY .	
	sectionname	The section name. Two wildcards are allowed:	
		? matches any single character	
		* matches zero or more characters.	
	module	A name in the form <i>objectname(libraryname)</i> . Sections from object modules where both the object name and the library name match their respective patterns are selected. An empty library name pattern selects only sections from object files.	
	filename	The name of an object file, a library, or an object in a library. Two wildcards are allowed:	
		? matches any single character	
		* matches zero or more characters.	

Description A section selector selects all sections that match the section attribute, section type, section name, and the name of the *object*, where *object* is an object file, a library, or an object in a library. Up to three of the four conditions can be omitted. If the section attribute is omitted, any section will be selected, without restrictions on the section attribute. If the section type is omitted, sections of any type will be selected.

If the section name part or the object name part is omitted, sections will be selected without restrictions on the section name or object name, respectively.

It is also possible to use only $\{ \ \}$ without any section selectors, which can be useful when defining blocks.

Note that a section selector with narrower scope has higher priority than a more generic section selector.

If more than one section selector matches for the same purpose, one of them must be more specific. A section selector is more specific than another one if:

- It specifies a section type and the other one does not
- It specifies a section name or object name with no wildcards and the other one does not
- There could be sections that match the other selector that also match this one, and the reverse is not true.

Selector I	Selector 2	More specific
section "foo*"	section "f*"	Selector 1
section "*x"	section "f*"	Neither
no code section "f*"	ro section "f*"	Selector 1
init_array	ro section "xx"	Selector 1
section ".intvec"	ro section ".int"	Selector 1
section ".intvec"	object "foo.o"	Neither

Table 41: Examples of section selector specifications

{ rw }

Example

```
/* Selects all read-write sections */
```

{ section .mydata* } /* Selects only .mydata* sections */
/* Selects .mydata* sections available in the object special.o */
{ section .mydata* object special.o }

Assuming a section in an object named foo.o in a library named lib.a, any of these selectors will select that section:

```
object foo.o{lib.a}
object f*{lib*}
object foo.o
object lib.a
```

See also Initialize directive, page 408, Do not initialize directive, page 410, and Keep directive, page 411.

Extended-selectors

Syntax	<pre>{ [extended-selector][, extended-selector] } where extended-selector is:</pre>		
	[first last]{	section-selector block name [inline-bloc overlay name }	ck-def]
	where inline-block-	def is:	
	[block-params]	extended-selectors	
Parameters			
	first	Places the selected name first in t	the region, block, or overlay.
	last	Places the selected name last in t	he region, block, or overlay.
	block	The name of the block.	
	overlay	The name of the overlay.	
Description	functionality for placing an overlay. It is also pos	section-selector does, extend g blocks or overlays first or last in a ssible to create an <i>inline</i> definition e control over section placement.	a set of sections, a block, or
Example		<pre>{ section .first }; /* De th t { first block First }; /*</pre>	ne section .first */
	or, equivalently using a	n inline definition of the block Fir	st:
	define block Table	e { first block First { sec	ction .first }};

See also Define block directive, page 405, Define overlay directive, page 407, and Place at directive, page 412.

Using symbols, expressions, and numbers

In the linker configuration file, you can also:

• Define and export symbols

The define symbol directive defines a symbol with a specified value that can be used in expressions in the configuration file. The symbol can also be exported to be used by the application or the debugger. See *Define symbol directive*, page 418, and *Export directive*, page 418.

• Use expressions and numbers

In the linker configuration file, expressions and numbers are used for specifying addresses, sizes, etc. See *Expressions*, page 419.

Check that directive

Syntax	check that <i>expression</i> ;	
Parameters	expression	A boolean expression.
Description	You can use the check that directive to compare the results of stack usage analysis against the sizes of blocks and regions. If the expression evaluates to zero, an error is emitted. Three extra operators are available for use only in check that expressions:	
	<pre>maxstack(category)</pre>	The stack depth of the deepest call chain for any call graph root function in the category.
	<pre>totalstack(category)</pre>	The sum of the stack depths of the deepest call chains for each call graph root function in the category.
	<pre>size(block)</pre>	The size of the block.
Example	<pre>check that maxstack("Program entry")</pre>	
See also	Stack usage analysis, page 78.	

Define symbol directive

Syntax	<pre>define [exported] symbol name = expr;</pre>	
Parameters	exported	Exports the symbol to be usable by the executable image.
	name	The name of the symbol.
	expr	The symbol value.
DescriptionThe define symbol directive defines a symbol with a specified v then be used in expressions in the configuration file. The symbols work exactly like the symbols defined with the optionconfig configuration file.The define exported symbol variant of this directive is a short directive define symbol in combination with the export symbol command line this would require both aconfig_def option at define_symbol option to achieve the same effect.		the configuration file. The symbols defined in this way
		ombination with the export symbol directive. On the ire both aconfig_def option and a
	Note:	
	• A symbol cannot be redefined	ned
	• Symbols that are either pref (double underscore) are res	fixed by x , where x is a capital letter, or that contain erved for toolset vendors.
Example	/* Define the symbol my define symbol my_symbol	r_symbol with the value 4 $*/$
See also	Export directive, page 418 and Interaction between ILINK and the application, page 94.	
Export directive		
Syntax	export symbol <i>name;</i>	
Parameters	name	The name of the symbol.

Description The export directive defines a symbol to be exported, so that it can be used both from the executable image and from a global label. The application, or the debugger, can then refer to it for setup purposes etc.

Example /* Define the symbol my_symbol to be exported */
export symbol my_symbol;

Expressions

Syntax

An expression is built up of the following constituents:

expression binop expression unop expression expression ? expression : expression (expression) number symbol func-operator

where *binop* is one of these binary operators:

+, -, *, /, %, <<, >>, <, >, ==, !=, &, ^, |, &&, | |

where unop is one of this unary operators:

+, -, !, ~

where number is a number, see Numbers, page 420

where *symbol* is a defined symbol, see *Define symbol directive*, page 418 and *--config def*, page 274

and where *func-operator* is one of these function-like operators:

<pre>minimum(expr,expr)</pre>	Returns the smallest of the two parameters.
<pre>maximum(expr, expr)</pre>	Returns the largest of the two parameters.
isempty(r)	Returns True if the region is empty, otherwise False.
<pre>isdefinedsymbol(expr-symbol)</pre>	Returns True if the expression symbol is defined, otherwise False.
<pre>start(r)</pre>	Returns the lowest address in the region.
end(r)	Returns the highest address in the region.
size(r)	Returns the size of the complete region.

where expr is an expression, and r is a region expression, see *Region expression*, page 403.

Description	In the linker configuration file, an expression is a 65-bit value with the range -2^64 to 2^64. The expression syntax closely follows C syntax with some minor exceptions. There are no assignments, casts, pre- or post-operations, and no address operations (*, $\&$, [], ->, and .). Some operations that extract a value from a region expression, etc, use a syntax resembling that of a function call. A boolean expression returns 0 (False) or 1 (True).	
Numbers		
Syntax	nr [nr-suffix]	
	where nr is either a decimal number or a hexadecimal number $(0x or 0X)$.	
	and where <i>nr-suffix</i> is one of:	
	K /* Kilo = (1 << 10) 1024 */	
	M /* Mega = (1 << 20) 1048576 */	
	G /* Giga = (1 << 30) 1073741824 */	
	T /* Tera = (1 << 40) 1099511627776 */	
	P /* Peta = (1 << 50) 1125899906842624 */	
Description	A number can be expressed either by normal C means or by suffixing it with a set of useful suffixes, which provides a compact way of specifying numbers.	
Example	1024 is the same as 0×400 , which is the same as 1κ .	

Structural configuration

The structural directives provide means for creating structure within the configuration, such as:

• Conditional inclusion

An if directive includes or excludes other directives depending on a condition, which makes it possible to have directives for several different memory configurations in the same file. See If directive, page 421.

• Dividing the linker configuration file into several different files

The include directive makes it possible to divide the configuration file into several logically distinct files. See Include directive, page 421.

If directive

Syntax	<pre>if (expr) { directives [} else if (expr) { directives] [} else { directives] } where expr is an expression, see Expressions, page 419.</pre>
Parameters	directives Any ILINK directive.
Description	An if directive includes or excludes other directives depending on a condition, which makes it possible to have directives for several different memory configurations, for example both a banked and non-banked memory configuration, in the same file.
	The directives inside an if part, else if part, or an else part are syntax checked and processed regardless of whether the conditional expression was true or false, but only the directives in the part where the conditional expression was true, or the else part if none of the conditions were true, will have any effect outside the if directive. The if directives can be nested.
Example	See Empty region, page 404.

Include directive

Syntax	include <i>filename</i> ;	
Parameters	filename	A string literal where both / and \ can be used as the directory delimiter.
Description	The include directive makes it possible to divide the configuration file into several logically distinct parts, each in a separate file. For instance, there might be parts that you need to change often and parts that you seldom edit.	

Structural configuration

Section reference

The compiler places code and data into sections. Based on a configuration specified in the linker configuration file, ILINK places sections in memory.

This chapter lists all predefined ELF sections and blocks that are available for the IAR build tools for ARM, and gives detailed reference information about each section.

For more information about sections, see the chapter *Modules and sections*, page 70.

Summary of sections

This table lists the ELF sections and blocks that are used by the IAR build tools:

Section	Description
.bss	Holds zero-initialized static and global variables.
CSTACK	Holds the stack used by programs.
.cstart	Holds the startup code.
.data	Holds static and global initialized variables.
.data_init	Holds initial values for .data sections when the linker directive initialize by copy is used.
DLIB_PERTHREAD	Holds variables that contain static states for DLIB modules.
.exc.text	Holds exception-related code.
HEAP	Holds the heap used for dynamically allocated data.
.iar.dynexit	Holds the atexit table.
.init_array	Holds a table of dynamic initialization functions.
.intvec	Holds the reset vector table
IRQ_STACK	Holds the stack for interrupt requests, IRQ, and exceptions.
.noinit	Holdsno_init static and global variables.
.preinit_array	Holds a table of dynamic initialization functions.
.prepreinit_array	Holds a table of dynamic initialization functions.
.rodata	Holds constant data.

Table 42: Section summary

.text Holds the program code. .textrw Holdsramfunc declared program code. .textrw_init Holds initializers for the .textrw declared section.	Section	Description
	.text	Holds the program code.
.textrw_init Holds initializers for the .textrw declared section.	.textrw	Holdsramfunc declared program code.
	.textrw_init	Holds initializers for the $\mbox{.textrw}$ declared section.

Table 42: Section summary (Continued)

In addition to the ELF sections used for your application, the tools use a number of other ELF sections for a variety of purposes:

- Sections starting with .debug generally contain debug information in the DWARF format
- Sections starting with .iar.debug contain supplemental debug information in an IAR format
- The section .comment contains the tools and command lines used for building the file
- Sections starting with .rel or .rela contain ELF relocation information
- The section . symtab contains the symbol table for a file
- The section .strtab contains the names of the symbol in the symbol table
- The section . shstrtab contains the names of the sections.

Descriptions of sections and blocks

This section gives reference information about each section, where the:

- *Description* describes what type of content the section is holding and, where required, how the section is treated by the linker
- Memory placement describes memory placement restrictions.

For information about how to allocate sections in memory by modifying the linker configuration file, see *Placing code and data—the linker configuration file*, page 73.

.bss

Description	Holds zero-initialized static and global variables.
Memory placement	This section can be placed anywhere in memory.

CSTACK

Description

Block that holds the internal data stack.

	Memory placement	This block can be placed anywhere in memory.
	See also	Setting up the stack, page 89.
.csta		
.CSLa	irt	
	Description	Holds the startup code.
	Memory placement	This section can be placed anywhere in memory.
.dat	2	
.uat	a	
	Description	Holds static and global initialized variables. In object files, this includes the initial values. When the linker directive initialize by copy is used, a corresponding .data_init section is created for each .data section, holding the possibly compressed initial values.
	Memory placement	This section can be placed anywhere in memory.
.dat	a_init	
	Description	Holds the possibly compressed initial values for .data sections. This section is created by the linker if the initialize by copy linker directive is used.
	Memory placement	This section can be placed anywhere in memory.
П	LIB_PERTHREAD	
	Description	Holds thread-local static and global initialized variables used by the main thread.
		This section is placed automatically. If you change the placement, you must not change its initialization. The initialization of this section must be controlled with the

Memory placement This section can be placed anywhere in memory.

See also *Managing a multithreaded environment*, page 130.

initialize by copy directive.

.exc.text

Description	Holds code that is only executed when your application handles an exception.
Memory placement	In the same memory as .text.
See also	Exception handling, page 174.

HEAP

Description	Holds the heap used for dynamically allocated data in memory, in other words data allocated by malloc and free, and in C++, new and delete.
Memory placement	This section can be placed anywhere in memory.
See also	Setting up the heap, page 90.

.iar.dynexit

Description	Holds the table of calls to be made at exit.
Memory placement	This section can be placed anywhere in memory.
See also	Setting up the atexit limit, page 90.

.init_array

Description	Holds pointers to routines to call for initializing one or more C++ objects with static storage duration.
Memory placement	This section can be placed anywhere in memory.

.intvec

Description	Holds the reset vector table and exception vectors which contain branch instructions to cstartup, interrupt service routines etc.
Memory placement	Must be placed at address range 0x00 to 0x3F.

IRQ_STACK

Description	Holds the stack which is used when servicing IRQ exceptions. Other stacks may be added as needed for servicing other exception types: FIQ, SVC, ABT, and UND. The cstartup.s file must be modified to initialize the exception stack pointers used.	
	Note: This section is not used when compiling for Cortex-M.	
Memory placement	This section can be placed anywhere in memory.	
See also	Exception stack, page 184.	

.noinit

Description	Holds static and globalno_init variables.
Memory placement	This section can be placed anywhere in memory.

.preinit_array

Description	Like .init_array, but is used by the library to make some C++ initializations happen before the others.
Memory placement	This section can be placed anywhere in memory.
See also	.init_array, page 426.

.prepreinit_array

Description	Like .init_array, but is used when C static initialization is rewritten as dynamic initialization. Performed before all C++ dynamic initialization.
Memory placement	This section can be placed anywhere in memory.
See also	. <i>init_array</i> , page 426.

.rodata

Description	Holds constant data. This can include constant variables, string and aggregate literals,
	etc.

	Memory placement	This section can be placed anywhere in memory.
.tex	t	
	Description	Holds program code, except the code for system initialization.
	Memory placement	This section can be placed anywhere in memory.
.tex	trw	
	Description	Holdsramfunc declared program code.
	Memory placement	This section can be placed anywhere in memory.
	See also	<i>ramfunc</i> , page 321.

.textrw_init

Description	Holds initializers for the .textrw declared sections.
Memory placement	This section can be placed anywhere in memory.
See also	<i>ramfunc</i> , page 321.

Stack usage control files

This chapter describes the purpose and the syntax of stack usage control files.

Before you read this chapter, see Stack usage analysis, page 78.

Overview

A stack usage control file consists of a sequence of directives that control stack usage analysis. You can use C ("/*...*/") and C++ ("//...") comments in these files.

The default filename extension for stack usage control files is suc.

C++ NAMES

When you specify the name of a C++ function in a stack usage control file, you must use the name exactly as used by the linker. Both the number and names of parameters, as well as the names of types must match. However, most non-significant white-space differences are accepted. In particular, you must enclose the name in quote marks because all C++ function names include non-identifier characters.

You can also use wildcards in function names. "#*" matches any sequence of characters, and "#?" matches a single character. This makes it possible to write function names that will match any instantiation of a template function.

Examples:

Stack usage control directives

This section gives detailed reference information about each stack usage control directive.

function directive

Syntax	[override] function [category] function-spec : stack-size [, call-info];	
Parameters	category function-spec call-info stack-size	See <i>category</i> , page 433 See <i>function-spec</i> , page 433 See <i>call-info</i> , page 434 See <i>stack-size</i> , page 434
Description	Specifies what the maximum stack usage is in a function and which other functions that are called from that function. Normally, an error is issued if there already is stack usage information for the function, but if you start with override, the error will be suppressed and the information supplied in the directive will be used instead of the previous information.	
Example	<pre>function MyFunc1: 32, calls MyFunc2, calls MyFunc3, MyFunc- function [interrupt] nm.</pre>	

exclude directive

Syntax	exclude function-spec [,	function-spec];
Parameters	function-spec	See function-spec, page 433
Description	Excludes the specified functions, and call trees originating with them, from stack usage calculations.	
Example	exclude fun1, fun2;	

possible calls directive

Syntax	possible calls calling-func : called-func [, called-func \ldots];	
Parameters	calling-func called-func	See <i>function-spec</i> , page 433 See <i>function-spec</i> , page 433
Description	Specifies an exhaustive list of possible destinations for all indirect calls in one function. Use this for functions which are known to perform indirect calls and where you know exactly which functions that might be called in this particular application. Consider using the #pragma calls directive if the information about which functions that might be called is available when compiling.	
Example	possible calls afun: bfun, cfun;	
See also	calls, page 328.	

call graph root directive

Syntax	<pre>call graph root [category] : function-spec [, function-spec];</pre>	
Parameters	category	See <i>category</i> , page 433
	function-spec	See function-spec, page 433
Description	Specifies that the listed functions are call graph roots. You can optionally specify a call graph root category. Call graph roots are listed under their category in the <i>Stack Usage</i> chapter in the linker map file. The linker will normally issue a warning for functions needed in the application that are not call graph roots and which do not appear to be called.	
Example	call graph root [task]: fun1, fun2;	
See also	<i>call_graph_root</i> , page 328.	

max recursion depth directive

Syntax	max recursion depth function-spec : size;	
Parameters	function-spec size	See <i>function-spec</i> , page 433 See <i>size</i> , page 435
Description	Specifies the maximum number of iterations through any of the cycles in the recursion nest of which the function is a member.A recursion nest is a set of cycles in the call graph where each cycle shares at least one node with another cycle in the nest.Stack usage analysis will base its result on the max recursion depth multiplied by the stack usage of the deepest cycle in the nest. If the nest is not entered on a point along one of the deepest cycles, no stack usage result will be calculated for such calls.	
Example	max recursion depth fun1: 10;	

no calls from directive

Syntax	no calls from module-spec to function-spec [, function-spec];	
Parameters	function-spec module-spec	See <i>function-spec</i> , page 433 See <i>module-spec</i> , page 433
Description	When you provide stack usage information for some functions in a module without stack usage information, the linker warns about functions that are referenced from the module but not listed as called. This is primarily to help avoid problems with C runtime routines, calls to which are generated by the compiler, beyond user control.	
	If there actually is no call to some of these functions, use the no calls from directive to selectively suppress the warning for the specified functions. You can also disable the warning entirely (diag_suppress or Project>Options>Linker>Diagnostics>Suppress these diagnostics).	
Example	no calls from [file.o]	to fun1, fun2;

Syntactic components

The stack usage control directives use some syntactical components. These are described below.

category

Syntax	[name]
Description	A call graph root category. You can use any name you like. Categories are not case-sensitive.
Example	category examples:
	[interrupt] [task]

function-spec

Syntax	[?] name [module-spec]
Description	Specifies the name of a symbol, and for module-local symbols, the name of the module it is defined in. Normally, if the function-spec does not match a symbol in the program, a warning is emitted. Prefixing with ? suppresses this warning.
Example	<i>function-spec</i> examples:
	xFun MyFun [file.o] ?"fun1(int)"
module-spec	
Syntax	[name [(<i>name</i>)]]
Description	Specifies the name of a module, and optionally, in parentheses, the name of the library it belongs to. To distinguish between modules with the same name, you can specify:
	• The complete path of the file ("D:\C1\test\file.o")
	• As many path elements as are needed at the end of the path ("test\file.o")
	• Some path elements at the start of the path, followed by "", followed by some path elements at the end ("D:\\file.o").

	Note that when using multi-file compilation (mfc), multiple files are compiled into a single module, named after the first file.
Example	module-spec examples:
	<pre>[file.o] [file.o(lib.a)] ["D:\Cl\test\file.o"]</pre>
name	
Description	A name can be either an identifier or a quoted string.
	The first character of an identifier must be either a letter or one of the characters "_", "\$", or ".". The rest of the characters can also be digits.
	A quoted string starts and ends with " and can contain any character. Two consecutive " characters can be used inside a quoted string to represent a single ".
Example	name examples:
	MyFun
	file.o "file-1.o"

call-info

Syntax	<pre>calls function-spec [, function-spec][: stack-size]</pre>
Description	Specifies one or more called functions, and optionally, the stack size at the calls.
Example	call-info examples:
	calls MyFunc1 : stack 16 calls MyFunc2, MyFunc3, MyFunc4

stack-size

Syntax	[stack] <i>size</i>
Description	Specifies the size of a stack frame.

Example	stack-size examples:
	24
	stack 28
size	

Description	A decimal integer, or $0x$ followed by a hexadecimal integer. Either alternative can optionally be followed by a suffix indicating a power of two (K=210, M=220, G=230, T=240, P=250).
Example	size examples:
	24
	0x18
	2048
	2K

Syntactic components

IAR utilities

This chapter describes the IAR command line utilities that are available:

- The IAR Archive Tool—iarchive—creates and manipulates a library (an archive) of several ELF object files
- The IAR ELF Tool—ielftool—performs various transformations on an ELF executable image (such as fill, checksum, format conversions, etc)
- The IAR ELF Dumper for arm—ielfdumparm—creates a text representation of the contents of an ELF relocatable or executable image
- The IAR ELF Object Tool—iobjmanip—is used for performing low-level manipulation of ELF object files
- The IAR Absolute Symbol Exporter—isymexport—exports absolute symbols from a ROM image file, so that they can be used when you link an add-on application.

Descriptions of options gives detailed reference information about each command line option available for the different utilities.

The IAR Archive Tool—iarchive

The IAR Archive Tool, iarchive, can create a library (an archive) file from several ELF object files. You can also use iarchaive to manipulate ELF libraries.

A library file contains several relocatable ELF object modules, each of which can be independently used by a linker. In contrast with object modules specified directly to the linker, each module in a library is only included if it is needed.

For information about how to build a library in the IDE, see the *IDE Project* Management and Building Guide for ARM.

INVOCATION SYNTAX

The invocation syntax for the archive builder is:

iarchive parameters

Parameters

The parameters are:

Parameter	Description
command	Command line options that define an operation to be performed. Such an option must be specified before the name of the library file.
libraryfile	The library file to be operated on.
objectfile1 objectfileN	The object file(s) that the specified command operates on.
options	Command line options that define actions to be performed. These options can be placed anywhere on the command line.

Table 43: iarchive parameters

Examples

This example creates a library file called mylibrary.a from the source object files module1.o, module2.o, and module3.o:

iarchive mylibrary.a module1.o module2.o module3.o.

This example lists the contents of mylibrary.a:

iarchive --toc mylibrary.a

This example replaces module3.o in the library with the content in the module3.o file and appends module4.o to mylibrary.a:

iarchive --replace mylibrary.a module3.o module4.o

SUMMARY OF IARCHIVE COMMANDS

This table summarizes the iarchive commands:

Command line option	Description
create	Creates a library that contains the listed object files.
delete, -d	Deletes the listed object files from the library.
extract, -x	Extracts the listed object files from the library.
replace, -r	Replaces or appends the listed object files to the library.
symbols	Lists all symbols defined by files in the library.
toc, -t	Lists all files in the library.

Table 44: iarchive commands summary

For more information, see Descriptions of options, page 451.

SUMMARY OF IARCHIVE OPTIONS

This table summarizes the iarchive options:

Command line option	Description
-f	Extends the command line.
output, -o	Specifies the library file.
silent	Sets silent operation.
verbose, -V	Reports all performed operations.

Table 45: iarchive options summary

For more information, see Descriptions of options, page 451.

DIAGNOSTIC MESSAGES

This section lists the messages produced by iarchive:

La001: could not open file filename

iarchive failed to open an object file.

La002: illegal path pathname

The path *pathname* is not a valid path.

La006: too many parameters to cmd command

A list of object modules was specified as parameters to a command that only accepts a single library file.

La007: too few parameters to cmd command

A command that takes a list of object modules was issued without the expected modules.

La008: lib is not a library file

The library file did not pass a basic syntax check. Most likely the file is not the intended library file.

La009: lib has no symbol table

The library file does not contain the expected symbol information. The reason might be that the file is not the intended library file, or that it does not contain any ELF object modules.

La010: no library parameter given

The tool could not identify which library file to operate on. The reason might be that a library file has not been specified.

La011: file file already exists

The file could not be created because a file with the same name already exists.

La013: file confusions, lib given as both library and object

The library file was also mentioned in the list of object modules.

La014: module module not present in archive lib

The specified object module could not be found in the archive.

La015: internal error

The invocation triggered an unexpected error in iarchive.

Ms003: could not open file filename for writing

iarchive failed to open the archive file for writing. Make sure that it is not write protected.

Ms004: problem writing to file filename

An error occurred while writing to file *filename*. A possible reason for this is that the volume is full.

Ms005: problem closing file filename

An error occurred while closing the file filename.

The IAR ELF Tool—ielftool

The IAR ELF Tool, ielftool, can generate a checksum on specific ranges of memories. This checksum can be compared with a checksum calculated on your application.

The source code for ielftool and a Microsoft VisualStudio 2005 template project are available in the arm\src\elfutils directory. If you have specific requirements for how the checksum should be generated or requirements for format conversion, you can modify the source code accordingly.

INVOCATION SYNTAX

The invocation syntax for the IAR ELF Tool is:

ielftool [options] inputfile outputfile [options]

The ielftool tool will first process all the fill options, then it will process all the checksum options (from left to right).

Parameters

The parameters are:

Parameter	Description
inputfile	An absolute ELF executable image produced by the ILINK linker.
options	Any of the available command line options, see Summary of ielftool options, page 441.
outputfile	An absolute ELF executable image.
T 11 16 1 10 1	

Table 46: ielftool parameters

For information about specifying a filename or a directory, see *Rules for specifying a filename or directory as parameters*, page 228.

Example

This example fills a memory range with 0xFF and then calculates a checksum on the same range:

SUMMARY OF IELFTOOL OPTIONS

This table summarizes the ielftool command line options:

Command line option	Description
bin	Sets the format of the output file to binary.
checksum	Generates a checksum.
fill	Specifies fill requirements.
ihex	Sets the format of the output file to linear Intel hex.
self_reloc	Not for general use.
silent	Sets silent operation.
simple	Sets the format of the output file to Simple code.
srec	Sets the format of the output file to Motorola S-records.

Table 47: ielftool options summary

Command line option	Description	
srec-len	Restricts the number of data bytes in each S-record.	
srec-s3only	Restricts the S-record output to contain only a subset of records.	
strip	Removes debug information.	
verbose, -V	Prints all performed operations.	
m 11 /m / 10 / 1	(a)	

Table 47: ielftool options summary (Continued)

For more information, see Descriptions of options, page 451.

The IAR ELF Dumper for ARM—ielfdumparm

The IAR ELF Dumper for ARM, ielfdumparm, can be used for creating a text representation of the contents of a relocatable or absolute ELF file.

ielfdumparm can be used in one of three ways:

- To produce a listing of the general properties of the input file and the ELF segments and ELF sections it contains. This is the default behavior when no command line options are used.
- To also include a textual representation of the contents of each ELF section in the input file. To specify this behavior, use the command line option --all.
- To produce a textual representation of selected ELF sections from the input file. To specify this behavior, use the command line option --section.

INVOCATION SYNTAX

The invocation syntax for ielfdumparm is:

ielfdumparm input_file [output_file]

Note: ielfdumparm is a command line tool which is not primarily intended to be used in the IDE.

Parameters

The parameters are:

Parameter	Description	
input_file	An ELF relocatable or executable file to use as input.	
output_file	A file or directory where the output is emitted. If absent and no	
	output option is specified, output is directed to the console.	

Table 48: ielfdumparm parameters

For information about specifying a filename or a directory, see *Rules for specifying a filename or directory as parameters*, page 228.

SUMMARY OF IELFDUMPARM OPTIONS

This table summarizes the ielfdumparm command line options:

Command line option	Description	
all	Generates output for all input sections regardless of their names or numbers.	
code	Dumps all sections that contain executable code.	
-f	Extends the command line.	
output, -o	Specifies an output file.	
no_strtab	Suppresses dumping of string table sections.	
raw	Uses the generic hexadecimal/ASCII output format for the contents of any selected section, instead of any dedicated output format for that section.	
section, -s	Generates output for selected input sections.	

Table 49: ielfdumparm options summary

For more information, see Descriptions of options, page 451.

The IAR ELF Object Tool—iobjmanip

Use the IAR ELF Object Tool, iobjmanip, to perform low-level manipulation of ELF object files.

INVOCATION SYNTAX

The invocation syntax for the IAR ELF Object Tool is:

iobjmanip options inputfile outputfile

Parameters

The parameters are:

Parameter	Description	
options	Command line options that define actions to be performed. These	
	options can be placed anywhere on the command line. At least one of	
	the options must be specified.	
inputfile	A relocatable ELF object file.	

Table 50: iobjmanip parameters

Parameter	Description	
outputfile	A relocatable ELF object file with all the requested operations applied.	

Table 50: iobjmanip parameters (Continued)

For information about specifying a filename or a directory, see *Rules for specifying a filename or directory as parameters*, page 228.

Examples

This example renames the section .example in input.o to .example2 and stores the result in output.o:

iobjmanip --rename_section .example=.example2 input.o output.o

SUMMARY OF IOBJMANIP OPTIONS

This table summarizes the iobjmanip options:

Command line option	Description	
-f	Extends the command line.	
remove_section	Removes a section.	
rename_section	Renames a section.	
rename_symbol	Renames a symbol.	
strip	Removes debug information.	

Table 51: iobjmanip options summary

For more information, see Descriptions of options, page 451.

DIAGNOSTIC MESSAGES

This section lists the messages produced by iobjmanip:

Lm001: No operation given

None of the command line parameters specified an operation to perform.

Lm002: Expected nr parameters but got nr

Too few or too many parameters. Check invocation syntax for iobjmanip and for the used command line options.

Lm003: Invalid section/symbol renaming pattern pattern

The pattern does not define a valid renaming operation.

Lm004: Could not open file filename

iobjmanip failed to open the input file.

Lm005: ELF format error msg

The input file is not a valid ELF object file.

Lm006: Unsupported section type nr

The object file contains a section that iobjmanip cannot handle. This section will be ignored when generating the output file.

Lm007: Unknown section type nr

iobjmanip encountered an unrecognized section. iobjmanip will try to copy the content as is.

Lm008: Symbol symbol has unsupported format

iobjmanip encountered a symbol that cannot be handled. iobjmanip will ignore this symbol when generating the output file.

Lm009: Group type nr not supported

iobjmanip only supports groups of type GRP_COMDAT. If any other group type is encountered, the result is undefined.

Lm010: Unsupported ELF feature in file: msg

The input file uses a feature that iobjmanip does not support.

Lm011: Unsupported ELF file type

The input file is not a relocatable object file.

Lm012: Ambiguous rename for section/symbol name (alt1 and alt2)

An ambiguity was detected while renaming a section or symbol. One of the alternatives will be used.

Lm013: Section name removed due to transitive dependency on name

A section was removed as it depends on an explicitly removed section.

Lm014: File has no section with index nr

A section index, used as a parameter to --remove_section or --rename_section, did not refer to a section in the input file.

Ms003: could not open file filename for writing

iobjmanip failed to open the output file for writing. Make sure that it is not write protected.

Ms004: problem writing to file filename

An error occurred while writing to file *filename*. A possible reason for this is that the volume is full.

Ms005: problem closing file filename

An error occurred while closing the file filename.

The IAR Absolute Symbol Exporter—isymexport

The IAR Absolute Symbol Exporter, isymexport, can export absolute symbols from a ROM image file, so that they can be used when you link an add-on application.

INVOCATION SYNTAX

The invocation syntax for the IAR Absolute Symbol Exporter is:

isymexport [options] inputfile outputfile [options]

Parameters

The parameters are:

Parameter	Description	
inputfile	A ROM image in the form of an executable ELF file (output from linking).	
options	Any of the available command line options, see Summary of isymexport options, page 447.	
outputfile	A relocatable ELF file that can be used as input to linking, and which contains all or a selection of the absolute symbols in the input file. The output file contains only the symbols, not the actual code or data sections. A steering file can be used to control which symbols that are included, and also to rename some of the symbols if that is desired.	

Table 52: ielftool parameters

For information about specifying a filename or a directory, see *Rules for specifying a filename or directory as parameters*, page 228.

SUMMARY OF ISYMEXPORT OPTIONS

This table summarizes the isymexport command line options:

Command line option	Description	
edit	Specifies a steering file.	
-f	Extends the command line.	
ram_reserve_ranges	Generates symbols to reserve the areas in RAM that the image uses.	
reserve_ranges	Generates symbols to reserve the areas in ROM and RAM that the image uses.	

Table 53: isymexport options summary

For more information, see Descriptions of options, page 451.

STEERING FILES

A steering file can be used for controlling which symbols that are included, and also to rename some of the symbols if that is desired. In the file, you can use show and hide directives to select which public symbols from the input file that are to be included in the output file. rename directives can be used for changing the names of symbols in the input file.

Syntax

The following syntax rules apply:

- Each directive is specified on a separate line.
- C comments (/*...*/) and C++ comments (//...) can be used.
- Patterns can contain wildcard characters that match more than one possible character in a symbol name.
- The * character matches any sequence of zero or more characters in a symbol name.
- The ? character matches any single character in a symbol name.

Example

```
rename xxx_* as YYY_* /*Change symbol prefix from xxx_ to YYY_ */
show YYY_* /* Export all symbols from YYY package */
hide *_internal /* But do not export internal symbols */
show zzz? /* Export zzza, but not zzzaaa */
hide zzzx /* But do not export zzzx */
```

Show directive

Syntax	show pattern	
Parameters	pattern A pattern to match against a symbol name.	
Description	A symbol with a name that matches the pattern will be included in the output file unless this is overridden by a later hide directive.	
Example	/* Include all public symbols ending in _pub. */ show *_pub	

Hide directive

Syntax	hide pattern	
Parameters	pattern	A pattern to match against a symbol name.
Description	A symbol with a name that matches the pattern will not be included in the output file unless this is overridden by a later show directive.	
Example	/* Do not inclu hide *_sys	de public symbols ending in _sys. */

Rename directive

Syntax	rename pattern1	pattern2
Parameters	pattern1	A pattern used for finding symbols to be renamed. The pattern can contain no more than one * or ? wildcard character.
	pattern2	A pattern used for the new name for a symbol. If the pattern contains a wildcard character, it must be of the same kind as in <i>pattern1</i> .
Description	Use this directive to rename symbols from the output file to the input file. No exported symbol is allowed to match more than one rename pattern.	

rename directives can be placed anywhere in the steering file, but they are executed before any show and hide directives. Thus, if a symbol will be renamed, all show and hide directives in the steering file must refer to the new name.

If the name of a symbol matches a *pattern1* pattern that contains no wildcard characters, the symbol will be renamed *pattern2* in the output file.

If the name of a symbol matches a *pattern1* pattern that contains a wildcard character, the symbol will be renamed *pattern2* in the output file, with part of the name matching the wildcard character preserved.

DIAGNOSTIC MESSAGES

This section lists the messages produced by isymexport:

Es001: could not open file filename

isymexport failed to open the specified file.

Es002: illegal path pathname

The path *pathname* is not a valid path.

Es003: format error: message

A problem occurred while reading the input file.

Es004: no input file

No input file was specified.

Es005: no output file

An input file, but no output file was specified.

Es006: too many input files

More than two files were specified.

Es007: input file is not an ELF executable

The input file is not an ELF executable file.

Es008: unknown directive: directive

The specified directive in the steering file is not recognized.

Es009: unexpected end of file

The steering file ended when more input was required.

Es010: unexpected end of line

A line in the steering file ended before the directive was complete.

Es011: unexpected text after end of directive

There is more text on the same line after the end of a steering file directive.

Es012: expected text

The specified text was not present in the steering file, but must be present for the directive to be correct.

Es013: pattern can contain at most one * or ?

Each pattern in the current directive can contain at most one * or one ? wildcard character.

Es014: rename patterns have different wildcards

Both patterns in the current directive must contain exactly the same kind of wildcard. That is, both must either contain:

- No wildcards
- Exactly one *
- Exactly one ?

This error occurs if the patterns are not the same in this regard.

Es014: ambiguous pattern match: *symbol* matches more than one rename pattern

A symbol in the input file matches more than one rename pattern.

Descriptions of options

This section gives detailed reference information about each command line option available for the different utilities.

all					
	Syntax		all		
	ТооІ		ielfdumparm Use this option to include the contents of all ELF sections in the output, in addition to the general properties of the input file. Sections are output in index order, except that each relocation section is output immediately after the section it holds relocations for.		
	Description				
			By default, no sectio	n contents are included in the output.	
		X	This option is not available in the IDE.		
bin					
	Syntax		bin ielftool		
	Tool				
	Description		Sets the format of the output file to binary.		
		VP	To set related options, choose:		
		99	Project>Options>Output converter		
ch	ecksum				
	Syntax		checksum {symbol[+offset] address}:size,algorithm[:[1 2][m][L W][r][i p]] [,start];range[;range]		
	Parameters		symbol	The name of the symbol where the checksum value should be stored. Note that it must exist in the symbol table in the input ELF file.	

offset	An offset to the symbol.
address	The absolute address where the checksum value should be stored.
size	The number of bytes in the checksum: 1, 2, or 4; must not be larger than the size of the checksum symbol.
algorithm	The checksum algorithm used, one of:
	• sum, a byte-wise calculated arithmetic sum. The result is truncated to 8 bits.
	• sum8wide, a byte-wise calculated arithmetic sum. The result is truncated to the size of the symbol.
	• sum32, a word-wise (32 bits) calculated arithmetic sum
	• crc16, CRC16 (generating polynomial 0x11021); used by default
	• crc32, CRC32 (generating polynomial 0x104C11DB7)
	• crc=n, CRC with a generating polynomial of n.
1 2	If specified, can be one of:
	• 1 - Specifies one's complement.
	• 2 - Specifies two's complement.
m	Reverses the order of the bits within each byte when calculating the checksum.
L W	Specifies the size of the unit for which a checksum should be calculated. Typically, this is useful to make the linker produce the same checksum as some hardware CRC implementations that calculate a checksum for more than 8 bits per iteration. Choose between:
	L, calculates a checksum on 32 bits in every iteration
	w, calculates a checksum on 16 bits in every iteration.
	If you do not specify a unit size, 8 bits will be used by default. Using these parameters does not add any additional error detection power to the checksum.
r	Reverses the byte order of the input data within each word of size <i>size</i> .

	i p	Use either i or p, if the <i>start</i> value is bigger than 0. If specified, can be one of:
		• i – Initializes the checksum value with the start value.
		• p - Prefixes the input data with a word of size <i>size</i> that contains the <i>start</i> value.
	start	By default, the initial value of the checksum is 0. If necessary, use <i>start</i> to supply a different initial value. If not 0, then either i or p must be specified.
	range	The address range on which the checksum should be calculated. Hexadecimal and decimal notation is allowed (for example, 0x8002-0x8FFF).
Tool	ielftool	
Description	ranges. The checks symbol will be gene calculated checksur	alculate a checksum with the specified algorithm for the specified um will then replace the original value in <i>symbol</i> . A new absolute erated; with the <i>symbol</i> name suffixed with _value containing the n. This symbol can be used for accessing the checksum value later xample during debugging.
	evaluated from left	option is used more than once on the command line, the options are to right. If a checksum is calculated for a <i>symbol</i> that is specified in checksum option, an error is issued.
Example	This example show address range 0x80	s how to use the crc16 algorithm with the start value 0 over the $00-0x8FFF$:
		cksum=checksum:2,crc16;0x8000-0x8FFF destinationFile.out
	2 bytes will be store	d from sourceFile.out, and the resulting checksum value of size ed at the symbolchecksum. The modified ELF file is saved as e.out leaving sourceFile.out untouched.
See also	Checksum calculati	ion, page 187
Y	To set related option	ns, choose:
	Project>Options>	Linker>Checksum

ELF

--create

Syntax		create libraryfile objectfile1 objectfileN		
Parameters		libraryfile	The library file that the command operates on. For information about specifying a filename, see <i>Rules for specifying a filename</i> or directory as parameters, page 228.	
		objectfile1 objectfileN	The object file(s) to build the library from.	
ТооІ		iarchive		
Description		Use this command to build a new library from a set of object files (modules). The ob files are added to the library in the exact order that they are specified on the commu- line.		
		If no command is sp	pecified on the command line,create is used by default.	
	X	This option is not av	vailable in the IDE.	

--delete, -d

Syntax	delete libraryfile objectfile1 objectfileN -d libraryfile objectfile1 objectfileN		
Parameters	<i>libraryfile</i> The library file that the command opera about specifying a filename, see <i>Rules for directory as parameters</i> , page 228.		
	<pre>objectfile1 The object file(s) that the command ope objectfileN</pre>	rates on.	
Tool	iarchive		
Description	Use this command to remove object files (modules) from an ex files that are specified on the command line will be removed f	<i>c</i> , ,	
	This option is not available in the IDE.		

--edit

Syntax		edit steering_file
Tool		isymexport
Description		Use this option to specify a steering file to control which symbols that are included in the isymexport output file, and also to rename some of the symbols if that is desired.
See also		Steering files, page 447.
1	X	This option is not available in the IDE.

--extract, -x

•		
Syntax		ryfile [objectfile1 objectfileN] [objectfile1 objectfileN]
Parameters	libraryfile	The library file that the command operates on. For information about specifying a filename, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.
	objectfile1 objectfileN	The object file(s) that the command operates on.
Tool	iarchive	
Description	object files is specifi	extract object files (modules) from an existing library. If a list of ed, only these files are extracted. If a list of object files is not files in the library are extracted.
×	This option is not av	ailable in the IDE.

-f

Syntax		-f filename
Parameters		For information about specifying a filename, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.
Tool		iarchive, ielfdumparm, iobjmanip, and isymexport.
Description		Use this option to make the tool read command line options from the named file, with the default filename extension xcl .
		In the command file, you format the items exactly as if they were on the command line itself, except that you can use multiple lines, because the newline character acts just as a space or tab character.
		Both C and C++ style comments are allowed in the file. Double quotes behave in the same way as in the Microsoft Windows command line environment.
	X	This option is not available in the IDE.

--fill

Syntax	fill pattern;range[;range]	
Parameters	range	Specifies the address range for the fill. Hexadecimal and decimal notation is allowed (for example, 0x8002-0x8FFF). Note that each address must be 4-byte aligned.
	pattern	A hexadecimal string with the $0 \times \text{prefix}$ (for example, $0 \times \text{EF}$) interpreted as a sequence of bytes, where each pair of digits corresponds to one byte (for example 0×123456 , for the sequence of bytes 0×12 , 0×34 , and 0×56). This sequence is repeated over the fill area. If the length of the fill pattern is greater than 1 byte, it is repeated as if it started at address 0.
Applicability	ielftool	
Description	Use this option to fill all gaps in one or more ranges with a pattern, which can be either an expression or a hexadecimal string. The contents will be calculated as if the fill pattern was repeatedly filled from the start address until the end address is passed, and then the real contents will overwrite that pattern.	
	If thefill option overlap each other.	n is used more than once on the command line, the fill ranges cannot
Y	To set related option	is, choose:
	Project>Options>I	.inker>Checksum

--ihex

Syntax		ihex
Tool		ielftool
Description		Sets the format of the output file to linear Intel hex.
	X	To set related options, choose:
		Project>Options>Linker>Output converter

	_		
no	_strtab		
	Syntax		no_strtab
	Tool		ielfdumparm
	Description		Use this option to suppress dumping of string table sections (sections of type SHT_STRTAB).
		X	This option is not available in the IDE.
ou	tput, -o		
	Syntax		-o {filename directory} output {filename directory}
	Parameters		For information about specifying a filename or a directory, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.
	Tool		iarchive and ielfdumparm.
	Description		iarchive
			By default, iarchive assumes that the first argument after the iarchive command is the name of the destination library. Use this option to explicitly specify a different filename for the library.
			ielfdumparm
			By default, output from the dumper is directed to the console. Use this option to direct the output to a file instead. The default name of the output file is the name of the input file with an added id filename extension
			You can also specify the output file by specifying a file or directory following the name of the input file.
		SP.	This option is not available in the IDE.



--ram_reserve_ranges

Syntax	ram_reserve_ranges[=symbol_prefix]
Parameters	symbol_prefix The prefix of symbols created by this option.
ТооІ	isymexport
Description	Use this option to generate symbols for the areas in RAM that the image uses. One symbol will be generated for each such area. The name of each symbol is based on the name of the area and is prefixed by the optional parameter <i>symbol_prefix</i> .
	Generating symbols that cover an area in this way prevents the linker from placing other content at the affected addresses. This can be useful when linking against an existing image.
	Ifram_reserve_ranges is used together withreserve_ranges, the RAM areas will get their prefix from theram_reserve_ranges option and the non-RAM areas will get their prefix from thereserve_ranges option.
See also	reserve_ranges, page 462.
	This option is not available in the IDE.

--raw

Syntax		raw
Tool		ielfdumparm
Description		By default, many ELF sections will be dumped using a text format specific to a particular kind of section. Use this option to dump each selected ELF section using the generic text format.
		The generic text format dumps each byte in the section in hexadecimal format, and where appropriate, as ASCII text.
	X	This option is not available in the IDE.

--remove_section

Syntax		remove_section {section number}	
Parameters		section	The section—or sections, if there are more than one section with the same name—to be removed.
		number	The number of the section to be removed. Section numbers can be obtained from an object dump created using ielfdumparm.
ТооІ		iobjmanip	
Description		Use this option to make iobjmanip omit the specified section when generating the output file.	
	X	This option is not a	vailable in the IDE.

--rename_section

Syntax		rename_section {oldname oldnumber}=newname	
Parameters		oldname The section—or sections, if there are more than one section with the same name—to be renamed.	
		oldnumber	The number of the section to be renamed. Section numbers can be obtained from an object dump created using ielfdumparm.
		newname	The new name of the section.
Tool		iobjmanip	
Description		Use this option to make iobjmanip rename the specified section when generating the output file.	
	X	This option is not available in the IDE.	

--rename_symbol

Syntax	rename_symbol	oldname =newname
Parameters	oldname newname	The symbol to be renamed. The new name of the symbol.
Tool	iobjmanip	
Description	Use this option to m output file.	ake iobjmanip rename the specified symbol when generating the
J.	This option is not av	vailable in the IDE.

--replace, -r

Syntax	-	replace libraryfile objectfile1 objectfileN -r libraryfile objectfile1 objectfileN	
Parameters	libraryfile	The library file that the command operates on. For information about specifying a filename, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.	
	objectfile1 objectfileN	The object file(s) that the command operates on.	
ТооІ	iarchive		
Description	object files specifie	to replace or add object files (modules) to an existing library. The d on the command line either replace existing object files in the e the same name) or are appended to the library.	
	This option is not a	vailable in the IDE.	

--reserve_ranges

Syntax		reserve_ranges[=symbol_prefix]	
Parameters		symbol_prefix The prefix of symbols created by this option.	
Tool		isymexport	
Description		Use this option to generate symbols for the areas in ROM and RAM that the image uses. One symbol will be generated for each such area. The name of each symbol is based on the name of the area and is prefixed by the optional parameter <i>symbol_prefix</i> .	
		Generating symbols that cover an area in this way prevents the linker from placing other content at the affected addresses. This can be useful when linking against an existing image.	
		Ifreserve_ranges is used together withram_reserve_ranges, the RAM areas will get their prefix from theram_reserve_ranges option and the non-RAM areas will get their prefix from thereserve_ranges option.	
See also		ram_reserve_ranges, page 459.	
	X	This option is not available in the IDE.	

--section, -s

Syntax	<pre>section section_number section_name[,]s section_number section_name[,]</pre>		
Parameters	section_numberThe number of the section to be dumped.section_nameThe name of the section to be dumped.		
ТооІ	ielfdumparm		
Description	Use this option to dump the contents of a section with the specified number, or any section with the specified name. If a relocation section is associated with a selected section, its contents are output as well.		
	If you use this option, the general properties of the input file will not be included in the output.		

You can specify multiple section numbers or names by separating them with commas, or by using this option more than once.

By default, no section contents are included in the output.

Example	-s 3,17	/* Sections #3 and #17
	-s .debug_frame,42	<pre>/* Any sections named .debug_frame and also section #42 */</pre>
		,



1

This option is not available in the IDE.

--self_reloc

Syntax	self_reloc	
Tool	ielftool	
Description	This option is intentionally not documented as it is not intended for This option is not available in the IDE.	r general use.

--silent

Syntax		silent -S (iarchive only)
Tool		iarchive and ielftool.
Description		Causes the tool to operate without sending any messages to the standard output stream. By default, ielftool sends various messages via the standard output stream. You can use this option to prevent this. ielftool sends error and warning messages to the error output stream, so they are displayed regardless of this setting.
	SP.	This option is not available in the IDE.

sir	nple	
	Syntax	simple
	ТооІ	ielftool
	Description	Sets the format of the output file to Simple code.
	X	To set related options, choose:
		Project>Options>Output converter
sre	ec	
	Syntax	srec
	ТооІ	ielftool
	Description	Sets the format of the output file to Motorola S-records.
	X	To set related options, choose:
		Project>Options>Output converter
sre	ec-len	
	Syntax	srec-len=length
	Parameters	<i>length</i> The number of data bytes in each S-record.
	ТооІ	ielftool
	Description	Restricts the number of data bytes in each S-record. This option can be used in combination with thesrec option.
	X	This option is not available in the IDE.

--srec-s3only

	/		
	Syntax		srec-s3only
	Tool		ielftool
	Description		Restricts the S-record output to contain only a subset of records, that is S0, S3 and S7 records. This option can be used in combination with thesrec option.
		X	This option is not available in the IDE.
stri	p		
	Syntax		strip

Description Use this option to remove all sections containing debug information before the output file is written.

Note that ielftool needs an unstripped input ELF image. If you use the --strip option in the linker, remove it and use the --strip option in ielftool instead.



To set related options, choose:

iobjmanip and ielftool.

Project>Options>Linker>Output>Include debug information in output

--symbols

Tool

-

Syntax	symbols libr	symbols libraryfile		
Parameters	libraryfile	The library file that the command operates on. For information about specifying a filename, see <i>Rules for specifying a filename</i> or directory as parameters, page 228.		
Tool	iarchive			
Description	Use this command to list all external symbols that are defined by any object file (module) in the specified library, together with the name of the object file (module) that defines it.			

In silent mode (--silent), this command performs symbol table-related syntax checks on the library file and displays only errors and warnings.



This option is not available in the IDE.

--toc, -t

Syntax	toc libraryfile -t libraryfile	
Parameters	libraryfile	The library file that the command operates on. For information about specifying a filename, see <i>Rules for specifying a filename or directory as parameters</i> , page 228.
ТооІ	iarchive	
Description	Use this command to list the names of all object files (modules) in a specified library. In silent mode (silent), this command performs basic syntax checks on the library file, and displays only errors and warnings.	
Y	This option is not available in the IDE.	

--verbose, -V

Syntax	-verbose V (iarchive only)	
Tool	archive and ielftool.	
Description	ription Use this option to make the tool report which operations it giving diagnostic messages.	
	his option is not available in the IDE bed	cause this setting is always enabled.

Implementation-defined behavior for Standard C

This chapter describes how the compiler handles the implementation-defined areas of the C language based on Standard C.

Note: The IAR Systems implementation adheres to a freestanding implementation of Standard C. This means that parts of a standard library can be excluded in the implementation.

If you are using C89 instead of Standard C, see *Implementation-defined behavior* for C89, page 483. For a short overview of the differences between Standard C and C89, see *C language overview*, page 161.

Descriptions of implementation-defined behavior

This section follows the same order as the C standard. Each item includes references to the ISO chapter and section (in parenthesis) that explains the implementation-defined behavior.

J.3.1 TRANSLATION

Diagnostics (3.10, 5.1.1.3)

Diagnostics are produced in the form:

filename,linenumber level[tag]: message

where filename is the name of the source file in which the error was encountered, linenumber is the line number at which the compiler detected the error, level is the level of seriousness of the message (remark, warning, error, or fatal error), tag is a unique tag that identifies the message, and message is an explanatory message, possibly several lines.

White-space characters (5.1.1.2)

At translation phase three, each non-empty sequence of white-space characters is retained.

J.3.2 ENVIRONMENT

The character set (5.1.1.2)

The source character set is the same as the physical source file multibyte character set. By default, the standard ASCII character set is used. However, if you use the --enable_multibytes compiler option, the host character set is used instead.

Main (5.1.2.1)

The function called at program startup is called main. No prototype is declared for main, and the only definition supported for main is:

```
int main(void)
```

To change this behavior, see Customizing system initialization, page 117.

The effect of program termination (5.1.2.1)

Terminating the application returns the execution to the startup code (just after the call to main).

Alternative ways to define main (5.1.2.2.1)

There is no alternative ways to define the main function.

The argv argument to main (5.1.2.2.1)

The argv argument is not supported.

Streams as interactive devices (5.1.2.3)

The streams stdin, stdout, and stderr are treated as interactive devices.

Signals, their semantics, and the default handling (7.14)

In the DLIB library, the set of supported signals is the same as in Standard C. A raised signal will do nothing, unless the signal function is customized to fit the application.

Signal values for computational exceptions (7.14.1.1)

In the DLIB library, there are no implementation-defined values that correspond to a computational exception.

Signals at system startup (7.14.1.1)

In the DLIB library, there are no implementation-defined signals that are executed at system startup.

Environment names (7.20.4.5)

In the DLIB library, there are no implementation-defined environment names that are used by the getenv function.

The system function (7.20.4.6)

The system function is not supported.

J.3.3 IDENTIFIERS

Multibyte characters in identifiers (6.4.2)

Additional multibyte characters may not appear in identifiers.

Significant characters in identifiers (5.2.4.1, 6.1.2)

The number of significant initial characters in an identifier with or without external linkage is guaranteed to be no less than 200.

J.3.4 CHARACTERS

Number of bits in a byte (3.6)

A byte contains 8 bits.

Execution character set member values (5.2.1)

The values of the members of the execution character set are the values of the ASCII character set, which can be augmented by the values of the extra characters in the host character set.

Alphabetic escape sequences (5.2.2)

The standard alphabetic escape sequences have the values a-7, b-8, f-12, n-10, r-13, t-9, and v-11.

Characters outside of the basic executive character set (6.2.5)

A character outside of the basic executive character set that is stored in a char is not transformed.

Plain char (6.2.5, 6.3.1.1)

A plain char is treated as an unsigned char.

Source and execution character sets (6.4.4.4, 5.1.1.2)

The source character set is the set of legal characters that can appear in source files. By default, the source character set is the standard ASCII character set. However, if you use the command line option --enable_multibytes, the source character set will be the host computer's default character set.

The execution character set is the set of legal characters that can appear in the execution environment. By default, the execution character set is the standard ASCII character set.

However, if you use the command line option --enable_multibytes, the execution character set will be the host computer's default character set. The IAR DLIB Library needs a multibyte character scanner to support a multibyte execution character set. See *Locale*, page 123.

Integer character constants with more than one character (6.4.4.4)

An integer character constant that contains more than one character will be treated as an integer constant. The value will be calculated by treating the leftmost character as the most significant character, and the rightmost character as the least significant character, in an integer constant. A diagnostic message will be issued if the value cannot be represented in an integer constant.

Wide character constants with more than one character (6.4.4.4)

A wide character constant that contains more than one multibyte character generates a diagnostic message.

Locale used for wide character constants (6.4.4.4)

By default, the C locale is used. If the --enable_multibytes compiler option is used, the default host locale is used instead.

Locale used for wide string literals (6.4.5)

By default, the C locale is used. If the --enable_multibytes compiler option is used, the default host locale is used instead.

Source characters as executive characters (6.4.5)

All source characters can be represented as executive characters.

J.3.5 INTEGERS

Extended integer types (6.2.5)

There are no extended integer types.

Range of integer values (6.2.6.2)

The representation of integer values are in the two's complement form. The most significant bit holds the sign; 1 for negative, 0 for positive and zero.

For information about the ranges for the different integer types, see *Basic data types*, page 298.

The rank of extended integer types (6.3.1.1)

There are no extended integer types.

Signals when converting to a signed integer type (6.3.1.3)

No signal is raised when an integer is converted to a signed integer type.

Signed bitwise operations (6.5)

Bitwise operations on signed integers work the same way as bitwise operations on unsigned integers; in other words, the sign-bit will be treated as any other bit.

J.3.6 FLOATING POINT

Accuracy of floating-point operations (5.2.4.2.2)

The accuracy of floating-point operations is unknown.

Rounding behaviors (5.2.4.2.2)

There are no non-standard values of FLT_ROUNDS.

Evaluation methods (5.2.4.2.2)

There are no non-standard values of FLT_EVAL_METHOD.

Converting integer values to floating-point values (6.3.1.4)

When an integral value is converted to a floating-point value that cannot exactly represent the source value, the round-to-nearest rounding mode is used (FLT_ROUNDS is defined to 1).

Converting floating-point values to floating-point values (6.3.1.5)

When a floating-point value is converted to a floating-point value that cannot exactly represent the source value, the round-to-nearest rounding mode is used (FLT_ROUNDS is defined to 1).

Denoting the value of floating-point constants (6.4.4.2)

The round-to-nearest rounding mode is used (FLT_ROUNDS is defined to 1).

Contraction of floating-point values (6.5)

Floating-point values are contracted. However, there is no loss in precision and because signaling is not supported, this does not matter.

Default state of FENV_ACCESS (7.6.1)

The default state of the pragma directive FENV_ACCESS is OFF.

Additional floating-point mechanisms (7.6, 7.12)

There are no additional floating-point exceptions, rounding-modes, environments, and classifications.

Default state of FP_CONTRACT (7.12.2)

The default state of the pragma directive FP_CONTRACT is OFF.

J.3.7 ARRAYS AND POINTERS

Conversion from/to pointers (6.3.2.3)

For information about casting of data pointers and function pointers, see *Casting*, page 306.

ptrdiff_t (6.5.6)

For information about ptrdiff_t, see *ptrdiff_t*, page 306.

J.3.8 HINTS

Honoring the register keyword (6.7.1)

User requests for register variables are not honored.

Inlining functions (6.7.4)

User requests for inlining functions increases the chance, but does not make it certain, that the function will actually be inlined into another function. See *Inlining functions*, page 65.

J.3.9 STRUCTURES, UNIONS, ENUMERATIONS, AND BITFIELDS

Sign of 'plain' bitfields (6.7.2, 6.7.2.1)

For information about how a 'plain' int bitfield is treated, see Bitfields, page 300.

Possible types for bitfields (6.7.2.1)

All integer types can be used as bitfields in the compiler's extended mode, see -e, page 242.

Bitfields straddling a storage-unit boundary (6.7.2.1)

A bitfield is always placed in one—and one only—storage unit, which means that the bitfield cannot straddle a storage-unit boundary.

Allocation order of bitfields within a unit (6.7.2.1)

For information about how bitfields are allocated within a storage unit, see *Bitfields*, page 300.

Alignment of non-bitfield structure members (6.7.2.1)

The alignment of non-bitfield members of structures is the same as for the member types, see *Alignment*, page 297.

Integer type used for representing enumeration types (6.7.2.2)

The chosen integer type for a specific enumeration type depends on the enumeration constants defined for the enumeration type. The chosen integer type is the smallest possible.

J.3.10 QUALIFIERS

Access to volatile objects (6.7.3)

Any reference to an object with volatile qualified type is an access, see *Declaring objects volatile*, page 309.

J.3.11 PREPROCESSING DIRECTIVES

Mapping of header names (6.4.7)

Sequences in header names are mapped to source file names verbatim. A backslash '\' is not treated as an escape sequence. See *Overview of the preprocessor*, page 381.

Character constants in constant expressions (6.10.1)

A character constant in a constant expression that controls conditional inclusion matches the value of the same character constant in the execution character set.

The value of a single-character constant (6.10.1)

A single-character constant may only have a negative value if a plain character (char) is treated as a signed character, see *--char_is_signed*, page 235.

Including bracketed filenames (6.10.2)

For information about the search algorithm used for file specifications in angle brackets <>, see *Include file search procedure*, page 221.

Including quoted filenames (6.10.2)

For information about the search algorithm used for file specifications enclosed in quotes, see *Include file search procedure*, page 221.

Preprocessing tokens in #include directives (6.10.2)

Preprocessing tokens in an #include directive are combined in the same way as outside an #include directive.

Nesting limits for #include directives (6.10.2)

There is no explicit nesting limit for #include processing.

Universal character names (6.10.3.2)

Universal character names (UCN) are not supported.

Recognized pragma directives (6.10.6)

In addition to the pragma directives described in the chapter *Pragma directives*, the following directives are recognized and will have an indeterminate effect. If a pragma directive is listed both in the *Pragma directives* chapter and here, the information provided in the *Pragma directives* chapter overrides the information here.

```
alignment
baseaddr
basic_template_matching
building_runtime
can instantiate
```

codeseg constseg cspy_support dataseg define_type_info do_not_instantiate early_dynamic_initialization function function_effects hdrstop important_typedef instantiate keep_definition library_default_requirements library_provides library_requirement_override memory module_name no_pch once public_equ system_include vector warnings

Default __DATE__ and __TIME__ (6.10.8)

The definitions for __TIME__ and __DATE__ are always available.

J.3.12 LIBRARY FUNCTIONS

Additional library facilities (5.1.2.1)

Most of the standard library facilities are supported. Some of them—the ones that need an operating system—require a low-level implementation in the application. For more information, see *The DLIB runtime environment*, page 99.

Diagnostic printed by the assert function (7.2.1.1)

The assert() function prints:

filename: linenr expression -- assertion failed

when the parameter evaluates to zero.

Representation of the floating-point status flags (7.6.2.2)

For information about the floating-point status flags, see fenv.h, page 395.

Feraiseexcept raising floating-point exception (7.6.2.3)

For information about the feraiseexcept function raising floating-point exceptions, see *Floating-point environment*, page 304.

Strings passed to the setlocale function (7.11.1.1)

For information about strings passed to the setlocale function, see Locale, page 123.

Types defined for float_t and double_t (7.12)

The FLT_EVAL_METHOD macro can only have the value 0.

Domain errors (7.12.1)

No function generates other domain errors than what the standard requires.

Return values on domain errors (7.12.1)

Mathematic functions return a floating-point NaN (not a number) for domain errors.

Underflow errors (7.12.1)

Mathematic functions set errno to the macro ERANGE (a macro in errno.h) and return zero for underflow errors.

fmod return value (7.12.10.1)

The fmod function returns a floating-point NaN when the second argument is zero.

The magnitude of remquo (7.12.10.3)

The magnitude is congruent modulo INT_MAX.

signal() (7.14.1.1)

The signal part of the library is not supported.

Note: Low-level interface functions exist in the library, but will not perform anything. Use the template source code to implement application-specific signal handling. See *Signal and raise*, page 127.

NULL macro (7.17)

The NULL macro is defined to 0.

Terminating newline character (7.19.2)

stdout stream functions recognize either newline or end of file (EOF) as the terminating character for a line.

Space characters before a newline character (7.19.2)

Space characters written to a stream immediately before a newline character are preserved.

Null characters appended to data written to binary streams (7.19.2)

No null characters are appended to data written to binary streams.

File position in append mode (7.19.3)

The file position is initially placed at the beginning of the file when it is opened in append-mode.

Truncation of files (7.19.3)

Whether a write operation on a text stream causes the associated file to be truncated beyond that point, depends on the application-specific implementation of the low-level file routines. See *File input and output*, page 123.

File buffering (7.19.3)

An open file can be either block-buffered, line-buffered, or unbuffered.

A zero-length file (7.19.3)

Whether a zero-length file exists depends on the application-specific implementation of the low-level file routines.

Legal file names (7.19.3)

The legality of a filename depends on the application-specific implementation of the low-level file routines.

Number of times a file can be opened (7.19.3)

Whether a file can be opened more than once depends on the application-specific implementation of the low-level file routines.

Multibyte characters in a file (7.19.3)

The encoding of multibyte characters in a file depends on the application-specific implementation of the low-level file routines.

remove() (7.19.4.1)

The effect of a remove operation on an open file depends on the application-specific implementation of the low-level file routines. See *File input and output*, page 123.

rename() (7.19.4.2)

The effect of renaming a file to an already existing filename depends on the application-specific implementation of the low-level file routines. See *File input and output*, page 123.

Removal of open temporary files (7.19.4.3)

Whether an open temporary file is removed depends on the application-specific implementation of the low-level file routines.

Mode changing (7.19.5.4)

freopen closes the named stream, then reopens it in the new mode. The streams stdin, stdout, and stderr can be reopened in any new mode.

Style for printing infinity or NaN (7.19.6.1, 7.24.2.1)

The style used for printing infinity or NaN for a floating-point constant is inf and nan (INF and NAN for the F conversion specifier), respectively. The n-char-sequence is not used for nan.

%p in printf() (7.19.6.1, 7.24.2.1)

The argument to a %p conversion specifier, print pointer, to printf() is treated as having the type void *. The value will be printed as a hexadecimal number, similar to using the %x conversion specifier.

Reading ranges in scanf (7.19.6.2, 7.24.2.1)

A - (dash) character is always treated as a range symbol.

%p in scanf (7.19.6.2, 7.24.2.2)

The p conversion specifier, scan pointer, to scanf() reads a hexadecimal number and converts it into a value with the type void *.

File position errors (7.19.9.1, 7.19.9.3, 7.19.9.4)

On file position errors, the functions fgetpos, ftell, and fsetpos store EFPOS in errno.

An n-char-sequence after nan (7.20.1.3, 7.24.4.1.1)

An n-char-sequence after a NaN is read and ignored.

errno value at underflow (7.20.1.3, 7.24.4.1.1)

errno is set to ERANGE if an underflow is encountered.

Zero-sized heap objects (7.20.3)

A request for a zero-sized heap object will return a valid pointer and not a null pointer.

Behavior of abort and exit (7.20.4.1, 7.20.4.4)

A call to abort () or _Exit() will not flush stream buffers, not close open streams, and not remove temporary files.

Termination status (7.20.4.1, 7.20.4.3, 7.20.4.4)

The termination status will be propagated to __exit() as a parameter. exit() and _Exit() use the input parameter, whereas abort uses EXIT_FAILURE.

The system function return value (7.20.4.6)

The system function is not supported.

The time zone (7.23.1)

The local time zone and daylight savings time must be defined by the application. For more information, see *Time*, page 127.

Range and precision of time (7.23)

The implementation uses signed long for representing clock_t and time_t, based at the start of the year 1970. This gives a range of approximately plus or minus 69 years in seconds. However, the application must supply the actual implementation for the functions time and clock. See *Time*, page 127.

clock() (7.23.2.1)

The application must supply an implementation of the clock function. See *Time*, page 127.

%Z replacement string (7.23.3.5, 7.24.5.1)

By default, ":" is used as a replacement for %Z. Your application should implement the time zone handling. See *Time*, page 127.

Math functions rounding mode (F.9)

The functions in math.h honor the rounding direction mode in FLT-ROUNDS.

J.3.13 ARCHITECTURE

Values and expressions assigned to some macros (5.2.4.2, 7.18.2, 7.18.3)

There are always 8 bits in a byte.

MB_LEN_MAX is at the most 6 bytes depending on the library configuration that is used.

For information about sizes, ranges, etc for all basic types, see *Data representation*, page 297.

The limit macros for the exact-width, minimum-width, and fastest minimum-width integer types defined in stdint.h have the same ranges as char, short, int, long, and long long.

The floating-point constant FLT_ROUNDS has the value 1 (to nearest) and the floating-point constant FLT_EVAL_METHOD has the value 0 (treat as is).

The number, order, and encoding of bytes (6.2.6.1)

See Data representation, page 297.

The value of the result of the size of operator (6.5.3.4)

See Data representation, page 297.

J.4 LOCALE

Members of the source and execution character set (5.2.1)

By default, the compiler accepts all one-byte characters in the host's default character set. If the compiler option --enable_multibytes is used, the host multibyte characters are accepted in comments and string literals as well.

The meaning of the additional character set (5.2.1.2)

Any multibyte characters in the extended source character set is translated verbatim into the extended execution character set. It is up to your application with the support of the library configuration to handle the characters correctly.

Shift states for encoding multibyte characters (5.2.1.2)

Using the compiler option --enable_multibytes enables the use of the host's default multibyte characters as extended source characters.

Direction of successive printing characters (5.2.2)

The application defines the characteristics of a display device.

The decimal point character (7.1.1)

The default decimal-point character is a '.'. You can redefine it by defining the library configuration symbol _LOCALE_DECIMAL_POINT.

Printing characters (7.4, 7.25.2)

The set of printing characters is determined by the chosen locale.

Control characters (7.4, 7.25.2)

The set of control characters is determined by the chosen locale.

Characters tested for (7.4.1.2, 7.4.1.3, 7.4.1.7, 7.4.1.9, 7.4.1.10, 7.4.1.11, 7.25.2.1.2, 7.25.5.1.3, 7.25.2.1.7, 7.25.2.1.9, 7.25.2.1.10, 7.25.2.1.11)

The sets of characters tested are determined by the chosen locale.

The native environment (7.1.1.1)

The native environment is the same as the "C" locale.

Subject sequences for numeric conversion functions (7.20.1, 7.24.4.1)

There are no additional subject sequences that can be accepted by the numeric conversion functions.

The collation of the execution character set (7.21.4.3, 7.24.4.4.2)

The collation of the execution character set is determined by the chosen locale.

Message returned by strerror (7.21.6.2)

The messages returned by the strerror function depending on the argument is:

Argument	Message
EZERO	no error
EDOM	domain error
ERANGE	range error
EFPOS	file positioning error
EILSEQ	multi-byte encoding error
<0 >99	unknown error
all others	error nnn

Table 54: Message returned by strerror()—IAR DLIB library

Implementation-defined behavior for C89

This chapter describes how the compiler handles the implementation-defined areas of the C language based on the C89 standard.

If you are using Standard C instead of C89, see Implementation-defined behavior for Standard C, page 467. For a short overview of the differences between Standard C and C89, see C language overview, page 161.

Descriptions of implementation-defined behavior

The descriptions follow the same order as the ISO appendix. Each item covered includes references to the ISO chapter and section (in parenthesis) that explains the implementation-defined behavior.

TRANSLATION

Diagnostics (5.1.1.3)

Diagnostics are produced in the form:

filename,linenumber level[tag]: message

where filename is the name of the source file in which the error was encountered, linenumber is the line number at which the compiler detected the error, level is the level of seriousness of the message (remark, warning, error, or fatal error), tag is a unique tag that identifies the message, and message is an explanatory message, possibly several lines.

ENVIRONMENT

Arguments to main (5.1.2.2.2.1)

The function called at program startup is called main. No prototype was declared for main, and the only definition supported for main is:

int main(void)

To change this behavior for the IAR DLIB runtime environment, see *Customizing system initialization*, page 117.

Interactive devices (5.1.2.3)

The streams stdin and stdout are treated as interactive devices.

IDENTIFIERS

Significant characters without external linkage (6.1.2)

The number of significant initial characters in an identifier without external linkage is 200.

Significant characters with external linkage (6.1.2)

The number of significant initial characters in an identifier with external linkage is 200.

Case distinctions are significant (6.1.2)

Identifiers with external linkage are treated as case-sensitive.

CHARACTERS

Source and execution character sets (5.2.1)

The source character set is the set of legal characters that can appear in source files. The default source character set is the standard ASCII character set. However, if you use the command line option --enable_multibytes, the source character set will be the host computer's default character set.

The execution character set is the set of legal characters that can appear in the execution environment. The default execution character set is the standard ASCII character set. However, if you use the command line option --enable_multibytes, the execution character set will be the host computer's default character set. The IAR DLIB Library needs a multibyte character scanner to support a multibyte execution character set.

See Locale, page 123.

Bits per character in execution character set (5.2.4.2.1)

The number of bits in a character is represented by the manifest constant CHAR_BIT. The standard include file limits.h defines CHAR_BIT as 8.

Mapping of characters (6.1.3.4)

The mapping of members of the source character set (in character and string literals) to members of the execution character set is made in a one-to-one way. In other words, the same representation value is used for each member in the character sets except for the escape sequences listed in the ISO standard.

Unrepresented character constants (6.1.3.4)

The value of an integer character constant that contains a character or escape sequence not represented in the basic execution character set or in the extended character set for a wide character constant generates a diagnostic message, and will be truncated to fit the execution character set.

Character constant with more than one character (6.1.3.4)

An integer character constant that contains more than one character will be treated as an integer constant. The value will be calculated by treating the leftmost character as the most significant character, and the rightmost character as the least significant character, in an integer constant. A diagnostic message will be issued if the value cannot be represented in an integer constant.

A wide character constant that contains more than one multibyte character generates a diagnostic message.

Converting multibyte characters (6.1.3.4)

The only locale supported—that is, the only locale supplied with the IAR C/C++ Compiler—is the 'C' locale. If you use the command line option --enable_multibytes, the IAR DLIB Library will support multibyte characters if you add a locale with multibyte support or a multibyte character scanner to the library.

See Locale, page 123.

Range of 'plain' char (6.2.1.1)

A 'plain' char has the same range as an unsigned char.

INTEGERS

Range of integer values (6.1.2.5)

The representation of integer values are in the two's complement form. The most significant bit holds the sign; 1 for negative, 0 for positive and zero.

See *Basic data types*, page 298, for information about the ranges for the different integer types.

Demotion of integers (6.2.1.2)

Converting an integer to a shorter signed integer is made by truncation. If the value cannot be represented when converting an unsigned integer to a signed integer of equal length, the bit-pattern remains the same. In other words, a large enough value will be converted into a negative value.

Signed bitwise operations (6.3)

Bitwise operations on signed integers work the same way as bitwise operations on unsigned integers; in other words, the sign-bit will be treated as any other bit.

Sign of the remainder on integer division (6.3.5)

The sign of the remainder on integer division is the same as the sign of the dividend.

Negative valued signed right shifts (6.3.7)

The result of a right-shift of a negative-valued signed integral type preserves the sign-bit. For example, shifting 0xFF00 down one step yields 0xFF80.

FLOATING POINT

Representation of floating-point values (6.1.2.5)

The representation and sets of the various floating-point numbers adheres to IEEE 854–1987. A typical floating-point number is built up of a sign-bit (s), a biased exponent (e), and a mantissa (m).

See *Floating-point types*, page 304, for information about the ranges and sizes for the different floating-point types: float and double.

Converting integer values to floating-point values (6.2.1.3)

When an integral number is cast to a floating-point value that cannot exactly represent the value, the value is rounded (up or down) to the nearest suitable value.

Demoting floating-point values (6.2.1.4)

When a floating-point value is converted to a floating-point value of narrower type that cannot exactly represent the value, the value is rounded (up or down) to the nearest suitable value.

ARRAYS AND POINTERS

size_t (6.3.3.4, 7.1.1)

See *size_t*, page 306, for information about size_t.

Conversion from/to pointers (6.3.4)

See *Casting*, page 306, for information about casting of data pointers and function pointers.

ptrdiff_t (6.3.6, 7.1.1)

See *ptrdiff_t*, page 306, for information about the ptrdiff_t.

REGISTERS

Honoring the register keyword (6.5.1)

User requests for register variables are not honored.

STRUCTURES, UNIONS, ENUMERATIONS, AND BITFIELDS

Improper access to a union (6.3.2.3)

If a union gets its value stored through a member and is then accessed using a member of a different type, the result is solely dependent on the internal storage of the first member.

Padding and alignment of structure members (6.5.2.1)

See the section *Basic data types*, page 298, for information about the alignment requirement for data objects.

Sign of 'plain' bitfields (6.5.2.1)

A 'plain' int bitfield is treated as an unsigned int bitfield. All integer types are allowed as bitfields.

Allocation order of bitfields within a unit (6.5.2.1)

Bitfields are allocated within an integer from least-significant to most-significant bit.

Can bitfields straddle a storage-unit boundary (6.5.2.1)

Bitfields cannot straddle a storage-unit boundary for the chosen bitfield integer type.

Integer type chosen to represent enumeration types (6.5.2.2)

The chosen integer type for a specific enumeration type depends on the enumeration constants defined for the enumeration type. The chosen integer type is the smallest possible.

QUALIFIERS

Access to volatile objects (6.5.3)

Any reference to an object with volatile qualified type is an access.

DECLARATORS

Maximum numbers of declarators (6.5.4)

The number of declarators is not limited. The number is limited only by the available memory.

STATEMENTS

Maximum number of case statements (6.6.4.2)

The number of case statements (case values) in a switch statement is not limited. The number is limited only by the available memory.

PREPROCESSING DIRECTIVES

Character constants and conditional inclusion (6.8.1)

The character set used in the preprocessor directives is the same as the execution character set. The preprocessor recognizes negative character values if a 'plain' character is treated as a signed character.

Including bracketed filenames (6.8.2)

For file specifications enclosed in angle brackets, the preprocessor does not search directories of the parent files. A parent file is the file that contains the #include directive. Instead, it begins by searching for the file in the directories specified on the compiler command line.

Including quoted filenames (6.8.2)

For file specifications enclosed in quotes, the preprocessor directory search begins with the directories of the parent file, then proceeds through the directories of any grandparent files. Thus, searching begins relative to the directory containing the source file currently being processed. If there is no grandparent file and the file is not found, the search continues as if the filename was enclosed in angle brackets.

Character sequences (6.8.2)

Preprocessor directives use the source character set, except for escape sequences. Thus, to specify a path for an include file, use only one backslash:

#include "mydirectory\myfile"

Within source code, two backslashes are necessary:

file = fopen("mydirectory\\myfile","rt");

Recognized pragma directives (6.8.6)

In addition to the pragma directives described in the chapter *Pragma directives*, the following directives are recognized and will have an indeterminate effect. If a pragma directive is listed both in the *Pragma directives* chapter and here, the information provided in the *Pragma directives* chapter overrides the information here.

alignment baseaddr basic_template_matching building_runtime can instantiate codeseg constseg cspy_support dataseg define_type_info do_not_instantiate early_dynamic_initialization function function effects hdrstop important_typedef instantiate keep_definition library_default_requirements library_provides library_requirement_override memory module_name no_pch once

public_equ system_include vector warnings

Default __DATE__ and __TIME__ (6.8.8)

The definitions for __TIME__ and __DATE__ are always available.

IAR DLIB LIBRARY FUNCTIONS

The information in this section is valid only if the runtime library configuration you have chosen supports file descriptors. See the chapter *The DLIB runtime environment* for more information about runtime library configurations.

NULL macro (7.1.6)

The NULL macro is defined to 0.

Diagnostic printed by the assert function (7.2)

The assert() function prints: filename:linenr expression -- assertion failed when the parameter evaluates to zero.

Domain errors (7.5.1)

NaN (Not a Number) will be returned by the mathematic functions on domain errors.

Underflow of floating-point values sets errno to ERANGE (7.5.1)

The mathematics functions set the integer expression errno to ERANGE (a macro in errno.h) on underflow range errors.

fmod() functionality (7.5.6.4)

If the second argument to fmod() is zero, the function returns NaN; errno is set to EDOM.

signal() (7.7.1.1)

The signal part of the library is not supported.

Note: Low-level interface functions exist in the library, but will not perform anything. Use the template source code to implement application-specific signal handling. See *Signal and raise*, page 127.

Terminating newline character (7.9.2)

stdout stream functions recognize either newline or end of file (EOF) as the terminating character for a line.

Blank lines (7.9.2)

Space characters written to the stdout stream immediately before a newline character are preserved. There is no way to read the line through the stdin stream that was written through the stdout stream.

Null characters appended to data written to binary streams (7.9.2)

No null characters are appended to data written to binary streams.

Files (7.9.3)

Whether the file position indicator of an append-mode stream is initially positioned at the beginning or the end of the file, depends on the application-specific implementation of the low-level file routines.

Whether a write operation on a text stream causes the associated file to be truncated beyond that point, depends on the application-specific implementation of the low-level file routines. See *File input and output*, page 123.

The characteristics of the file buffering is that the implementation supports files that are unbuffered, line buffered, or fully buffered.

Whether a zero-length file actually exists depends on the application-specific implementation of the low-level file routines.

Rules for composing valid file names depends on the application-specific implementation of the low-level file routines.

Whether the same file can be simultaneously open multiple times depends on the application-specific implementation of the low-level file routines.

remove() (7.9.4.1)

The effect of a remove operation on an open file depends on the application-specific implementation of the low-level file routines. See *File input and output*, page 123.

rename() (7.9.4.2)

The effect of renaming a file to an already existing filename depends on the application-specific implementation of the low-level file routines. See *File input and output*, page 123.

%p in printf() (7.9.6.1)

The argument to a p conversion specifier, print pointer, to printf() is treated as having the type void *. The value will be printed as a hexadecimal number, similar to using the x conversion specifier.

%p in scanf() (7.9.6.2)

The p conversion specifier, scan pointer, to scanf() reads a hexadecimal number and converts it into a value with the type void *.

Reading ranges in scanf() (7.9.6.2)

A - (dash) character is always treated as a range symbol.

File position errors (7.9.9.1, 7.9.9.4)

On file position errors, the functions fgetpos and ftell store EFPOS in errno.

Message generated by perror() (7.9.10.4)

The generated message is:

usersuppliedprefix:errormessage

Allocating zero bytes of memory (7.10.3)

The calloc(), malloc(), and realloc() functions accept zero as an argument. Memory will be allocated, a valid pointer to that memory is returned, and the memory block can be modified later by realloc.

Behavior of abort() (7.10.4.1)

The abort() function does not flush stream buffers, and it does not handle files, because this is an unsupported feature.

Behavior of exit() (7.10.4.3)

The argument passed to the exit function will be the return value returned by the main function to cstartup.

Environment (7.10.4.4)

The set of available environment names and the method for altering the environment list is described in *Environment interaction*, page 126.

system() (7.10.4.5)

How the command processor works depends on how you have implemented the system function. See *Environment interaction*, page 126.

Message returned by strerror() (7.11.6.2)

The messages returned by strerror() depending on the argument is:

Argument	Message
EZERO	no error
EDOM	domain error
ERANGE	range error
EFPOS	file positioning error
EILSEQ	multi-byte encoding error
<0 >99	unknown error
all others	error nnn

Table 55: Message returned by strerror()—IAR DLIB library

The time zone (7.12.1)

The local time zone and daylight savings time implementation is described in *Time*, page 127.

clock() (7.12.2.1)

From where the system clock starts counting depends on how you have implemented the clock function. See *Time*, page 127.

Descriptions of implementation-defined behavior

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